

From about June 18 to Aug 17, 1956. I attended the Dartmouth summer research project on Artificial Intelligence. About 56 pages were done in this notebook - these later became a ~~the~~ Multilist report of ≈ 55 pp. Also || with this work, a notebook with pages numbered [(83 eq.)] was kept. The pp. of this N.B. with ~~the~~ are referred to e.g. by (83).16. This N.B. has 104 pp. and goes from -1) to 102). Also an attempt at a prelim. report was written - this was ≈ 20 pp.

PP
 ≈ 46 eq
 ≈ 85 incl.
 12 only
39 pp

So 175 pp. written these 2 months; Rather large. The research project wasn't very suggestive.

The main things ~~are~~ of value

- 1) Write and get report reproduced (very imp.)
- 2) Met some interesting people in this field.
- 3) Got idea of how poor most that in this field is.
- 4) Some ideas:
 - a) Search problem may be imp.
 - b) These guys may eventually invent a T.M.,

Simply by working more and more interesting special problems. Simon and Newell; Minsky: best candidates - ^{Trench}More is a question mark.

- 5) Interested some of the people in T.M.

- May be able to get it programmed on 709 or IAS computer.

A bunch of random, imp. ideas on T.M.: Clean them up later.

.03 1) If we have $3x_i + 2y_i = z_i$ and this works for a few values of i , (x_i, y_i, z_i) being some n-tuple, then it should be possible to realize more quickly that $3[1] + 2[2] = [3]$ is a sort of thing that one can plug a relation $(1, 2, 3)$ into directly, without having to go back to

$$(3x_i + 2y_i = z_i) = [1|2|3|4|5|6|7] \times (3, x_i, +, 2, y_i, =, z_i)$$

- i.e. thru th. set. There should be a more direct way to retain $3, +, 2, =$, in their proper positions. How to do it: see 289.01

.21

Consider $x_1 + 2x_2 = 3x_3$. For (several values of i gn. as examples)

$$= [1|2|3|4|5|6] \times (x_i, +, 2, =, 3)$$

We must give "2" a much higher priority than th. set of all integers — similarly with "3". Otherwise th. predictions for $x + 2x = \square\square$ would be

$\infty x, \dots, 1x, 2x, 3x, \dots, \infty x.$

In general in this sort of thing, th. priority of a set must be $<$ that of any of its elements components. This can be assured if one has th. special operation, in which / a set of objects can be x'fnd into a set of individual objects; each individual U, and each

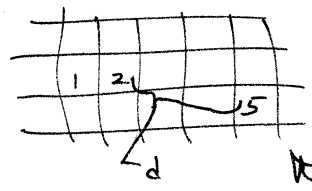
T.M.

individual $U_{appri} \rightarrow$ th. U_{appri} of th. set of objects.

(2)

A ~~Mode~~ Mode by which I would like NMTM to work in changing U's of str, ntp, ngm, in time. If an ~~abs~~ abs (\equiv pugm) has been useful for a period of time, then it ~~is~~ is not useful for a long time — then I would want its U to go down rather slowly, since abs's that have once been useful, tend to be useful again. — even tho they may be rather un-useful for a while.

15 On th. "closeness" rule for str. It is felt that the distance "d" should be a major contributor to th. decrement of U_{appri} with distance, but that th. difference betw. 5 and 2 should be something irrelevant. I.e. th. 5 (skipping 3 and 4) can



be performed by permuting, repeating or omitting operations on th. ntp that th. str. operates on. — it should be done before th. str. is used.

Essentially, then, a str. may consist of only a set of ~~no~~ squares with integers in them, but no repetitions or omissions.

It is, perhaps, possl. to get along with binary str only. E.g., to get

$x_1 \oplus x_1 \oplus x_1$, one would first use

$$\boxed{1|1|2} \text{ on } (x_1, x_1), \text{ then } \rightarrow \boxed{1|1|2} \times (x_1, \boxed{1|1|2} \times (x_1, x_1)).$$

A trouble mite be experienced with getting th. ngm $x_i \oplus y_i \oplus z_i$ from th. ntp (x_i, y_i, z_i) 96.06

10. | from 287.21 we can use $N_1 = 3B + 2Z$

and $N_2 = 2X + 3Y + Z$

$$\begin{bmatrix} 1 & 2 \end{bmatrix} \times (N_1, N_2) = 3X + 2Y + Z ; \quad 3B + 2Z = 2X + 3Y + Z$$

This ambiguity is rather irrelevant and disconcerting. — perhaps this idea of str. isn't so hot.

A rather strange, tho perhaps useful way to get these substitutions: multiply (x_i, y_i, z_i) by

the set $\begin{bmatrix} 1 & 2 & 3 \end{bmatrix}$, $\begin{bmatrix} 1 & 2 & 3 \end{bmatrix}$, $\begin{bmatrix} 1 & 2 & 3 \end{bmatrix}$, etc.

— i.e. the set of all ^{ternary} linear str., in which we have the order 1, 2, 3. we get an ~~in~~ N_i^0

~~we can multiply~~ which we can use with

$$N_1 = (3B + 2Z) = 1N$$

$\begin{bmatrix} 1 & 2 \end{bmatrix} \times (N_1, N_i^0) =$ all poss. substitutions of x_i, y_i, z_i in that order. Unfortunately, there will be ~~too~~ many irrelevant ~~substitutions~~ u, v, w resulting, but maybe they can be elim. somehow.

Actually, if one were given the ~~problem~~ $q.e.d.$:

$$3X_3 + 2Y_3 = \begin{bmatrix} \quad & \quad \end{bmatrix} \text{ one could give only}$$

~~answers~~ Z_3 or B_3 or Z_3 or B_3 as

correct answers.

Gen. Methodological Notes:

On page ^{Dart N.B.} 99) are a bunch of imp. things that must be done for future T.M. work. Th. foll. seems to be clear: It would be well to work up some good MTM's that used ~~some~~ ~~abstract~~ abstraction sets and hier order sets. Concurrent with this ~~is~~ (before or after) work out details of NMTM.

They work out self-improvement program.

Keep eyes open for kinds of methods that need not be built in - open loop, but would arise when T.M. began trying to improve itself.

Also, work out present T.M. in somewhat greater detail than in Th. report. Use $U_{ij} \rightarrow \frac{c_{ij}}{T}$ and actually ~~compute~~ compute Th. changes in R . U 's, using some simplified approximation method.

List various important groups of problems from 99).

- E.G.
- 1) NMTM
 - 2) Self-improvement T.M.
 - 3) Detailed mechanics of present T.M., with no abs. sets, or ~~sets~~ hier order sets. General improvements.
 - 4) a) abs. sets
b) hier order sets.
 - 5) Physical realization of T.M.
 - 6) General remarks on ^{overall.} T.M. philosophy.

Dart N.B.

↓

T.M.

exceptions over $C_1 - C_2$). The only trouble, is that in using ngrams from C_1 , one is "safe" - but one ~~to~~ would be losing out on specialized ngrams that hold over C_2 only.

A ngram that has been used much over $C_1 - C_2$ and little over C_2 , should be given a rather small U.

An example of ~~is~~ is $\boxed{1}$ in $\mathbb{R}^n =$ "even set" of examples. $\boxed{1}$ is of low U in C_1 , ~~because~~ because of "s", but it is O.K. for C_2 .

Present orientation: That it would be O.K. to use C_1 (averaging over all cases). If there are some ngrams that work over C_2 but not over parts of C_1 , then T.M. should make ~~more~~ more complex ngrams that ~~can~~ apply only to ~~some~~ possl. members of C_2 . This orientation is still tentative.

How to assign "U" is a big question. I think that this should be tentatively done by using C_2 only.

~~Two~~ Two closely related problems:

1) suppose we have $\boxed{10}$ with a pgram count of 10, and 2 "cases". [see α90 and α91 for defs. of "pgram" and "pgram count"] [cases are q. elements (see α89 for def. of "q. element") to which that pgram applies] suppose we have $\boxed{5}$ with a pgram count of 5; and 4 cases. If we are given a q. element with

$\boxed{1}$ in it, how shall we predict?

2) If we have a pgram with $\boxed{1}$ and $\boxed{10}$ in it, how shall we predict $\boxed{1}$? Look at the relative wts. of these 2 pgrams.

Discu: Both of the above ~~can~~ will not happen often. Reason: If a q. element is always followed by the correctly completed element, such the pgrams causing trouble will be ~~soon~~ immediately eliminated. If not, such pgrams should have a high mortality rate. (A pgram dies when a counter example is found.)

T.M.

in R. most recent, extensive work on B.G. - from ~ X889 to X892, and X893 to X898 - also much comment on ~~pages~~ pages following X898. I think f(M) = "aprip" of n. ngmst.

Anyway, say one has ~~pick~~ r examples of this \square ngmst. One ^{also} has v ~~ex~~ examples (v >> r) of a ngmst. averaged over all sgs. One then has $F_1(v, x)$ probability of at least x "good" members of R's first ngmst.

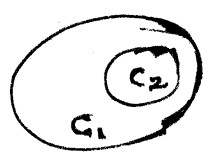
One has $F_2(v, x)$ proby of at least x "good" membs. of R. second ~~ngmst~~ ngmst. How do F_1 and F_2 cross-couple?



The over-all average / ^{case = C1} is at least as reliable as R. strict \square case. $\equiv C_2$ This is because th. only imp. things are counter examples. C_1 has all of C_2 ~~as~~ as special cases, so C_1 has ~~over~~ all C_2 's opportunities for counter examples, plus many extra.

Th. ~~real~~ real q. is whether to ~~pay~~ pay much attention to ~~cases of C2 with no counter~~ ngms that have no ~~counter~~ counter examples / ^{in C2} but do have ~~counter~~ counter examples in C_1 .

It would seem that R. above argument is ~~more~~ most directly applicable to Math T.M. (~~with~~ with no probabilistic elements) - but is probably also applicable to R. more general case.



There seems to be somewhat of a paradox here. Apparently in C_1 , we get a larger sample size than in C_2 . But C_1 ~~is~~ is ~~more~~ at least as reliable for prediction as C_2 is. Th. problem seems to be

that ~~we must have~~ all ~~ngmsts~~ ngmsts. satisfying C_1 are O.K. for C_2 . ~~and~~ In fact, ngmsts that hold over C_1 can be expected to be more reliable than those which hold only over C_2 (~~and~~ and with some