

created a new prgm.

→ these are now pp. 114-115.
Dart. Report

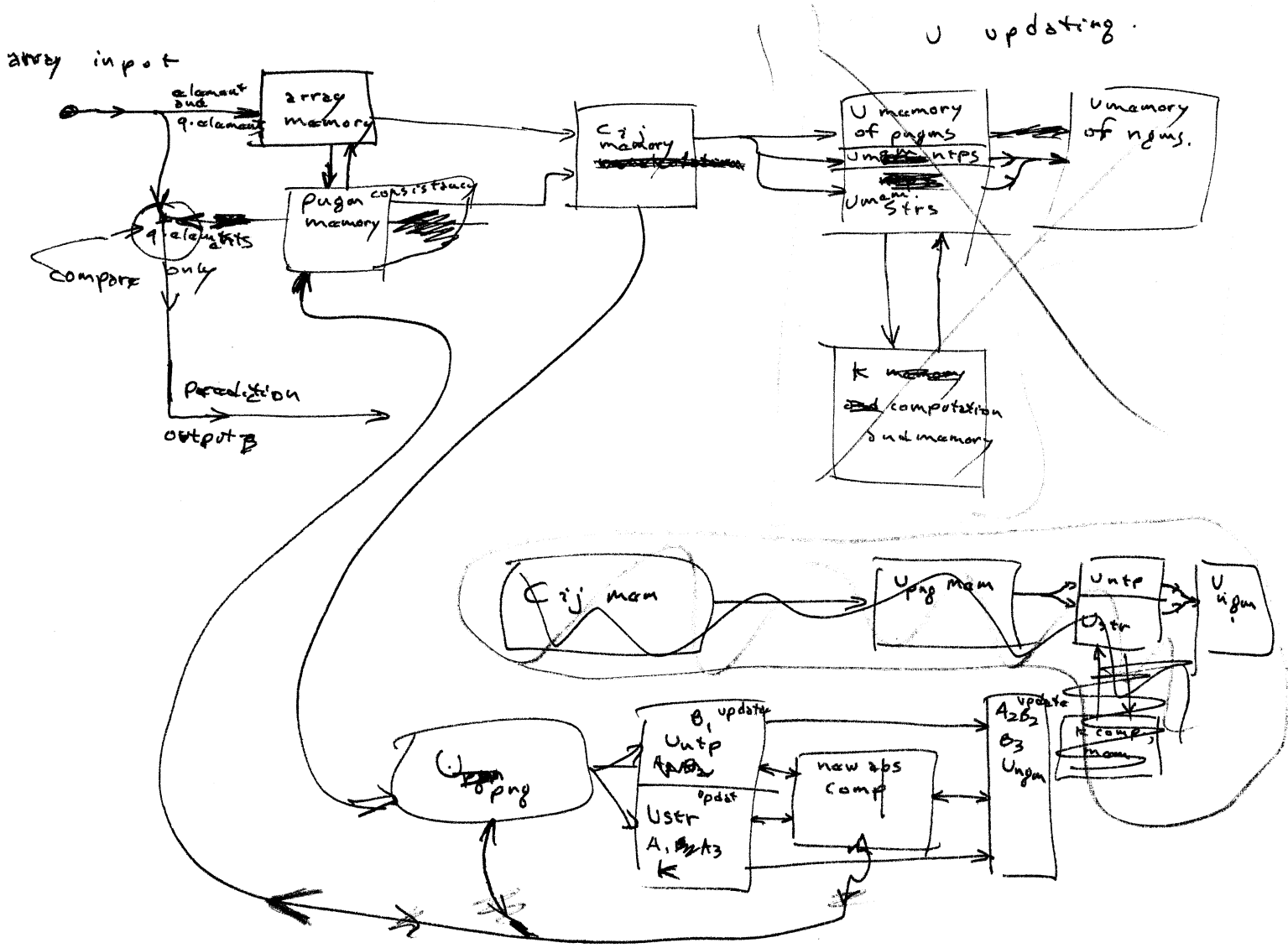
02 in $\alpha 122.06$: If all of the new prgms are inconsistent, do any of the abs. need to be updated? ^{Yes} Well - just what happens when new prgms are created?

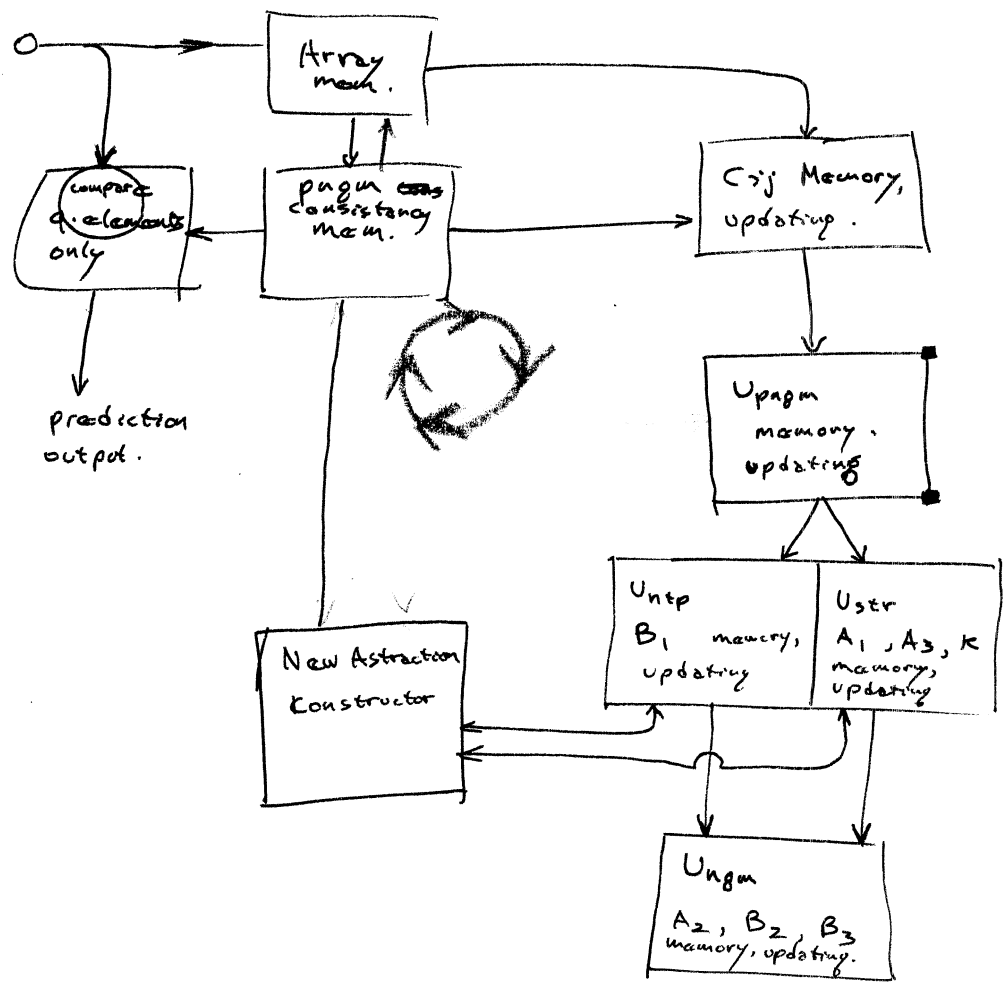
- 1) str x ntp
- 2) . . .
- 3) . . .
- 4) . . .

← (see $\alpha 114$)

Many new abs. are created. Even if the old ~~c_{ij} 's~~ are the same, there are new i 's and j 's - ~~all of the new c_{ij} 's are zero~~ There are new c_{ij} 's - they are all zero. They do, however, effect all of the U 's, since some new U_{ij} 's = 0 - these U_{ij} 's did not, however, exist before.

Block diagram of $\alpha 121, \alpha 118, \alpha 122$





A possl. block diagram. Flow for search for applicable pnm / is red loop. with updating

This flow diagram isn't too clear.

The use of " C_{ij} " in the exposition, for the counts of all pgms ~~is~~ is bad. Some pgms are not expressible as the product of a str. and ntp.

Re hash of T.M. ~~loop~~ operation rules:

1) T.M. gets element:

2) Store in array mem.
 Scan thru ~~pgm~~ mem. to see if any C_{ij} 's have been changed due to ^{an} ~~old~~ pgm becoming inconsistent.

~~Change these bits.
 don't mention possibility of skipping.~~

c) If no C_{ij} 's change stop and wait for new input.
 If ~~see~~ one or more C_{ij} 's change, do $\alpha 118.05$ to .11 and wait for new input.

make into 2 steps

2) T.M. get q. element.

a) store in array mem.
 b) scan thru pgm memory to see if any C_{ij} 's have been changed due to an old pgm becoming inconsistent.

SN

This is unnec., but it will \downarrow the confusion in the loop that has been written. Unnec. because the updating of C_{ij} when a new element comes in is unnec. This can be done when and only when a new q. element comes.

c) search thru pgm mem. for appropriate consistent pgm.
 d) If proper pgm is found, change proper C_{ij} .
 e) If ~~any C_{ij} 's have been changed~~ 2. b) or 2. c), have resulted in any change in C_{ij} , do $\alpha 118.05$ to .11, then go to 2f) directly. If not, go to 2f).

f) If \neq 2. d) yielded suitable proper pgm, stop and wait for new input. If 2. d) did not yield suitable

pages, go to step 2)

g) Construct ^{several} new pages of hy a priori U that ^{are not in the page. mean, that} apply to P.
interrogation sq. in q. Do this by constructing new strs, ngs, utps and combining them, using methods of (2.14.32 ff.)

SN Whenever New abs. are created - whether or not Cij's are changed, we must update all abs. — However, whenever we create new abs., we also ^{always} create new pages - ~~whenever~~ even if their Cij's ~~are~~ may be zero, this does affect ~~all of the pages~~ many of the abs.

The moral is - that the ~~top~~ program of $\alpha 121 \alpha 118 \alpha 122$ is probably O.K., tho not too easy to understand.

SN on "search" problems: Just run thru some problems to see just what kinds of "characters" one uses in directing search. These characters may be of P. kind that either chose ^{new} sub-goals, or tell a reasonable step to try, towards ~~gen.~~ sub-goal. — or, when a gen. sub-goal is gn, tell whether a step should be taken, or a new sub-goal placed. Also time estimates on various paths are imp. "Sequenced analysis" may be imp. — just as a method of directing one's attention to P. rite points.

I don't think this problem is particularly diff. - apart from P. routine of getting new tricks from old.

My approach to the ^{sub-}problem " ~~find~~ construct a ~~new~~ page that fits P. q. element " has been in the spirit of Ashby: i.e. describe ~~all~~ a large no. of possibilities, then use exhaustive search, since the problem is ~~the~~ well defined".

Thurs Aug 9, 1956

(74)

In th. problems in which I applied this method, th. search wasn't long enough to make any difference, but eventually, I will have to get around to th. "search" problem. What I may need is only a quick way of searching, rather than a general method of solving all search problems - tho probably th. latter is required.

Exposition: Mention that most of th. equations are nec. only "in spirit," ~~that extremely difficult~~ to give an idea as to just what ~~variables~~ ^{the} measurable, the various U 's depend upon. In a real machine, only the roughest approx. solns. would be used. Probably powers of 2 are adequate in representing their values.

serious trouble! After doing " = " T.M. finds a certain set of abss. useful. However, as soon as n is introduced, all of th. old pgms become of zero U , since they are inconsistent, and so do all of th. ~~other~~ other abss. ~~that were~~ that are based on th. U 's of th. pgms. This difficulty can be elim. in several ways - a good one is using AIMTM, in which a pgram is not discarded until many counter examples are found.

However, probably the best thing to do is leave it alone, and start from scratch with $n \begin{matrix} 1 & 0 & 0 & \dots \\ 0 & 1 & 0 & \dots \end{matrix}$

Also, explain technique of forming relevant pgms. - by a) starting with a possl. pgram, and trying to construct it by basic operations (the inversion problem)

or b) Iterating all possl. pugms with $U < a$ certain amt. and choosing out th. relevant ones. (Exhaustive search, ruffly in order of $\text{apri } U$.)
 We have a simple clue in a), however, since we have the $\text{apri } U$'s ~~with~~ of th. components. — also if a "useless" pugm is constructed, it may very well be useful later. I.e. this is analogous to inventing sin^x by constructing a table of sin^x . It is practical if sin^x is used very often. In th. present case this ~~is~~ may or may not be so.

This is imp. and interesting: That th. finding of a relevant pugm, of th. best possl. $\text{apri } U$, is a "well defined problem of th. second kind" \rightarrow more exactly, th. "1/2 kind"! an inversion problem — like ~~finding~~ finding a sequence of string x fms operations that proves a theorem. It is a problem in which sub-goals are imp., and in which one has, to some extent, an integrated serving signal, that tells one how by th. U apri is, of th. ~~state~~ present state of x fmn.

28 Also write Introduction, saying that this is not a progress report in th. ord. sense, but merely serves to indicate ~~where~~ in what direction I am working — th. general methodology and approach used. While some details are given, the details are only to make the methods clear. ~~many~~ The details, themselves, are not necessarily those in the present state, ~~will be modified~~ will be modified con and will continue to be considerably modified, as the work progresses.

NUMBER "V.S." sets of sets . . . etc.

It is somewhat likely, that humans store th. concept "3", not as th. set of all triples, but as the result of a "counting" operation.

~~III~~ — Matching objects to th. integers.

— However, this does, correspond to th. equivalence class of all triples — i.e. 2 things equivalent to th. integer 3, must be equivalent to each other.

At any rate, I think that there ought to be a more natural, directly intuitive approach to ~~III~~ all of th. concepts that I want T.M. to learn.

IMP. IDEA: Instead of trying to

get T.M. to ^{invent} ~~discover~~ th. concepts of "negative numbers," "imaginary nos.," etc, give it problems,

→ th. soln. does not directly involve these objects, but could most easily be obtained using them.

Then see if T.M. could ~~be made~~ form the internalized abs. that would yield the ~~these~~ correct answers. I can "look inside" T.M. and see just what ~~ing.~~ seq. is nec. so that it would

tend to form, internally, abs methods equivalent to neg. nos. and im. nos.

I don't think that T.M. is really proceeding properly. What I have done, is taken ~~these~~ some combination and xform methods, that I think are reasonable, and tried to make them do what is required.

What should be done, is to try to trace out my intuitive processes in solving these problems, then find a set of epist. rules that form an ^{adequate} "basis" for them.

General Discussion: From about ~~120~~ 120 to 130 there has been an enormous waste of time and effort that may have to be rewritten.

T.M. should be gn. q. elements like $\begin{matrix} = 1 \\ 1 \end{matrix}, \begin{matrix} \sim 1 \\ 0 \end{matrix}, \begin{matrix} = 50 \\ 50 \end{matrix}, \text{ etc.}$, which make it imp. for it to develop tetragrams, etc., which turn out to be useless, in the long run.

We want to develop strs like $\boxed{1|2}, \boxed{1|1|2}$, $\boxed{2|1|1|2}$, etc., which are useful in the long run.

A disturbing feature, is that about a month ago I realized this difficulty and its "solution", and wrote some about it.

There is the possy. of just continuing in the present direction.

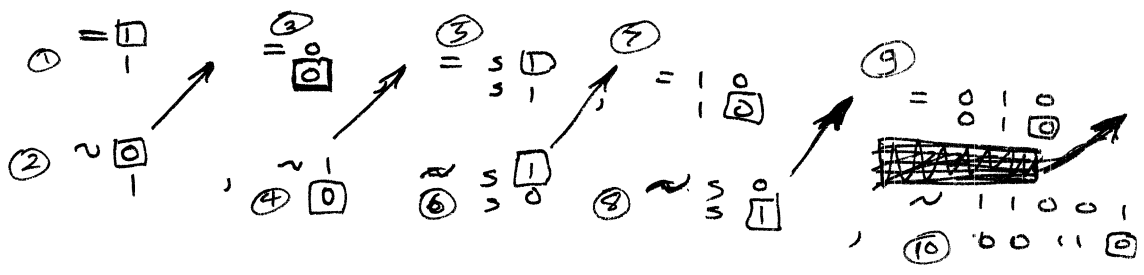
(SN) I certainly don't like this discontinuous behavior, in which a whole bunch of abs suddenly change their 0 values / where ^{to near zero} a few pgrams become inconsistent.

Any way! A poss. plan for Pt. chapter on "examples".

- 1) List the comb. rules and equations that are to be used. If extra equs. are used, indicate that $\exists a, b, c$ are to be added to.
- 2) Make outline of th. tag. seq. ~~then~~ then determine just which pugs, ngms ntps and str become imp.

.10

So first Pt. tag. seq.:



Suppose T.M. get

$$= 001011$$

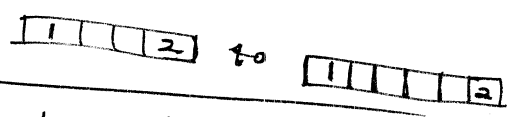
$$001010 \text{ after } = 110$$

$$110, \text{ so it}$$

has to make a rather large "jump". ~~Therefore~~

.25

i.e. from th. str.



On Inversion (cont. and other inversions) to 50.01

IF $str \times N_t = P$ with N_t and P of by U ,
 then $U_{str} = \frac{1}{N} \sum_{i=1}^n U_{p_i} - \text{const.}$ - with no special
 rule for ~~this~~ this inversion. So inversions of
 this sort, are simply a rapid method to find objects
 of by U , not a "really" new method.

Inversions are "Economy Measures".

Possibly Great Idea: Mixing ~~methodology~~ and ~~the~~ methods of forming new words, with R_n words themselves.

Suppose we store the word "3" as a set of 3 objects, plus a set of instructions that tell how to match it with other sets of objects.

The relation "1 greater than" is a set of instructions telling how to get R_n set of instructions for $n+1$ from the set for n .

In general, this is a set of instructions for getting a new object from an old. Expressed extensively, we would simply write all ordered pairs of R_n old and new objects.

In practice - (intensive definition), we would give the rule. An example is the str. S . We can

express it as the \cong relation $(S, N_{t_2}, S \times N_{t_2}) (i = -\infty \dots +\infty)$

or as the method by which one goes from N_{t_2} to $S \times N_{t_2}$.

A set of instructions for getting a new word from an old, is a relation, but it is also a device used to create new ~~words~~ abs. As a relation it is a word - as a device, it is part of the methodology.

This is related, perhaps, to the str. scheme of (67) and (68)

101 from 78). 25 :

① = 0
|

1, 0, etc. tried - N.G.

= 1, 1 tried - (6 poss. cases - all inductive.)

~~scribbled out~~

1.5 = 1 = 1 and 1 now only consistent digits.

How to solve Problems:

- ① state problem.
- ② solve problem or break into sub-problems
- ③ solve sub problems
- or write down next problem
- ④ solve ~~any~~ unsolved problems, break them into sub-problems or write down new problems.
- ⑤ same as ④
- ⑥ same as ⑤

② ~ 0
|

prediction is 1

②.5 ~ 0
|

① is inconsistent.

③ = 0
|

= 1 inconsistent.

0 and 1 are poss. digits: ambiguity, so no prediction. prediction is 1

③.5 = 0
~~scribbled out~~

0 and 1 are inconsistent. there are no consistent digits.

④ ~ 1
|

~ 1 and 1 are consistent

th. prediction is 1

4.5 $\sim \begin{matrix} 1 \\ 0 \end{matrix}$ and $\begin{matrix} 1 \\ 1 \end{matrix}$ are incons.

5

The Tag. seq. on 78.10 is too fast, and has too little info, per element. Perhaps I should make the Tag. seq. fairly ad.hoc, to illustrate the various methods of combination.

① = $\begin{matrix} 0 & 1 \\ 0 & 1 \end{matrix}$ ② = $\begin{matrix} 1 & 0 \\ 1 & 0 \end{matrix}$; $\begin{matrix} 1 \\ 1 \end{matrix}$ is incons.
 $\begin{matrix} 1 \\ 1 \end{matrix}$ is consistent.

\therefore prediction is 1

③ ~~matrix~~ = $\begin{matrix} 1 & 0 \\ 1 & 0 \end{matrix}$ $\begin{matrix} 0 \\ 0 \end{matrix}$ is consistent.
prediction is 0

④ $\sim \begin{matrix} 0 & 0 \\ 1 & 1 \end{matrix}$ $\begin{matrix} 1 \\ 1 \end{matrix}$ is incons.

⑤ $\sim \begin{matrix} 1 & 0 \\ 1 & 1 \end{matrix}$ $\begin{matrix} 1 \\ 1 \end{matrix}$ is cons. prediction is 1

⑥ ~~matrix~~ $\sim \begin{matrix} 1 & 1 & 0 \\ 0 & 0 & 1 \end{matrix}$ $\begin{matrix} 1 \\ 1 \end{matrix}$ is incons. ~~matrix~~ and $\sim \begin{matrix} 1 \\ 1 \end{matrix}$ are cons.
 $\begin{matrix} 1 \\ 0 \end{matrix}$ is cons. ~~matrix~~ is incons. ~~matrix~~
prediction is 1

⑦ = $\begin{matrix} 1 & 1 & 0 & 1 & 0 \\ 1 & 0 & 0 & 1 & 0 \end{matrix}$ no cons. digits.
= $\begin{matrix} 1 \\ 1 \end{matrix}$, = $\begin{matrix} 1 \\ 1 \end{matrix}$, $\begin{matrix} 0 \\ 0 \end{matrix}$
cons cons incons.
prediction is 1



$\begin{matrix} \boxed{1} \\ | \end{matrix}$ and $\begin{matrix} | \\ \boxed{1} \end{matrix}$ are imp. ngs. Also $\begin{matrix} \boxed{2} & \boxed{1} \\ | & | \end{matrix}$ and to some extent, $\begin{matrix} \boxed{1} \\ | \end{matrix}$.

SN It might be well to make U of an ngm. a mean of \mathcal{R}_n pgrams that it could be xformed into. This simplifies updating ~~and~~ a great deal, since there \mathcal{R}_n exists only about 1 imp. updating eqn. for U 's.

⑧ $\sim \begin{matrix} | & 0 & | \\ \square & | & 0 \end{matrix}$ no cons. dpgms.

Use more elements, otherwise there is no way for t.M. to learn about \mathcal{R}_n environs of \sim and $=$.

What about many = examples, with \square randomly placed - than \rightsquigarrow a few \sim elements. than $\sim \begin{matrix} | & 0 \\ \square & | \end{matrix}$ should be O.K. followed by $= \begin{matrix} \square & 0 \\ | & 0 \end{matrix} \dots$

O.K.:

① $= \begin{matrix} | & 0 & | \\ \square & | & 0 \end{matrix}$, ② $= \begin{matrix} 0 & 0 & | \\ 0 & 0 & | \end{matrix}$ ~~is cons.~~ $\begin{matrix} | \\ | \end{matrix}$ is cons. as is $0 \begin{matrix} | \\ | \end{matrix}$ and $0 \begin{matrix} | \\ | \end{matrix}$ prediction is 1

③ $= \begin{matrix} 0 & 0 & | & 0 \\ 0 & \square & | & 0 \end{matrix}$ ~~is cons.~~ $\begin{matrix} 0 \\ 0 \end{matrix}$ is cons. $\begin{matrix} 0 \\ \square \end{matrix}$ is cons. \therefore no pred.

④ $= \begin{matrix} 0 & 0 & | & 1 & 0 \\ 0 & 0 & | & 1 & 0 \end{matrix}$ $\begin{matrix} 0 \\ | \end{matrix}$ is incons. $\begin{matrix} | \\ | \end{matrix}$ is cons. prediction is 1

5

$$= \begin{matrix} \blacksquare & 0 & 0 & 1 \\ \square & 0 & 0 & 1 \end{matrix}$$

\square is cons.
pred. is 1.

6

$$\sim \begin{matrix} 1 & 0 \\ 0 & 1 \end{matrix}$$

7

$$\sim \begin{matrix} 0 & 0 & 1 \\ 1 & 1 & 0 \end{matrix}$$

8

$$\sim \begin{matrix} 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 \end{matrix}$$

9

$$\sim \begin{matrix} 1 & 1 & 0 \\ 0 & 0 & 1 \end{matrix}$$

10

~~10~~

$$\sim \begin{matrix} \square & 1 & \blacksquare \\ 1 & 0 & 0 \end{matrix}$$

no cons. digms. at all

$$\sim \begin{matrix} \square \\ 1 \end{matrix} \text{ is cons,}$$

~~$\begin{matrix} \square & 1 \\ 1 & 0 \end{matrix}$ is cons~~

~~$\begin{matrix} \square & 1 \\ 1 & 0 \end{matrix}$ is cons~~

~~$\begin{matrix} \square & 1 \\ 1 & 0 \end{matrix}$ is cons~~

Prediction is 0.

\square , \square by U digms

11

$$= \begin{matrix} 1 & 0 & 1 & 1 \\ \square & 0 & 1 & 1 \end{matrix}$$

~~$\begin{matrix} 2 & 1 \\ 1 & 2 \end{matrix}$~~ , $\begin{matrix} 1 & 2 \\ 2 & 1 \end{matrix}$ by U str

$(0, 1), (1, 0)$, imp pairs

$(\infty, 0), (0, \infty)$ imp. str. pairs.

no digms are cons

$$= \begin{matrix} 1 \\ \square \end{matrix} \text{ is cons.}$$



so $\begin{matrix} 1 \\ 1 \end{matrix}$ gets by U.

$$\begin{matrix} 1 \\ \square \end{matrix} \text{ in cons.}$$

$$\begin{matrix} 1 \\ \square \end{matrix} \text{ in cons.}$$

$$= \begin{matrix} \square \\ \square \end{matrix} \text{ in cons.}$$

so

~~$\begin{matrix} 1 \\ \square \end{matrix}$~~ , $\begin{matrix} 1 \\ 1 \end{matrix}$, become get by U

~~$\begin{matrix} 1 \\ \square \end{matrix}$~~ $(=, 1), (1, =)$ is by U.

12 $\sim \begin{matrix} 1 & 0 & 1 \\ 0 & 1 & 0 \end{matrix}$

try $\boxed{1|2}$, $\boxed{2|1}$

times $(n, 0)$ \uparrow
 \therefore by U.

13 $= \begin{matrix} 1 & 0 & 0 & 1 \\ 1 & 0 & 0 & 1 \end{matrix}$

0 tried since $00 = \boxed{11} \times 0$.

$= \begin{matrix} 0 \\ 0 \end{matrix} = \boxed{1|2} \times (=, 0)$

14 $\sim \begin{matrix} 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 1 \end{matrix}$

$= \boxed{1|2} \times (=, 1 \times 0)$

$\sim \begin{matrix} 1 \\ 0 \end{matrix}$ is consistent. $(n, 1), (1, n)$ are by U.

15 $\sim \begin{matrix} 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 0 \end{matrix}$

$\boxed{1|2}$, $\boxed{2|1}$

Also $\boxed{1|2} \times \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{matrix} 0 & 1 \\ 1 & 0 \end{matrix} \rightarrow \begin{matrix} 0 & 1 \\ 1 & 0 \end{matrix}$ appears and is consistent.

to 86.25 (directly)

Th. Q. is "is $\begin{matrix} 1 & 1 \\ 0 & 0 \end{matrix}$ closer to $\begin{matrix} 1 & 2 \\ 2 & 1 \end{matrix}$ than $\begin{matrix} 0 & 1 \\ 1 & 0 \end{matrix}$ is to $0, 1, \begin{matrix} 1 \\ \text{str.} \end{matrix}$ etc.?"

Th. dist. betw. Th. strs. is 2, but Th. 2 priors of $\begin{matrix} 0 & 1 \\ 1 & 0 \end{matrix}$ might be rather low, since it may have to be built from scratch.

Do not so; we have, at a certain level, Th. dups $\begin{matrix} 0 & 1 \\ 1 & 0 \end{matrix}$ - so we can get pairs by the cart. prod. of this set with itself - then mult by $\begin{matrix} 1 & 2 \\ 2 & 1 \end{matrix}$ or $\begin{matrix} 2 & 1 \\ 1 & 2 \end{matrix}$ to get $\begin{matrix} 0 & 1 \\ 1 & 0 \end{matrix}$, etc.

So at this stage of Th. game, we cannot afford Th. jump in fact, whether we like it or not,

Rh. tetragram ~~1010~~ $\begin{matrix} 01 \\ 10 \end{matrix}$ is of such hy $\text{epri } U$,
 that it would get formed along with, say $\boxed{1112}$,
 provided 0 and 1 were around. In such a case,
 the str would not \rightarrow say $\boxed{11112}$, since it would
 be un nec, since Rh. ptetragrams would all be available.

However Rh. str. $\boxed{11112}$ would have very hy U ,
 since it has ^{had} many cases in Rh. post - from both
 $=$ and \neq . Rh. tetragram $\begin{matrix} 01 \\ 10 \end{matrix}$ has had only cases
 from \neq , and these aren't too frequent, because
 of its individual str. - i.e. there are 4

\neq trigrams: $\begin{matrix} 11 \\ 00 \end{matrix}, \begin{matrix} 00 \\ 11 \end{matrix}, \begin{matrix} 10 \\ 01 \end{matrix}, \begin{matrix} 01 \\ 10 \end{matrix}$.

We get about 2 ~~cases~~ cases of these per placement,
 so that this is only $\sim \frac{1}{2}$ as frequent as some of
 Rh. str. like $\boxed{1112}$. \leftarrow The Rh. ^{particular} str. is
not used every time - so Rh. tetgrams may
 have hyer U after all!

Well - if they want to have hyer U , there isn't much
 one can do about it!

Nor that is nec to do about it! with \oplus , all
 old ptetragms will become curious. Rh. trigrams $\begin{matrix} 10 \\ 01 \end{matrix}, \begin{matrix} 00 \\ 11 \end{matrix}, \begin{matrix} 11 \\ 00 \end{matrix}, \begin{matrix} 01 \\ 10 \end{matrix}$
 will still be hy U , as will $\boxed{11112}$, \Rightarrow $\boxed{1112}$, etc.
 These $\text{d}gms$ will suggest

etc.

~~When~~ When we get trigrams like $\begin{matrix} 11 \\ 01 \end{matrix}$ and $\begin{matrix} 1 \\ 0 \\ 1 \end{matrix}$
 it will be quite reasonable to try $\begin{matrix} 11 \\ 01 \\ 11 \end{matrix}$

~~This will be~~ For ~~cases~~ cases like \otimes and even \oplus , Rh's
 will continue to be true. T.M. will always start using
 $\boxed{11112}$, etc, which will have enormous U . ^{in retrospect,} but
 eventually Rh. hexagms or tetragms will form, and
 then hexagms or octogms will form from them by simple

Juxtaposition.

A way to avoid this and make the str. method rec. is to occasionally give examples like:

Musical notation showing a treble clef, a key signature of one sharp (F#), and a melody. To the right, a binary sequence is written: 001111 and 01010 . A box is drawn around the second '0' in the second row.

Anyway, we get back to the old saw, that if a technique is rec., it will be found — if not, it is less likely.

So: continue with the examples: Go to ⊕, ⊗ and +, trying to throw in occasionally, and noting (to reader) that Hexagony will be used, etc.

After the examples, note to reader that the generalize structural extrapolation from $\langle 112 \rangle$ to $\langle 1 \dots n \dots 12 \rangle$ will not take place until the concept of number occurs, in the sense that is useful or rec., elsewhere as well; that "the concept of no." is being "worked upon".

from 84).23

Musical notation with a circled '16' to the left. The notation is heavily scribbled over.

~ 00111
11000

$\begin{matrix} 1 & 1 \\ \square & 0 \end{matrix}$ is consistent and is used, prediction is 0.

$\square N$ I have been assuming that I.M. stores in its program memory only programs that have at least 1 case — whether or not they both consistent and in consistent programs. { This fact ^{has been} made part of the rules — see $\alpha 122.10$

Musical notation with a circled '17' to the left. To the right, a binary sequence is written: 111001 and 001010 . A box is drawn around the second '0' in the second row. Further right, a sequence of boxes is shown: $\langle 1 \dots 2 \rangle$ is used (? ?) why?

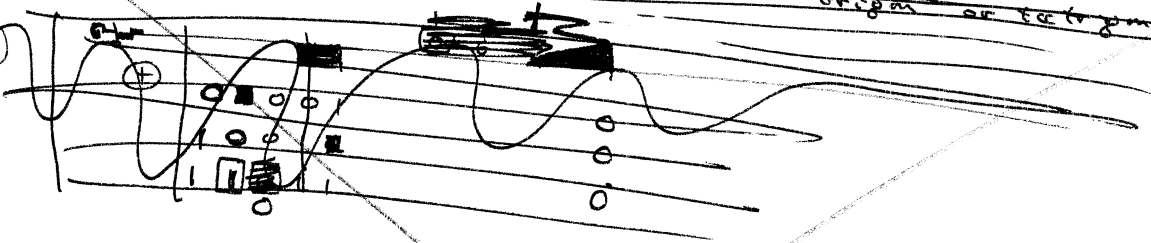
18

$$\oplus \begin{matrix} 110 \\ 010 \\ 100 \end{matrix}$$

$$\begin{matrix} 1 \\ 2 \end{matrix} \times (1, 1) = \begin{matrix} 1 \\ 1 \end{matrix} \rightarrow \begin{matrix} 1 \\ 1 \end{matrix}$$

since no other trigms or tetragms fits.

19



There are no cons. trigms of R. form $\begin{matrix} 1 \\ 0 \\ 1 \end{matrix}$, because $\begin{matrix} 1 \\ 0 \\ 1 \end{matrix}$ has occurred. $\begin{matrix} 0 \\ 1 \end{matrix}$ is, however, cons., as yet, and has reasonably by apri U. - so R. prediction is 0.

19

$$\oplus \begin{matrix} 0010 \\ 0010 \\ 1100 \end{matrix}$$

$$\begin{matrix} 0 \\ 1 \end{matrix} \rightarrow \text{incons.}$$

$$\begin{matrix} 0 \\ 0 \end{matrix} \rightarrow \text{incons}$$

$$\begin{matrix} 1 \\ 1 \end{matrix} \rightarrow \text{still incons. yet.}$$

$$\text{also } \begin{matrix} 1 \\ 1 \end{matrix} \begin{matrix} 0 \\ 0 \end{matrix} \text{ is cons.}$$

See 93) 30 for 18 and 19

EN If only pgrams with at least 1 case appear in R. pgram mem, we will not have T.M. recognize that a certain new str (say) would have been good in R. past, as well as now. T.M. only uses new ~~str~~ str to make new pgrams.

Not so: In getting R. U of a str., it is mult. by every ntp. that it can, an R. U of R. resultant pgrams is taken. This U should be calculated somehow on R. basis of just how likely a str, ntp. product will yield a useful pgram. - Th. putting on of \square may come later - i.e. suppose we make a rule that only ~~trigs~~ ngrams are produced by Str x ntp., and that to get pgrams, one must add \square later.

This point should be made clear in R. exposition.

SN

perhaps we want

$$U_{str_i \times Ntp_j}^{apri} = U_{str_i} \times U_{Ntp_j}$$

This may conform to a better \circ so what one wants. For the case in which U's are added, R.M. is simple, but if a str, say, has by U, then ~~any~~ ntp. mult. by it will give a ~~page~~ of by U apri - in fact, there will be rather poor discrimination. This will have to be looked into to a greater extent.

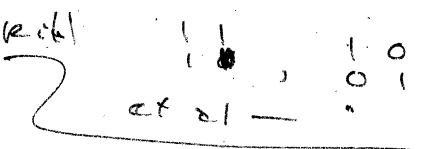
I think that now I can continue with a few examples of \oplus , then finally, +.

Note in exposition, that in (18) and (19) T.M. hasn't yet learned what \oplus is. I should make a ruff study to find out just how many examples must be given before

T.M. learns something. Apparently, the first thing that must be done is UU learning. This can only be done by

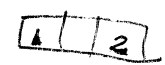
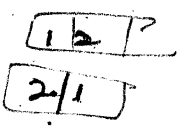
- 1) Making all R. old / ^{bad} / ~~pages~~ / incons.
- 2) Making sure that any ^{ultimately wrong} / ~~new~~ / ~~pages~~ that will be constructed, are rendered incons.

In R. case of \oplus , we want to kill



all 8 of them. We will want to keep

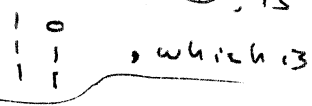
1, 0, 0, 1,



et al.

all nms like Any nms that is probably incons. - should have been in cons. long ago. 3 digits by, that doesn't contain \oplus , is

unless it is like



of low \circ prop.