

TM3

$$\frac{101.7}{3} = 33.9$$

$$\approx 33.8 \text{oz} = 1 \text{ liter!}$$

$$\frac{1000}{33.8} = 29.585798 \neq 28.35!$$

01: 451.40: Now, at a certain cond; seems to be getting a lot of wt.; ~~would~~ wouldn't we want to give it a hyper function of trials $\binom{2}{2}$ - so its influence would be more accurate (larger ssz).

02: If so, we'd have to modify trials... But drop this for linear moment!

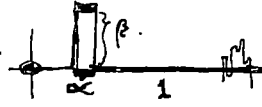
03: Suppose we don't do .00, but just run \bar{x} system for a long time. In theory, we approach \bar{x} U.D.

04: Since it ~~traces~~ all discrete ~~model~~ entries & tries all continuous values on \bar{x} etc. As we slowly \uparrow cc, TM gets better and better. In lines .03-.04 we considered $cc=00$ & it could find, but ~~for~~ because all models had $cc=00$ & no "ssz error" (ssz is no of Mt-Corpus ~~not~~ corpus size) so \bar{x} is how much error we have for ssz $cc=00$... particularly, small ssz: and could considerations of .00-.02 help it?

As each cond. accumulates data, it might be possible to compress certain off-line cards so \rightarrow to ~~minimize~~ ~~amount~~ maximize \bar{x} amt. of value. reduction in \bar{x} results per unit trial - This is somehow kind of work on SM in deciding how much money to invest in each stock so as to minimize $\frac{\sigma}{\mu}$ (or $\frac{\sigma^2}{\mu}$?).

05: So \bar{x} is; how well does it do in $cc=cjs$? (P.11)

Also Note: for a corpus of much size (say SM) pc of corpus will be very small, so cjs would be enormous: I must find way to deal w. this!

07: (1.3) Try simple $y(x)$ function: 

If pc of "soln" is, say $\alpha \cdot \beta \cdot P_1$. If it takes time T for \bar{x} trial $cjs = \frac{T}{\alpha \cdot \beta \cdot P_1}$.

08: Now let's do Mt-Corpus search. Say it takes time Δ to search find pc of a trial (to "test" it). If we spend time T in search we will spend time $\beta \cdot T$ on trials for cond; & we will do $\frac{\beta \cdot T}{\Delta}$ trials. $\alpha \cdot \frac{\beta \cdot T}{\Delta}$ is no. of trials that will be "successful" & will obtain $pc = \beta$.

Clearly, this is indep of β & \therefore seems wrong. T. Mc search seems divided into 2 parts:

24: 1) Choosing discrete part α \Rightarrow knocking out continuous parts: Actually, after "test" α region is found, we can use standard continuous optimization methods. To make α larger, use a small ssz (of corpus). Then when α is found, ~~increase~~ ~~ssz~~ ~~simple~~ find it (peak), then \uparrow ssz & look for new peak, use old peak as starting points \rightarrow for a very nice trim on P15 \rightarrow 453.20

So: It looks like main problem is finding the "x" region, & \therefore ~~is~~ cc of this would seem to not be much related to the size of $\alpha \cdot \beta$ - \therefore not $\propto \frac{1}{\alpha \cdot \beta}$

09: On the other hand: Once one has picked a cond., a single random trial in \bar{x} space, will give us an \bar{y} which is the pc of that cond. - T. problem is - P15 \bar{x} pc is (for larger corpus \rightarrow larger ssz) very small. T. prob of "hitting" α is also very small (usually) \rightarrow $\frac{1}{\alpha \cdot \beta}$ is constant not $\frac{1}{\alpha \cdot \beta}$ squares α . So while \bar{x} Expected value is fine from Bayes; T. Variance and estimate of \bar{y} & hence in one's production, is enormous! So this is why we do .24-.29!

10: If we don't do \bar{x} (.24-.29) then we have to do a num Mt-Corpus trials to get the variance of \bar{x} random pop. to \approx the var. of \bar{x} actual distribution in \bar{x} space, - which is usually quite expensive. - 453.00

On R functions & Genzup Room. There is one way on 430.15 - .33; 431.20, 431.30

Note: 430.10:
I did work out
a nice MB ratio
using 430.15 -
Market data
~ 400.00!
430.10 is mainly about
2 priors based on self!

1 To do a search over a space of functs w. discrete & continuous params.
This is to get functs of by p.c. as "it" well - for induction (e.g. SVM).

2 Say one does it MonteCarlish! We ^{randomly} choose $lead_i \in \mathcal{W}$; W p.c. = P_i ; P_i is p.c. of
discrete part of $Cond_i$'s desc. Say its a D param d.f. (T. discuss. probly doesn't care how many
contin params (Ruvro are). T. param uses variable exp from 0 to 1, (0, 1).
It takes $c = 2^{10}$ to generate a trials, to see how good it fits is ($c = \dots \approx P_i \approx 1/P$)

5 Say $Y(x) \in \mathcal{R}^n = Y(x)$ (x is contin param on $(0, 1)$). $Y(x)$
We are usually interested in $\exists y$: i.e., $\int_0^1 Y(x) dx / \int_0^1 dx$
say $\exists y_i$ is expected value of y for $cond_i$.

10 To make a prediction, ^{say} R runs over just r different discrete $cond_i$'s.
 Z_i = prediction using present data, using ~~discrete~~ discrete params of $cond_i$ & continuous param $\exists y_i$.
Then we have r different preds. Per predn $\exists y_i$ is given w/ P_i . These preds are t. d.f.
of r. output of our predn system. (See (33) for error.)

15 Say we have many ~~cond~~ $cond$'s ($r \gg 1$) & we only have v trials per $cond$.
It would seem like a good idea to do more trials ~~for~~ for $cond$'s in which t. single y_i was large.
In the usual MC case Lsrch., we spend a fixed amount of time Δ on each Mt. C. trial.
So then it's we make do $\frac{\Delta}{b_i}$ trials for each time $cond_i$ is chosen (I say $\frac{\Delta}{b_i}$ because b_i
doesn't consider "startup costs" for a set of trials - in general, f. no. of trials for $cond_i$ will be
slightly
a more complicated func than $\frac{\Delta}{b_i}$)

20 If w_i is small and δ we have a no. of trials $\gg v$, then we will have several
 Δ periods per $cond$ - so t. no. of trials for $cond_i$ is $\propto P_i$, v. no. of times P_i is (randomly) selected.

25 Th. P. on .25-.23 on predn is wrong. Say $b_i = b$ is t. same for all i .
(Samp inc. time for each trial). When we choose x_i for $cond_i$ & its output is Y_i , we then make
a prediction, using ^{b_i} discrete params of $cond_i$ & continuous param, x_i & we get the
predn w/ Y_i . We do many ~~trials~~ trials, so t. no. of trials $cond$ does is $\propto P_i$, &
t. no. of micro trials it does is $\propto \frac{P_i}{b}$. So we just w/ all predns using x_i by v correct predns N_i .
The w/ \propto to P_i comes automatically since we f. no. of trials for $cond_i$ is $\propto P_i$.
Spec
452.00

This is wrong!
See (33)

Q: 447.40: SM cont.: But, say I don't bet every day but only on certain days, Even so I have ^{empirical} sequences of fractions of money made (i.e. $\frac{\text{sum of } P_i}{\text{previous bank}}$) & some hairy, I convert discrete probs — so I have a PC seq. & the PC of coding to a strategy, plus need an aprp — which enables us to compute "true expected yield" of the strategy.

Say I have r different strategies: Strategy has aprp \sqrt{Pr} (due to code cost) & multi yield Y_r for t time intervals (of unit).

The expected yield from this set of strat is $\sum_r P_r Y_r$.

If the set of strat has only 1 param, α w. uniform aprp from 0 to 1 then

expected yield is $\int_0^1 dx K(x)$.

Review past stuff in Feb extending of good content forms for ALP

A perhaps V.G. way to do this that would also write tests for TM:

List a large bunch of strategies: Then "factor" them & write a grammar for them, then use it. Grammar generates new strategies. Formally strategy in the original set, given α w/ α its score

Some strategies w/ param schemes or parts of them:

- 1) linear predn. 2) N.I. predn some standard ways: ANN, RANN.
- 2) If system is in state S in morning, then bet using system S in afternoon.

For example: If system found k points in Q minutes on opening; Then bet for fixed time T_0 or until end of day.

3) Cross Predn schemes:

4) Ways to save time in searching.

5) Perhaps look at old copies of "Stocks & Commodities" — or get their CD Rom.

6) Try listing kinds of Prediction (or Compression) methods — from e.g. "Pattern Class'n" & other ~~revisions~~

D25.03

TM3

SM 21

WWW.02515casino.com
1.800.922.4515

Give me odds for the Betting
in CARACAS (Caribbean Island, Dutch)
But ~~25 to 200~~ 25 to 200 per bet. - RR
\$ 200 is 200 so Big! (Make out for me!)
\$1 to 5000 for Sports betting (Non-RR)

447

10:446.40

Another similar problem: Should I spend ~~on~~ trying to cure Cancer or on Trading
General Relativity?

12: On defining n against BZZ : in 446.00 it is considered ^{names} ~~surprises~~ in the corpus that ends in β . The set of all n against corpus that ends in β is surprise. - i.e. β . By looking at β it may be able to derive definitions "But are ① cheap ② cover much of β ③ do not include ~~many~~ many n 's that ~~end~~ don't end in β ."

Another view: In BZZ map are treating the corpus as an "encyclopedia" and the problem is to obtain a cheap code (w. that corpus as input) for the next symbol. Cheap in cost (not hardly cost). Stated this way, ^{it seems} ~~identical~~ (to finding the pc of the next symbol of the corpus! (??))

11
0
BZZ Play: Discern. w. Steve Wilton: To code ~~an~~ a / length n or n corpus, use ~~all~~ all n shifts of the corpus in lexical order (probably I'd want to do corpus backwards, so I could predict forward!)
To do this incrementally, we just add a symbol to the corpus, then insert the augmented corpus in the list of shifted corpi.
To get pc's it will be best to put each shifted copy in lexical order w.r.t. to ~~the~~ all but last symbol (8₂ last symbol being the one to be predicted). Or, if one does a credit shift & gets them in backwards lex order, this would also be o.k.

Another way to look at ~~BZZ~~ codes internally: say the terminal symbol "xy" occurs, and we have a k symbol alphabet (Rustar), & the corpus is of length n . Including "xyz". The new z is a n -gram of length n , with 2^{n-1} members - all ending in z . "yz" is similarly a marker of the $n-2$ member n -gram of all n symbol strings that end in "yz".

24
It's more of a kind of sequence. Use BZZ to see if I can do useful predn.
Use long sequence, so that startup costs are minimal. Also look at actual profits obtained, so I'm sure no bits are being wasted.

26
Try it with variance smoothing: Obtain σ^2 in one of several ways & a quantization ^{into k levels} in view of expected variance. Some ways to get σ^2 ① from Option prices ② x window ③ Median filter: so σ^2 is used as d.f. - RR is shout work better than Gaussian.
Q is there a way to do x window median filtering?

3
31
Try .26 w. various window widths & various stocks. Contrast w. 1, 2, 3, 4 corr (know prob).
Using the quantization method of .26-.31, do prediction based on ^{previous} 1 or 2 or 3 lengths n grams (simple Markov Model). Compare resultant Entropy w. Price obtained by BZZ.

Try .21 ff in cross predict using BZZ: I may be able to modify lexical analysis of 10 ff
SN on SM ~~See how better in a strategy using Kelly betting (heavy)~~ See I have a strategy n is σ^2 of form that says bet of any time or Bank. Then I can discuss Kelly corollary assoc. w. that strategy, & probability estimates for each bet. This enables me to use a coding cost for the strategy & use AIP to decide how much w.t. to give a bet. After the based \rightarrow 449.00
rounding & while. Spec

~~HM3~~

On limits of regularity in BZZ: Consider a linear corpus that is a linear seq. of symbols (not necessarily)

The best token is β . Consider a regularity that is a set of all substrings (to data) that have terminated in β .

Call this regularity $\tilde{\beta}$. (Since all the regms in $\tilde{\beta}$ are of length ≥ 1 and none of them have ever occurred before, so they can't be "productive")
No! regms that are outside/within of β that $\tilde{\beta}$ have short chains $\tilde{\beta}$ are productive (in the sense of having occurred at least once before). Some ways to construct such regularities:

1) T. elements of $\tilde{\beta}$ are all of length $\leq k$. Augment β with sequences that are "doubtless" from t to j in element, but otherwise are identical to some seq. in β .

2) 02-03 is wrong: regms in $\tilde{\beta}$ are of various lengths, but no 2 elements of β are the same length. Defn. $\tilde{\beta}$ elements "backward" - so that first symbol is always β . $\tilde{\beta}$ of length r . ($r \in \mathbb{N}$)

Using this "backward" form: A way to construct a regularity (say $\tilde{\beta}$) with another regularity (say $\tilde{\alpha}$) is to take $\tilde{\beta}$ and $\tilde{\alpha}$ and combine them by giving intervals $i, j \in \mathbb{Z}$ where $i < j$ and $\tilde{\alpha}$ is in (i, j) we AND $\tilde{\beta}$ is U or for other regions we OR $\tilde{\beta}$ is V. Say $i = -3, j = 7$. Then $\tilde{\beta}$ is U AND $\tilde{\alpha}$ is V from 1 to 2 and 8 to ∞ ; and OR from 3 to 7, we get $\tilde{\beta}$ set but from 3 to 7 inclusive.

any symbols are legal parts of $\tilde{\beta}$ regularity.

\dagger found is an early sketch with problem: In $\tilde{\alpha}$ so (so different mechanism) I had regms $\tilde{\alpha}$ regms, but I didn't appreciate that regms were "inverses" rather than "forms", i.e. that "wild regms" were a complete unobscured induction. What I probably should do: In the induction problems in my TSO: write up the "solutions" to problems that I want to use; then see how I can derive ways to construct (combine) regularities to get the needed solutions.

In 00-14 I think I was trying to expand the scope of applicability of BZZ, both here, and so,

I should perhaps first decide which new kinds of regms I want to be able to discover - then build these regms, regms, regularities etc to express a given class of new kind of regms.

SM the BZZ model is an open Markov model. Each symbol puts us in a known state. We can then derive sequences that put us in a known state, or "assign a different token" to symbols that put us in a certain state.

Hidden Markov Models (HMM) must be more general than Open Markov Models since OMM is a special case of HMM. This is peculiar! HMM would seem to have essentially more powerful set of models than OMM.

[Model OMM is like HMM in that "state" is never directly observable, it's inferred probabilistically; it's never known for sure]

Yet OMM can be generalized to wild regms, which is unreal! This causal doesn't seem to "pass thru" HMM, yet it must include HMM!

In explaining how "Induction etc, deals a. learning; & Sci method, "Sci method is how the Sci Community

"Learns" (to some extent). Hvt. T. Sci community is more like org. evolution in that it really has no "top goals" - it makes up sub-goals as it develops. These sub-goals can lead to where or to what very peculiar places. I suspect that the goal of "understanding the Universe as a whole" may be meaningless.

If the Sci Community is given a problem, I think I know how it solves it or should solve it.

T. Q of what it decides to solve is usually inductive, but very randomish. Individual have means that

its new sub-goals are patterned after what seemed useful in the past? Deciding on what new data to (acquire/pursue) is decided in a similar way.

Annals
Reviews
Sci Community
Exp evolv
All 3 learn
in their own
way. Tury to
Sci Community
(rns: e. 2. 11. 6)
"Sci Method"

447.00

REV:

The Bumpin P-500 D23.027PM
No. of bumps & w. numer. 24 bumps Max } to cost P-433
00 - .40 } 2 bumps
I.e. a not bad quick revision of $\sim 400.00 - 440.40$

Rows: 437.00ff - 438, 439, 440, 441
429.31 - .70 brief
421.12 - 18, 15 - 40
422.00 - .40
403.24

10:44:1.40: One way to organize f. ideas of 441.00 - .40: Consider what f. tool is the most direct approach

to strong AI: Tell how to do it: then show how (most of) f. ideas of 441.00 - .40 fit in.

The macro is: Phase 1, then Phase 2: (Tell in detail how Phase 1, 2 work). is incomp. table

In "Phase 1" we have an a priori multivariate model of f. data PD1: PD1 is a univ. D. f. incomp. table

To find good models for PD1, we give PD2 which is a PD to guide L search for good approx. Best Nat. dist way to search, search, BUT OK for Phase 1

In of we may run into "Encyc" problem! Given "free" by date of date, E (like Encyc. or Internet) to get best ans.

to Question w. CC S C B. If one has many Q's w. answers in E, one may describe "Index" "to E".
Expln: Encyc problem occurs when we don't want to have to code entire corpus to answer a Q. Perhaps "just record E" as part of Q

The TSO for PD2 is f. seq. of Q's, A's & hyper functions O's, but has been successful in types of f.

Phase 1's strategy is to try techniques u to Phase 2 were successful in f. past. Phase 2 tries for "Best soln available to CB"

The TSO for PD1 will be (mg & dates of Maple, Mathematica: eg. eval, eval literally, eval uncorrectly).
Solve ^{inverses forms} ~~diff~~ derivatives, integrals. These are all d-proba Q's to start.

T. problem of solving to S-proba is simple. - change of TSO: One of big problems: What TSO to use?
How to represent S - funts so that previous d-funct solns & their components are useful for the TSO?
How to use d-funct solns. to help create S-funct solns.

In both PD1 & PD2, context is imp part of discovery of chart codes ("Search").
w.o. context, each token has a unconditional probability. By using definitions of nouns, pc's of tokens become very small as size of (search, problem) if context is ignored.

implement use of nouns for pred context \rightarrow predn. Hvr, defining only nouns will only give noun context.
for more general context, we need to be able to define verbs: for most general kind of context, we

need to define wfd nouns. Hvr, definitions has advantage of speed, but disadvantage of possibly wrong "parse".

EN Re: 05 ff: If we don't want to define ^{wfd nouns} ~~verbs~~ as contexts, then it would seem that dates of these sorts Hvr, if nouns defined as language, "context" (informal nprn) \rightarrow automatically considered

can define nouns then we don't need verbs to define noun context. Similar exists hold for defining nouns & wfd nouns. (See 441.23 for way to define nouns & wfd nouns (Ray last is any pd on string: d-construct (by) a stochastic, or by defining a wfd machine or function.

One apparently good idea: that Definitions were not necessary for induction but they did a certain amt. of "pre parsing" of f. corpus to reduce same in later parsing. The disadvantage is that f. dates did not

always parse f. corpus properly \rightarrow so they should be used only if they have by PC: or if corpus is small
Set of ^{dates} ~~dates~~ that have almost all of f. pc or decision" at a branch.

E.g. B22 doesn't use dates - but it is very fast. This is perhaps because: only kinds of Recognition via B22 + B3
Coastal uses is nouns (no nouns or wfd nouns)

One of Big Ideas of last month: Use of B22 for PD1 & PD2: Hvr, note .28!

Also note need to deal w. "degeneracy" in codes for functions & "symmetric" & other equiv.

functional forms: This is not a problem peculiar to B22, hvr, it is important in all of TM's

all of probty evaln, equiv evaln, search.

Items in 441.00 - 12 not covered (mentioned) in 442.00 - 33!

Alternative Approach to "STRONG AI": 415.20 - 33 ~~lists~~ lists ~ 14 induction methods/AI methods.

Each has deficiencies wrt implementing "STRONG AI". How to Mix them to get ^{Strong} SAI?

One way: For each approach, list good things it does: its useful skills.
Also tell how it misses SAI: what it needs to get SAI.

One way is to see how each of strengths of the others can be brought out: deficiencies of each method.

20: 432 $\frac{3}{4}$!

It may be poss. to prove 432 $\frac{3}{4}$ by HM induction!

Write out expression as $\prod_{i=0}^n R(i)R(n-i)$; $R(i) = \prod_{j=0}^i \frac{4j+2}{4j+2}$

perhaps write expressions in terms of factorials.

$$\frac{2n+1}{2n} = \frac{4n+2}{n+2} = \frac{2(2n+1)}{n+2}$$

$$= \frac{2}{2} \cdot \frac{2}{3} \cdot \frac{3}{4} \cdot \frac{2}{5} \cdot \frac{3}{6} \cdot \frac{2}{7} \cdot \frac{3}{8} \cdot \frac{2}{9} \cdot \frac{3}{10} \cdot \frac{2}{11}$$

$$R(n) = \left(\frac{2 \cdot 7}{5} \right) \cdot \left(\frac{3}{4} \right) + 1 + \frac{n+1}{4n+2} + \frac{n}{4n-4} \cdot 2$$

$$R(n)R(0) + \frac{R(n) \cdot R(n-1)}{R(n)} \frac{R(0)R(1)}{R(0)} + \dots = R(n)R(0) \left(\frac{R(n-1) \cdot R(1)}{R(n)R(0)} + \frac{R(n-2) \cdot R(2)}{R(n-1)R(1)} + \dots \right)$$

$$= R(n)R(0) \cdot \left(\frac{n+1}{4n-2} \cdot \frac{2}{2} + \frac{n}{4n-6} \cdot \frac{3}{3} + \frac{n-1}{4n-10} \cdot \frac{4}{4} \right)$$

show this = $\frac{4n+2}{4n+2}$

$$\frac{n+i}{4n-j} \cdot \frac{j}{3-i} = \frac{n+i}{4n-4i+6} \cdot \frac{4i-6}{3-i} \quad j = 4i-6$$

$$= \frac{n-i+2}{4n-4i+2} \cdot \frac{4i-2}{7-i} = n$$

$$= \frac{n-i-1+3}{4i+2} \cdot \left(\frac{n+3}{4i+2} - 1 \right) = \frac{1}{4i-2}$$

26

Note: Something like $n!R(n)$ is actually redundant since we can put things in any order as well as arrange parents various ways:

	1	2	3	4	5	6
n	0	1	2	3	4	5
$R(n)$	1	1	2	5	14	42
$n!$	1	1	2	6	24	120
$F(n) = n!R(n)$	1	1	4	30	336	5040
$R(n)/R(n-1)$	1	1	1/2	1/3	1/4	1/5
$G(n) = (n!) \cdot R(n)$	1	2	12	120	1680	
$G(n)/G(n-1) = 2(2n-1)$	1	2	6	10	14	

$n!R(n)$ is actually redundant since we can put things in any order as well as arrange parents various ways:
 I want to solve!
 See 458.25 R
 47M 6.26 R
 EN $\frac{1}{x-1} = \sum_{n=0}^{\infty} \frac{n!}{x(x+1)\dots(x+n)}$
 Pol: "Operational Calculus" p 252 (example 1)
 Pol using Gen Fun: p 251, 253, 272
 Taylor exp: $\frac{d}{dx} \left(\frac{1}{x+1} \right) = -\frac{1}{(x+1)^2}$
 $\frac{d^2}{dx^2} \left(\frac{1}{x+1} \right) = \frac{2}{(x+1)^3}$

$G(n)/G(n-1) = 2(2n-1)$ No great surprise, since we have simpler form for $R(n)$
 actually $(n-1)!R(n)$ is a better redundancy

$$(n-1)!R(n) = \frac{(2n)!}{n!} = (n+1) \dots (2n)$$

n	1	2	3
$(n-1)!R(n)$	2	12	1680

47M 6.26
 458.20

$$80 \times 130 \quad 9 \times 9 \times 100 \times 2 \quad 16200 = \frac{16.2k}{x} \text{ cont. pt.}$$

$$\approx \frac{360 \times 80 \times 2}{x}$$

TM3

≡ 3TM

spec

o: 445.40

What I wanted the probability of the next symbol being β . If I were using the context

γ , (the ngram preceding β) to obtain probability, I could obtain the same result by defining the ngram $\gamma \cdot \beta$, and not use context [- I'm not completely sure of this effect, however, I have not worked out the detailed math to see just how (or if) it works.] end of F.N.

There are at least two kinds of Context. The first "External Context", tells where the problem came from a/o how it arose. Some examples: chemistry; Physics, Maths, differential Equations. Or the problem arose as a subproblem in converting a INV problem to a OZ problem, or the problem arose after I tried to integrate using "Integration by parts".

The second "Internal Context" relates the previous corpus (or part of it) to the symbol to follow. Most generally, it is the problem of sequential prediction. Less generally, the ^{last} ~~previous~~ n known symbols is a context for predicting the symbol to follow. More generally the previous (ngram) may be a member of an ngram set (ngms) that can be used for predicting the ^{the} probability of the next symbol. The most general kind of context is the ^{weighted} ~~weighted~~ ngram; this can correspond to using the Universal Distribution for prediction, if we use the entire previous corpus as context.

A common kind of context ngram is used in raster scan graphics. If a line is k pixels long, then any pixel will have as context: ① pixels immediately before it ② ^{a pixel} that is k pixels before it. ~~ANALOGOUS~~ This latter corresponds to the set of all ngrams of length k+1, that start with a particular pixel. If each pixel can have r values, there will be r^k ngrams in that ngram set.

Ngrams correspond to deterministic "formal languages". They can be described by generative grammars or by recognition functions.

Weighted ngrams correspond to probabilistic languages. ~~They are~~ They are probability density functions on the set of all strings. They can be described by stochastic generative grammars or by functions that map strings into probabilities.

To obtain the probability that the next symbol in a corpus will be X_i , we have an associated ngram, N_i . For d-induction N_i will be a d-ngram.

It could be defined by a function $F_i(x)$ say the range of F_i - the set of outputs, for all legal inputs. Associated with F_i will be its a priori probability, α_i - the prior description.

For S-induction on that same corpus, N_i will be a S-ngram - a ~~single~~ weighted set of ngrams - with each weight dependent on the description (length of that ngram). If the N_i ngram for d-induction was obtained as the range of a function F_i (the set of outputs of F_i with legal inputs) then we can associate with each (legal output of F_i , the total pc of all inputs to F_i that obtain that output - ~~creating~~ ^{creating} a p-ngram.

→ 4T4 12.00
cont for 11 150.00
↓
450.00

TM3 Rev

10:444.40

There are many ways to try to solve the optimization problem of 444.24.

In "Phase 1" we will ^{begin by} using L search, with a guiding PD that is to some extent PD₁.

After we have solved several problems, we will use a guiding PD that is based on the sequence of ~~successful~~ O's ~~that have been~~ successful solutions.

O_k will be a function that successfully optimized 444.24 for Q₁, ..., Q_k.

One way to do this is to use BZZ (or an augmentation of it) to extrapolate the sequence O₁, O₂, ..., O_k. BZZ will give probabilities to all possible continuations of this sequence.

Then this gives PD₂, the proby distribution on O_{k+1} — which are the candidates we test, using 444.24. We try the cand. in ~~the~~ PD₂ order and spend a maximum CC on each test is proportional to its ~~PD~~ PD₂.

[N.B. I do mean PD₂ :]

It is the PD that assigns CC's ^{limits to candidates} that guides to such, not to order in which trials are made } Arbitrariness!
↑
no enough

The initial TSQ (≡ Training Sequence) will be the learning of definitions from Maple, Mathematica, etc. Learning: evals, eval. numerically, differentiate, Integrate, solve literally, solve numerically... can be an adequately difficult task — ~~some of~~ some of these functions can be arbitrarily difficult.

The initial training sequence will involve d-induction (determined induction), in which there is only one correct answer for each question. Mathematical and (general) problems can be of this kind. Though d-induction certainly does not cover the range of problems that we want TM to be able to solve, nevertheless, if we are able to teach TM to solve really difficult problems of this sort, it will be an encouragement for the general methodology we have been using.

Next on agenda is p-induction (probabilistic induction), in which each question may have several different answers — each with an associated probability.

We want to devise TSQ's for p-induction that are able to use the concepts/definitions used in d-induction as useful concepts and sub-concepts. Some suggestions on how this might be done are in the section on Context.

p-induction ≡ s-induction

Actually, this is a BIG topic

Context: The need for contextual considerations arose in the "Alg. Notation Lang. Program" in Sep 88. The brood of solutions of successive problems increased very rapidly — so it was clear that ~~the~~ difficult problems would become impractical to solve. By taking "local context" (i.e. recent previous nems) into account, it was felt that the pc of ~~the~~ the next trial bit would be large enough to make the solutions of difficult problems possible.

30

[An aside (x Footnote): More recently I realized that if I was using nem definitions for

induction, that simple nem contexts would not increase my production accuracy: Suppose

SPC
448.00

TM 3 Rev:

: Summary of Notes as tops of pp.

- 1) Rev: see to part 443 for list.
- 2) Org. Evoln. 414, 426, 427, 429, 439
- 3) TSO, 419, 425,
- 4) Other 1 page listings: 413: How to Construct Recognition using BZZ

418, 420 off the Redundancy factor: Nuclear Res topic: Internal cross refs.

425 Contacts (Nuclear Res ~~topic~~ after 425)

439 SM: Importance of Gamble in strategy management.

405 OSL: This may not be such a new idea: BZZ does it & Predictive PD2 does it

Top indices:

- 405 OSL
- 403 TSO writing
- 403 Rev
- 399 File of 589
- 439 SM
- 399 OSL Rev
- 429 Org Evol. (outlets)
- 429 Rev
- 427 OSL Rev
- 426 Org Evol
- 425 TSO
- 425 Contacts (input refm.)
- 421 Rev
- 420 "Not a problem"
- 419 TSO
- 419 "redundancy factor"
- 414 OSL Rev
- 413 how low to R recognition objects.

O.k. so: a more detailed view of 443.04 - 40: I will then add relevant sections of ~ 400.00 off > 0 refs to new ideas.

First: The problem is, given a sequence of Q/A's: $Q_1, A_1, \dots, Q_n, A_n$: To find a set of such that

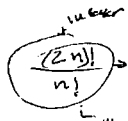
$$\alpha \equiv \sum_{j=1}^n z_j \prod_{i=1}^j O^i(A_j | Q_i) \text{ is as large as possible } \dots \text{ in the available time (or } \epsilon_B = \text{ computation cost limit) } \equiv \text{ "Computation Bound"}$$

z_j is the a priori probability of the unknown probabilistic function, $O^i(\cdot | \cdot)$.

To solve this problem, we start with a "known" PD₁: which is the a priori z_j on the set of probabilistic functions, $O^i(\cdot | \cdot)$. PD₁ can be realized in several ways:

- ① The AZ language of the IDS (A report)
- ② The Fort-line OOPS lang. that Jeurgan uses
- ③ Either or these augmented by \approx BZZ, which exploits redundancy by observing repetitions of ngrams.
- ④ Probably many other ways.

The appropriate PD₁ should be selected by the "Tramer" so it seems useful for the Tramer's sequence being used.



$\frac{(2n)!}{2^n n!} = 1 \cdot 3 \cdot 5 \dots (2n-1)$

$n!$	1	2	6	24	120	720
n	0	1	2	3	4	5
$R(n)$	1	1	2	5	14	42
$n! R(n)$	2	12	120	1680		
	6	10	14			

$3 \ 5 \ 7 \leftarrow NR!! \leftarrow \text{This may show } R(n) \text{ is integer!}$

$\frac{(2n)!}{n!} = n(n-1)(n-2)\dots(2n)$

- 20: So: ideas: 437:00ff!
- ① 2 PDs (PD1; PD2)
 - ② FN2 of SB9 = Related to (PD2, PD1): Phase2 of TM; Adaptive Esch. [This is an impt. set of ideas!]
 - ③ Context: kinds of; Generals of: related (2 to) General induction: Use to do (w/ "Scaling" problems)

- ④ ISQ's for Phase1: from MPTs defns, & MTM etc.
- ⑤ Use of \approx BZZ for Phase1; for PD2 (or PD1). enhancement of it.
- ⑥ Impt. of "Symmetries" in ~~various~~ aplys of functs that are solvs to prob. (Multi use of BZZ)
- ⑦ Orig. Evolv: Various ideas! Use of BZZ for GA!
- ⑧ MTM solms \rightarrow NMTM solms: related to "How to Express S-Points"; Also note General Context problem (0.03); The 3 vector soln to ~~the~~ n(1) problem in Phase2. (Scan .23) \rightarrow

⑨ STRONG AI: Try to come by modifying, mixing, & generally various recent, & "Mature" approaches to induction 415, 20-33 lists some approaches. \rightarrow (20/2)

That idea of last comp. of continues forms for predn

⑩ One Approach: List various Methods: Tell just why each is deficient as approach to strong AI. Also, tell what each is good for:

Then see which of them are ideal w/ deficiencies of any of them.

⑪ Inputs: How to ^(dec) parameter: ① list elements ② any distance function from canonical element, thresholded. (Koskine, EuclN, Manhattan, ...)

② combine inputs (Boolean; Concatenation) ④ Grammars: CFG; CSG; Boolean Trunc Grammar (General form for Decision Criterion) "Recognizers" Any via machine defines ed on will be in \approx a wild guess

Grammar or Language "is ~~the~~ General form - can be universal:

Q: Are inputs a universal soln. for d-induction? Say we have a certain UTM machine; The set of all poss outputs defines an input. This seems like a very general way to use inputs. The proc of δ set is r. proc of δ machine (or function) that defines δ set. E.g. use τ AZ formalism to assign ρ 's to functions; Then a given function defines τ range of that function. Or - conversely δ . range of δ function defines τ . input (This "range" is for all poss. inputs). A definition of this kind is used to define (generative defn) d-languages. (CFG; CSG) We consider τ (i.e. τ input) to be all poss. outputs of τ generator function

\rightarrow By "New Ideas" of this last Month - Poss. use of BZZ for Phase 1: Use for PD1 & PD2. As is: BZZ has no "Definitions"

TMS

SN: At various stages of TSO, we may want TM to have different "wts" for importance of different parts of TSO. This can be approximated by giving wts factors that TM will use as exponents of a 5500 PC of subcorp (or ln of those wts. for b costs).

SN: For recent "Review": For each item reviewed: Give final state of that problem: just what needs to be (solved/done)? Keep this list in mind! ABCde

26: 438.04: => (MTM -> NMTM tso problem) It certainly seems clear that conc used for MTM should be useful for NMTM. We can teach TM prob, statistics directly as "Academic Problems". Many (?) problems in this area can be formulated as MTM problems! - e.g. many prob (any or most) problems in books texts on prob/statistics can be reformulated so they only have one correct soln. Th. MTM (d-concs) used to solve such problems should be useful to solve ~~many~~ problems of S-prodn.

A NICE thing about this is that TM could learn many different ways to construct S-functions. Recall to convert before from

Another NICE thing about MTM (and) is that much of math can be used in this form - in particular, we can give MTM problems in logical reasoning - Recall prior problems in which "logical reasoning" is useful to help solve them - connection between "textbook logic" & "logic to be used for solving real problems" is "analogical" & learnable by TM.

Note that ~~NMTM~~ NMTM's solns need not be in \exists IU form; any form is usable.

E.g. The \exists parameter-forms of h(A) functions (See 438.07-08 for Bibli. refs)

D15-02 N.B. - While (perhaps) all of this is d-mathem (in formalism), I strongly suspect that S-induction is needed in many cases for d-functions solving. Some of this probabilistic work can be done via PDZ in "Phase" - we can augment this PDZ by (perhaps) by making improvements to BZZ - like neurons & (finally) wtd neurons.

dividing
data
www.
mcmtest.org.com

NB BZZ.zip: written in "Euphonia"; it's a set of files to implement -> (BZZ).

So see what it is: there are 2 Euphonia C: X/Tr. (~~some~~)
Euphonia is normally an interpreted lang. - But speeds up w. conversion to C.
In c:/BZZ dir I have a decompiled "BZZ.zip" file Not easy to understand!
[Think Main file is BZZ.E. - only one "E" Atom directory,
- (But it could be that Microsoft would be done by a .DLL file)

- SN: 1) try BZZ on BZZ for various tests, errors, etc. See if it can \uparrow compression
- 2) association in A, -f; X \neq it's input. Sources of ~~some~~ degeneracy
- 3) Can I use Euphonia functions in Fort, to speed up Euphonia?
- 4) BZZ doesn't actually do "production", it only knows data. In this sense, it is ideal for "operating", which does some things.

D.12.03 Row .00 - .07
T.13 SM .14 ← (T. Importance of Gambler's Ruin in strategy devine)
Org on .00

0 :438.02 : Org. Evolu: In a large \geq viable individuals, combination is almost always unique! Variation is wrt. alleles which are similar sets of "genes" : How can we characterize this aspect of org evolu? Another brick is "Dutch Book", in which in many/most cases, a potentially unviable cond. is not born in 1st place - saving work cc in trial cost.

07 ~ +17.00 -26 on/improvements in $\sum_{i=1}^n$ cpts / w/ data Eng. text. - 1st what is to "know" d.f. ? (rank v.s. Probab).

08:437.16 SM Same ideas on summ ~~with~~ functions: I can suspect 2 funcs are identical if they sum to same for several diffnt (or one big) random inputs. Can B use this (does it work)?

- SM To Agreement BZZ, consider Regs of finity res Detected (used) by ~~the~~ Wolff & by ~~the~~ Ming Li in his Gambles (gen)

Also, try to think of other prodns Regs, techniques, that ~~can~~ could be included to improve BZZ

Remendar, h/m/s is abt improvement of BZZ & w/d Neests is k. most general type of prodn. poss!?

4 SM In general, when ones "Bank", $B \leq B_0$, multi yield ≤ 1 . Say $\lambda =$ mult yield.

then $\lambda = \lambda(B)$ $\lambda(B_0) = 0$; $\lambda(B \rightarrow \infty) = \lambda_0$ (which is sum finite value) ← this is or fully conf=1.

so $\lim_{B \rightarrow \infty} \lambda = \lambda_0$. As B approaches B_0 , k (cholly are) \downarrow & become 0 at $B = B_0$.
 $k \rightarrow 1$ as $B \rightarrow \infty$. We want to under k \geq funct of B so our prob of Gambler's ruin (i.e. $B = B_0$) becomes small.

As $B \rightarrow \infty$, h/vr, prob of ~~the~~ GR (Gambler's Ruin) becomes small, i. probly we should use as bet criterion

$k \geq 1$. T. reason is: $k \geq 1$ if we ^{approach} ~~use~~ max expected gain (rather than max $\ln(E_{2M})$). ~~for~~ H/vr. for max expected gain, $k \rightarrow \infty$ (i.e. we bet as much as we can borrow).

So for proper bet (not abt GR), we may want to use $k > 1$: Noting that while $\leftarrow E(\ln(E_{2M})) \downarrow$ $E(E_{2M}) \uparrow$ (both $>$) — we may want $E(E_{2M})$ to be max, but w. some ~~the~~ limits on prob of GR.

So considerations of GR ordinarily are ~~as~~ implies considerations of gain or $\ln(E_{2M})$ or whatever funct of E_{2M} one is interested in.

The use of Gambler's Ruin is "best" because of "Dutch Book" arg. . . .

"Dutch Book" may not be relevant in real Mkt. & iterations, h/vr. — including prob of GR.

Perhaps $\max E(\ln(E_{2M}))$ is $\approx \max_{\text{bet}} E(E_{2M})$ if one considers one will be making many bets. From a "Dynamic Pymg" pt. of view, $\max \ln(E_{2M})$ ^{GR} ~~is~~ in each bet

32 may be best!

One way to bet: use k as large as can be (> 1 for large and B) i. yet prob of GR abt of prodns \geq a const. ← (But note arg. of (30 - 32)!) —

TM3

R-24

Org. evolv. 10

00: (Spec) 437.19: ~~410.20-26; 425.00-09~~ of (very) dense of concs like solve integrals, differential eqs, simplify. From Maple or μ Proced.

02 Org. Evolv. compared to (Sumac or Backhaul) 414.10 426.00; 427.00, 428.28, 439.00

02 03 04 \odot 524 up has some TM has solved $d-TSQ$: How to use the Phase traces ~~is~~ as Σ begins of Σ $S-TSQ$: 437.20 is relevant. Note: Many earlier nodes. Also see 420.02, 421.20, 429.09, 440.06, 437.20. \odot 414.10-415.40: Possi. Modifs. in Σ of Various "Micro-inductive Systems" to create "STRONG AI". 415.20-33 lists some inductive systems. ART Life is Good Cond.

07 08 \odot Gathering $h(T)$ functions for PDZ: 429.09-19; 430.15-33; (431.20ff, 23ff, 30ff) 433.00-30

• 433.31-434.28: "A kind of understanding" of context? Revisited ~~carefully~~: I don't yet seem integrated view" in it.

10 \odot Great Idea! \rightarrow see 437.27 for earlier version! Try BZZ \rightarrow degeneracy reduction on GA type problems! See of it \uparrow soln. rates.

Try it on some problems that regular GA has been used on. Try to find problems in which complicated functions lead to problems. Perhaps Koza's books would help. Also look at the Review Book on GP (By German writers) - also other GA review books I know.

It would perhaps work fine for problems in which Σ evolve. function took lots of time, so one could afford expensive ways to choose ~~(trials)~~ trials.

Another possy would be to try it on (strong) ANL. (Asy. Minimization).

07 08 20 On 2 Parameter level: for Σ SUMAC: Retain ~ 100 fairly disparate codes of corpus (so they are at distinct \rightarrow D w/ BZZ) Use BZZ to search for "near" extensions of "codes Mus star".

07-18 by itself if doesn't do much: I want codes "close to old successful codes" that are not simply extensions of Σ old codes. One way to get this: say old codes are N bit strings. Use BZZ to make extensions of old codes that are of length $\sim N$ & try Revo out. Complete corpus.

Or, for Backhaul, ~~for~~ for all code lengths $\leq N$ bits: Try extensions of the Σ first $N-k$ bits of old codes, using extensions of the old corpus.

I guess T. problem of SUMAC update is (other than Backhaul): To make trials that are "close" to known, ~~with~~ Σ successful corpus. Have a type of being ~~with~~ acceptable codes to augmented corpus. The R system (Σ Σ (of "Report")) is one way to deal w. this.

30 Re: Augmentation of BZZ, using d -neigists! A neigist is used for moda, just as Σ neigist is used. Hvr., standard Revo. have to agree on how far to stretch for suitable neigists. The further to stretch goes, the longer it takes for Σ to do decoding - but we get ~~more~~ more compression.

Row 421.12-18; 422.00-04
429.31-34
437.00; 438.00
I don't know (implications problems)
12 Rows 30 pp
Then summarize how far Σ use
Gather w. each of them.

TM3

$-40 \times \frac{3}{5} = -72 + 32 = -40$

Group (USA) #36
New Mexico
Botswana 2004
Benchoche or 1994?
Budapest (1996)
SP 1995

The "2 kinds of PD's" V.S.
The ~~idea~~ Paradigm of a ~~single~~ developing, changing during the career of G. Scientist. Does PD1 (copy) change? Does PD2 change? Do both change?

One approach: As the experience of the scientist develops, his corpus (\equiv TSQ) grows. The ~~idea~~ is Now to Universal D.F. applied to the known corpus of good pd for continuation. This assumes ~~some~~ $C_B \in \infty$.

For finite C_B 's, we the Scientist have a "Summary Mechanism" (\equiv Sumac) of each P.M. history: The Sumac looks a bit like an ~~idea~~ but it is able to do "Back tracking". So it ~~is~~ is not exactly a "Continuing" Sumac."

PD2 guides "Lsra" over the current ~~idea~~ & current Corpus. Note that for PD1-PD2 idea is a ruff model of human Problem Solving; - (It's for "Phase I" only)

D7.02 $\frac{399.17}{10}$ PD1 is a Practical Univ. D.F. - it is the P.D. on future contents of the known Corpus is incomplete. PD2 is a method of approximating PD1 as a computable function. PD2 grows with Corpus, just as PD1 does. One such method is a PD guiding an Lsra to find short codes for X.

Corpus, C_1 . The Corpus for PD2 can be regarded as part of the ~~idea~~ overall (Corpus of TM).

PD2 Corpus is closely related to a "current problem" that TM works on.

C_1 is the normal Corpus, C_2 is the PD2 Corpus. Perhaps the Corpus $C_1 \cup C_2$ should be used for problems! [In general, the complete trace of TM's activities should become part of C_2]

D10.02 C_2 .] So: TM has been Solving \geq TSQS! $C_1 \in C_2$, we can realize them into a single ~~idea~~ TSQ, BUT we furnish "External" Context into that $C_1 \in C_2$ are for. Somewhat Different Universes. This into into be in form of INDEX's, furnished by Trainer.

At the start, TM keeps the corps separate, but later may decide to Merge them converging degrees. Note that PD2 has a "special corpus", in that it includes part of traces of TM, so

it really has to be locked in a special way to prevent TM's Unwanted "Illegit. Self-Reinforcement"

39 744 cubes = 5 96 = 6 x 24 = 6 days: #24/day
414 cubes = #1.

4:30
8:10 P
4:30
2:34 P

00: 434.40: 433.31 - 434.28 Gives a recent take out. Context ideas of 402.12-22.

With t. "Self context" ideas of 434.33, this looks like a kind of understanding of context.

I do want to list various kinds of context & how to use/discover them.

[SN] There is a usual connection of "Bayes Seq" & "Laps rule" in most kinds of implementable "context".

So is a very simple probability Model ... except that context is inherently complicated.

1) "Universal".

So: starting w 402:12-22:
a) raw token frequencies (no context, uncond. prob) - ^{New} But tokens can be defined.

1) ngrams preceding t. Token to be predicted as conditions of p of next token

2) ngrams (see 434.03 - 14 for discussion) for "soft" ngrams see 434.33 - 40 - (The distinction of soft ngrams is important.)

2) The R functions of SA of report are ngrams. (They are for SA of Q's assoc w. Rest R exp) 425.10 it is attempt to define probabil N's: 434.33-40 on soft ngrams. Is relevant & maybe the better one 425.10

3) Other Grammars: a) The "context" is usually dependent for t. token to be predicted. An imp. Counter case is NPM. "Kasler" data. (Use of FFT to detect that kind of "closeness")

b) The "state" of a system producing all or part of t. corpus. - say a finite state (or F. finite state) machine: Hidden Markov Models

19 → 4) N.P.: T. wtd. Ngram is t. most general kind of predictor (it = ALP) on HMM's. It is t. most general kind of internal context. (But note 446.22-30)

We need good methods of deriving such ngrams (P. ngrams).

Contexts (like for Raster scan corpus) (but have gaps before t. relevant part of t. corpus at t. token to be predicted, (0/5) (3/2)) can be represented by an ngram in which all ngrams immediately before t. token will be predicted.

Note: While the adequacy of t. wtd. ngram is important as a means of

Universal means of representing context, it is not the only way. Any good well defined

predn. method (e.g. Boolean w. deltas of tokens or BZZ) can also be used

421.12
Start of review of how phased works
N.P.
Autd Ngram
to most general predictor
... t. most general context

TMS

Context: continued.

00:45-50: Ways to elz. context is prog. These elzns: Usual context (forms of context to predicted token), One Grammar is in ndim raster scan, when parts near to predicted token have to be found in

a special way (fft makes it easier to detect context of this sort).
Ngram
ngram set = (ngramst).

103 **D103** contexts: expanded to ~~macro~~: This makes size (in fact) larger! - perhaps very imp.

ANALYSIS is more the most important destination of a request. - otherwise, if size $\rightarrow \infty$, it would be poss. to use individual ngram components of an ngramst. as contexts.

Synthesis In general, we will be forming ngrams & breaking them up into smaller ngrams/or ngrams.

This is the analysis/synthesis aspect of the progress of science.

Analysis is breaking up of ngrams, synthesis is creation of Ngrams. This process

continues w. the "Breaching Up" & "Synthesis" occurring different ways each time they are invoked.

We need techniques (functions?) that do Synthesis & Analysis.

We try these functions on: data in ~~order~~ (in a priority) order, & see ~~what~~

14 Which of them give compression (= produce).

WB My original application of Context was to Modify the pc's of "Macros" (rather than Tokens).

This is, of course a more general than Tokens, but is about the same if one has already defined a token to be a Macro. Mainly, the original necessity of Context considerations for

20 "Tokens" (= Macros) was the "Scaling" problem. That w.o. context considerations (i.e. uncond. pc for all tokens/macros) - that pc's of solving problems \downarrow very

rapidly as one progressed in the TSO.

The recent discn. of Context seems v.g. in sense of suggesting several imp techniques to describe phase 1 - maybe "more or less" adequate. I think the context discussed were

28 all d-functs, d-ngrams. They are somewhat's, since they have varied approx percentages of they obtain different pc's for production, depending on their past successes.

30 I do want to review this in some detail - I don't want to forget any of it. Is it all of it

93331-434.28?

While context can predict any future config., it's usually easiest to get contexts to predict tokens (i.e. primitive or defined symbols).

402.12-22 is an early discn. of 6 kinds of contexts

35 On "soft" ngrams: A poss. (w/delta): We have an ngram in which each poss. ngram (of all poss. ngrams) is given a pc. (which may or may not sum to 1) - These pc's give us to be used in predicting

a certain token (or set of tokens). To obtain such a set of pc's (w/delta), do it backwards:

Use corpus backwards & see (for each token) what the previous preceding config. is. From Bayesian pt. of view, these pc's can be inverted to get the pc of the token in view of the past.

A "soft" = "w/delta" ngram is most genl. kind of Contexts, concept, PD \rightarrow 435.19

TM3.....

- 00 (500) (431.40): Tl. sources of f. data (Boxes) ~~are~~ stipulated ^{to be} ~~mut.~~ ^{mut.} exclusive pts for T_s.
 Each w. its own a, u, g₂.
- 02 S. c₁(g) of report looks reasonable. If there are no failures, we can get a = 1 i. f. the Gummel-Gummel d.f. Hk data. If there are no successes, we can get a = 0 i. f. the Gummel-Gummel d.f. Hk data. If there are no successes, we can get a = 0 i. f. the Gummel-Gummel d.f. Hk data.
- 03 Gummel d.f. Hk data. If there are no successes, we can get a = 0 i. f. the Gummel-Gummel d.f. Hk data. If there are no successes, we can get a = 0 i. f. the Gummel-Gummel d.f. Hk data.
- 02 i. 03 assume some kind of default apply for a, u, g₂. With different apply, no success in data would not give a = 0. Also 100% success would not give a = 1. We might get a value of a by STEINish Analysis - pooling lots of past data with other ^{kind} pairs (430/21).
 Perhaps we can just get our "a" values from Stein i. get mixed by other, ~~more~~ more conventional means.
- 10 STERN is one way of dealing w. t. "Zero frequency" problem. The foreg. (w. STEIN) seems like a good way to deal w. the situation in which we have a bunch of u, l pairs (430/21) and we want to assign a prob. (a, u, g₂) to each pair.
- 12 Hvr, if we try different a different set of {R_{pbm}}_s, {R_{psr}}_s funcs, we get a different ^{subset} of ~~a~~ $\langle a, u, g_2 \rangle$ for each u, l pair but (prob. will usually be a different no. of u, l pts) **NO!!**
 - Prob. no. of data pts is r. same. - How do we compare Bayesian fit, to Reason 2 situations? We end up w. a different no. of successes & failures for t. a different R function sets. **NO**. t. no. of successes & failures remain invariant. This info is assoc. w. ~~the~~ t. same set of empirical trials on both sets of "R" funcs.
- 20 This means that we can multiply the probs of success (to get a prob. density) by the probs of failure, (which are actual) - T. resultant product is comparable to any other product of same no. of successes & failures. - This is a t. Basic Update problem in Phase 2: (iteration - 24) ^{is/zn of it}
- So: -12 iff. very imp. when we are trying to fit a bunch of possible {R_{pbm}, R_{psr}} functions to a batch of data on success/failure of various trials/S.
- 25 T. development of '00ff & its preceding stuff 430.10 - 40 ^{is} a way to do some of Phase 2. ~~Abstract~~ How to Generate R_{pbm, psr} conts & how to do t. search efficiently, remains to be discovered! But (25 P) can use parts of R funcs obtained in solving t. d-problems of QATM - phase 1. - which may be a imp. way to use MTM development to help solve NMTM problems.
- 30
- 31 Still very imp. problems. How B22 discovers contexts; ~~what are~~ ^{what are} limitations of context? ~~undiscoverable~~ ^{undiscoverable} by B22 & how can we expand them? (usually to Universal status.)
- 32 T. Univl. context sets: How does t. past (probabilistic) determine t. future? In B22 t. contexts are always Noms & t. count to t. tokens to be predicted. A simple formula can take care of it.
- 35 is an ngmsl. - ex. T. previous token was an ^(odd) integer. Another/work be a state of a ~~system~~ system (ex. a Hidden Markov Model) that "state data" ~~could~~ ^{need} not be in a finite state Machine - states of a UMC are ok. - i. t. TRMC, T. read head has "States". Also context can be more distant (as in raster scan)
- 37
- 38 From 33: Any marked predn (future or s-funct of past) ~~defines~~ ^{defines} a kind of "Context".

or some computer (orbits) w. t. taken to be predicted.

(SN) Definition reduces cc, but are only done when parsing is quite clear & B22 is fairly large. B22 will do OSL, but ~~does not~~ ^{does not} have a final predictor in OSL. // Since they obtain a specific parsing of Cooper - that may be wrong. (437.00)

-C:\PB35\TM432.BAS

```

dim a##(6500)
a##(0)=1 : n=0
10 x##=0
for j=0 to n
x##=x##+a##(j)*a##(n-j)
next j
a##(n+1)=x##
print n+1,x##,6/(4-x##/a##(n))/(n+1)-1-1/(n+1)
if n=100 then end
n=n+1
goto 10
    
```

$$n+1 \quad \frac{6}{(4 - \frac{2n}{2n+1})^{n+1}} - (1 + \frac{1}{n+1})$$

$$n+1 \rightarrow n \quad \left(n, \frac{6}{(4 - \frac{2n}{2n+1})^n} - (1 + \frac{1}{n}) \right) = \phi$$

$$\frac{6}{(\quad)^n} = 1 + \frac{1}{n}$$

$$\frac{6}{4 - \frac{2n}{2n+1}} = n+1$$

$$4 - \frac{2n}{2n+1} = \frac{6}{n+1}$$

$$\frac{2n}{2n+1} = 4 - \frac{6}{n+1} = \frac{4n+4-6}{n+1} = \frac{4n-2}{n+1} = 2 \left(\frac{2n-1}{n+1} \right)$$

$$\frac{2n+1}{2n} = \frac{4(n+1)-2}{n+1} = \frac{4n+2}{n+1} = 2 \left(\frac{2n+1}{n+1} \right)$$

4328.

$\frac{2n+1}{2n} = \frac{4n+2}{n+1} \Rightarrow \frac{2n+1}{2n} = \frac{4n+2}{n+1} \Rightarrow \frac{2n+1}{2n} = \frac{2(2n+1)}{n+1}$
 T. I forgot zero for $n+1$ then 42, 48
 It's likely to be exactly true for all n.

This gives 2 much faster way to compute the sequence.

$$\frac{4n+2}{n+2} = \frac{2(2n+1)}{n+2} = \frac{2^n \prod_{k=1}^n (2k+1)}{\prod_{k=1}^n (k+2)} \quad \text{So later } 2^n$$

So we could write $a(n)$ as $2^n \div (n+2)! \cdot 2 \times$ product of even numbers.

Product of odd nos $= (2n)! / 2^n \cdot n!$

$$\frac{2^n (2n+1)!}{n! \cdot (n+2)!}$$

So $\prod_{n=1}^{\infty} \frac{4n-2}{n+1}$ is always an integer ≥ 1 .
 as $\prod_{n=1}^{\infty} \frac{2n-1}{n+1} = \prod_{n=2}^{\infty} \frac{2n+1}{n+2}$ } seems very unlikely!

does this imply $\prod_{n=1}^{\infty} \frac{4n-2}{n+1} = 1$? If looks like that

$$\frac{2}{2} \cdot \frac{6}{3} \cdot \frac{10}{4} \cdot \frac{12}{5}$$

$$\prod_{n=0}^{\infty} \frac{4n+2}{n+2} \quad \text{is interpreted as } \prod_{n=0}^{\infty} \frac{2n+1}{n+2}$$

$$\frac{(4n-2)(4n+2)}{(n+1)(n+2)}$$

$$\frac{2 \cdot 2 \cdot (2n-1)(2n+1)}{(n+1)(n+2)}$$

count no of even nos. in num \div in denom. in $\prod \frac{2(2n+1)}{n+2} = \prod \left(4 \frac{2n+1}{2n+4} \right)$

$$\prod_{n=1}^{\infty} \frac{2(2n+1)}{n+2} = 1$$

$\rightarrow 442.00$

D 6.0E
TM3

432E

67	2.203373E+37	3.911764
68	8.621893E+37	3.913043
69	3.374855E+38	3.914285
1	1	1
2	2	2
3	5	2.5
4	14	2.8
5	42	3
6	132	3.142857
7	429	3.25
8	1430	3.333333
9	4862	3.4
10	16796	3.454545
11	58786	3.5
12	208012	3.538462
13	742900	3.571429
14	2674440	3.6
15	9694845	3.625
16	3.535767E+7	3.647059
17	1.296448E+8	3.666667
18	4.776387E+8	3.684211
19	1.767263E+9	3.7
20	6.564121E+9	3.714285
21	2.446627E+10	3.727273

196 230 280
142 5.46 2.140
1.453 (1.1710)

TMB

4 PM Ya-humy : Grace : N: 30.03 ... Sunday

422.40 - not good!
420.90! There are 2 series: one $< R^{(n)}$ one $> R^{(n)}$.

The $<$ seq. ~~A(n)~~ ~~A(n)~~ ~~A(n)~~

For $>$

$$A(2n+1) = (A(n))^2$$

$B(2n+1) = (B(n))^2 \cdot (2n+1)$ ← No! It happens that $B(n)$ is a lower bound for R as well!

$A_0 \equiv \alpha$ say $1, 3, 7, 15, 31, \dots$
 $\alpha, \alpha^2, \alpha^4, \alpha^8, \dots, \alpha^{2^k} = (\alpha^{2^k})^{\frac{1}{2}}$

i.e. For square forms, when it exists, it is the best of $2n+1$ terms, for b. cases, you need

$k = 0, 1, 2, 3, 4, 5$
 $n = 2^k - 1 = 0, 1, 3, 7, 15, 31$

$$A_k(n) \propto \alpha^{2^k} \alpha^{2^{k-1}} \alpha^{2^{k-2}} \dots \alpha^{2^1} \alpha^{2^0} = \alpha^{(2^k-1)+1} = \alpha^{2^k}$$

$$B(n) = A(n) \cdot n$$

$B(0) = \beta$
 $\beta = 0.319327$
 $\beta = 0.319327$

It seems clear that for large numbers of 1 seq. squaring is a much higher source of growth than will by $2n+1$.

I looked up: say 1 1 2 5 14 46 140 ...

int. nat (integer sequences) — Apparently

it's not in f. (I tried looking up "1" but it didn't help).

Is it related to the seq. giving no. of funcs of a vars of depth k ? (Which fit into "closed form" in Szab. (1990).

We know that R major Majorana $X_n = \frac{1}{2} |X_{n-1} + iX_{n-2}|^2$

We can find the largest root of each n spec. char eq. by the lower bound on $R(n)$.

$$x^2 - x - 1 = \frac{1 \pm \sqrt{1+4}}{2} = \frac{1 \pm \sqrt{5}}{2} = 1.618$$

$$x^3 - x^2 - x - 2$$

$N = 68$

D6.05: Dim $A(100)$ $\beta(100)$

$$A(n) = 1 : N = 0$$

$x=0$: For $j = 0$ to N :

$$x = x + A(j) + A(N-j)$$

Next j : print $N+1, x, x/A(N)$

$$A(N+1) = x$$

If $N = 100$ then END

$$N = N + 1 : GOTO x$$

has 15 decimal precision
on 10 computers ± 4932

overflow after $N = 69$ 3.52 + 38, j 3.94

I understand on 420! 5+14=19 (not 21)!

scrap 432 $\frac{1}{2}$ for printout

$N = 200$	5.12200	3.99
$N = 1000$	2.170 + 597	3.994
200	3.32 + 19199	3.997
400	1.54 + 2408	3.9985
650	1.056 + 3908	3.999077

± 37
 ± 307
 ± 14932

$$\frac{6}{(4 - 2(n)/2(n-1))^n}$$

$N = 6000$ 2.78 ... $E+3606$

$$1.00016 = 1 + \frac{1}{2n}$$

For $N > 730$ it blows up!
For $N > 650$ it blows up!

$$\frac{6}{(4 - 2(n)/2(n-1))^n} \approx 1 - \frac{1}{n+1} \approx 0$$

 $\approx \pm \sim 10^{-6}$
 $\approx 10^{-20}$ for $n < 20$
 $\approx 10^{-19}$ if $n = 90$
 $\approx 10^{-17}$ if $n > 100$

For $n = 1, 2, 4$

$$\frac{6}{(4 - \frac{2(n)}{2(n-1)})^n} = \frac{6}{(4 - \frac{1}{n+1})^n} = 4 - \frac{6}{1 + \frac{1}{n+1}} = 4 - \frac{6n+6}{n+2} = \frac{4n+3-6n}{n+2} = \frac{-2n+3}{n+2}$$

If I did it at exact precision in Maple or Mathematica, it would probably be exact in all cases.

T013

Electricity cost: 8.2¢/kwh. 4:30 PM. Backfire.

4:30.40: So: T, Q, B, is \dots t. MITM w. T's q from \dots in Mipke \rightarrow maybe hard MITM problems, 2 good beginning (study problem) for MITM in general, is Phoca 2 in particular? Is it a good "study problem" for Long. English?

[SN] How to change only \dots d-function in S-funct! In each step, there is a "next token" that is chosen. If there was a hidden token in d-pm will give pc = 1 to 1 of the choices each time. If we make a less sharp d.f. we get an S-function. — Not so easy! If we start w. \dots d-function, then \dots after ϵ first choice (in S-function) of a token that is not that chosen by the d-funct. Then for each subsequent choice of tokens, we can no longer use the choice in d-function as a guide, since \dots all \dots state of the Machine has changed. (Motor control kinds of function traces, onto certain kinds of outputs. This is not so.) However, assoc w. each d-function, we have a pc for each token, that was obtained when that d-funct was discovered. One way is a "introduction" function.

4:30.31 \dots purely \dots best way to deal w. this problem of S-funct construction/deriv: \dots a \dots w. desired solns. Then \dots Do they \dots various kinds of S-TSP's: Then \dots from this mix of solns. Then see how to \dots (deduct, act) of \dots solns to d-problems can be used to \dots S-problems. (4 ff \dots like V.G. ideas. — V.G. Approach to Solns of S-Individ problems)

4:30.15 ff \dots (see also discuss today to 4:30.15) A better way would be for me to write out solns (\equiv "h(T)" functions) & try to factor them 2/o factors \dots reasoning back (od to them, so Th is able to \dots (find/search for) them \dots

4:30.29 Rq: 4:30.15 - 2g: Use th. Epim funts to map ϵ problems to \dots kps space. " " Rps " " " " " kps " " Then use continuity points \dots out. combined $k_p + k_s$ space \rightarrow $\langle \dots \rangle$ vector functions on that kps space.

Also, "World Multibim Scaling" be useful? (I know little about "MDS") — Maybe lookin Google?

4:30.28 If $\alpha \beta$ is reasonable: Then for each ϵ of $\alpha \beta$ cases one must have at least \dots ! I think \dots of ϵ of $\alpha \beta$ cases \dots must have at least one success & at least \dots failure.

\rightarrow (C): \dots if ϵ first success for one of ϵ cases is at t_0 , then isn't \dots also done for failures for $t < t_0$? ?? !! Go over eq. (9) p 16 of "Report" very carefully! Remember that "h"-type d-functions can have \dots meanings: I think \dots chosen for $h(T)$ \dots is \dots prob of first success at time t . Is $(1-h)$ \dots expression for ϵ vs. \dots of $\alpha \beta$? Perhaps if I did this I wouldn't \dots to have both success & failure cases!

Say we have a "case count" plot for a large no. of trials. \dots times \dots Both successes & failures: \dots 433.00

Tm3

Electrical Engineering
Costs Range ~ ~~8.2 f/kwh~~ 8.2 f/kwh. 4:30 PM. Back fire.

430.40: So: T. Q3, is ~~it~~ t . MTDI w. T3 of from delays in Mips \rightarrow maybe hard MTDI problems,
2 good beginning (study problem) for MTDI in favour, is Phase 2 in particular?
Is it a good "study problem" for Eng. English?

[SN] How to change any d -function on S-funct: In each posn, there is a "next token"

that is chosen. If there were a bidist token d posn will give posn 1 to 1 of choices each time.
If we make a less sharp d . we get an S-funct. — Not so easy! If we start w.

a d -funct. Then ~~the~~ d -function let us first choice (in d -funct) of a token that is
not that chosen by d -funct. Then for each subsequent choice of tokens,
we can no longer use the choices in d -function as guide, since d all over state of Machine
has changed. (That certain kinds of function trees, are certain and others. This is not so.)

How, as we w. each d -function, we have a posn for error token, that was observed when that d -funct
was discovered. One way is a "infodistance" function.

430.31: purely d . best way to deal w. this problem of S-funct construction/deriv:

430.31: a top w. desired solns. That is, Do this for various kinds of S-TSP's:

Then gene. from this mix of solns. Then see how to tokens (S-defns, etc)
of t . solns to d -problems can be used to help construct solns to S-problems.

(A ff-scheme, long V.G. idea. — V.G. Approach to Solns of S-Inducta problems)

430.41: It is realized (a bit) in 430.15 ff. (see also discuss leading to 430.15)
A Better way would be for me to write out solns (\equiv "h(T)" functions) & try to factor them

2/o factor t. measuring Dist (od to them, so TM is able to find/search for them effectively)

23: 430.29 Re: 430.15 - 29! Use th. Rpm funts to map t problems to k hypdim space.

Then use continuous funts ~~to map~~ onto combined $k_p + k_s$ space

\rightarrow $\langle 2, 1, 1, 0^2 \rangle$ vector funt from a Part hypdim space.

Also, "Word Multidim Scaling" be used? (I know little about "MDS") — Maybe lookin Google?
Multi-dimensional Scaling

30: 430.28 If $\alpha \beta$ is reasonable: Then for each of t cases one must have at least one success!

I think for each of t cases we must have at least one success & at least one failure

Q: if t first success for any of t cases is at t_0 , then isn't there also date
for failures for $t < t_0$? ??!! Go over eq. (9) p 16 of "Report" very carefully!

Remember that "h"-type distribns can have 2 different meanings: I think t uncertainty chosen for $h(T)$

is first probab of first success at time t . Is $(-h)$ finite expansion
for t vs. t of eq (9)? Perhaps if I didn't have to have both success &
failure cases!

So we have a "case count" plot for a large no. of trials. Both successes, failures:
or cumulative no. of cases, as a function of time.

5800
433.00

TMB

$$R(s) = 1; R_1(s) = R(s) \cdot (s-1) = (1 - 1 \cdot 200 \text{ on } \dots \text{ "best problem"}$$

$$R_2(s) = R_0 \cdot R_1 \cdot (s-1) \cdot R_0 = 2$$

Guess: End at 130pm

431.1611
 May 60
 xerox copy

00: 429.40: **ON NMM**: S. Powers: One way to implement Ram w. "3IU" is to have 2 inputs define f : discrete params of f function. — T "R input (s) (R can be \rightarrow ("R" input) filter. Continuous params of f : s. function. We use Ht Carlo method to get PC of corpus wrt - such as s-funct. If f : function has several continuous params (i.e. several "R" inputs) we get PC of corpus ~~in some~~ Ht Carlo way: It may be useful, sometimes to find γ .

06: actual points) \hat{a} (their) with $R(s)$. — say for by precision PC's — like SM, act \rightarrow like (00) Ht \rightarrow 110
 So .00 could be one PST for induction; BZZ (circumventing) can be another. So TM must learn to use γ into each problem.

06: A everpresent problem in (00) — fitting of continuous params: The apriori of these params — which (among other things) tells how many params to use. I guess one can usually get γ needed "apriori" of continuous params by looking at γ : past. If no γ exist, then perhaps one's should not be doing predn in this domain \odot . To do this for f : general problem of finding optimum curves w. continuous params: 431.10

15: 429.19: One (ruff) way to do this: we have a set of "R" functs that look at problems, i.e. put them in γ categories, $[R^i]_{PBM}$ we have w. R functs for PST's. $\rightarrow [R^i]_{PST}$; we can empirically get $\langle \alpha, \mu, \sigma \rangle$ for each R^i_{PBM}, R^j_{PST} pairs by averaging over cases: At first, γ w. of R's in each of γ : 2 categories will be quite small (~ 2000 PSTs, - 30 errors (pbm problem types))
 Ram as (52) γ we can however.

20: Given a set of (PBM, PST, T) PB: To derive γ 2 sets of R functs, $[R^k_{PBM}], [R^l_{PST}] \ni \langle \alpha, \mu, \sigma \rangle_{k,l}$ can be found that is a good predictor. Perhaps γ "h" equations (eq 9) of γ : 2.2 of 30 Oct 2003 report. will be useful: we want to Max γ is. γ : γ
 T. problemok 15-24: say we let α, β of $[R^k_{PBM}]$ is γ of $[R^l_{PST}]$. α = cardinality of γ ; β = cardinality of γ . γ : 5... 55... 55...

24: Then we have $\alpha\beta$ pairs: α, β status $\langle \alpha, \beta \rangle$. This is γ $\alpha\beta$ params — which is very large, so we'd have trouble optimizing (eq 9) w. that amt. of info in it: \rightarrow 431.30
 28: — If would be far better to find a way w. only, say γ $\langle \alpha + \beta \rangle$ params, say: \rightarrow 431.23

30: A nice thing about .15 Ht is that it is able to use f : γ functs γ other sub-functions (that were obtained γ : QA induction of γ : 2 of re part) to get useful "s-functs". \rightarrow See 431.20
 31: γ (N) A (params) \rightarrow common ways to make d-R functs: Taking γ γ stoch, or any 3IU (like BZZARAZ) is threshold it. — This is γ what's to gain converting IDV problem to γ problem.

34: 429.35: ON Generality of BZZ: Can it be used on 2 dim. raster pictures? Rissanen may keep discussed Ram his early paper on use of context for predn. T. relevant amt of context could be discovered using γ : autocorr of γ : original signal. — No autocorr for just a few values (for SM, say) is rather unappealing. A 2D raster is equiv. to 21 request (see 448.16)

TM3

ORGEVOLN: 28

Σ 111

00: (500/428.90) ends. This effectively reduces to CJS. So the effect on search speed is a factor of $> r$ but ^(usually probably) $< r^2$.

02: 425.02: In Maple, etc.: TM will learn about solve(x, expression), Integrate(x, expression).

Mean: We can then go off into a TSG. to teach how to solve/integrate progressively more difficult expressions. If this is done properly, TM should eventually be able to solve, integrate really difficult cases — perhaps beyond present state of symbolic Math.

08 One Q: How far can we go in Phase 1 Also: Can I put this ATM to solve more general kinds of Math problems of Ar by diffy?

09 → A Big Q is to how to go from MTM problems of .02 etc, to NMTM problems etc

10 In particular, the kinds of "s-funct." problems needed by Phase 2: — Involving parametrization of DEF h(↑) functions: $h(\text{param}, \text{PST}, T)$. h can be derived by 2 or 3 common

Param Problems

12 Continuous params: μ, α (often, but not necessarily μ).

→ To do .10 it would be well for me to "try it myself" ("By hand"). It would seem that T: "R" funct. ξ (of "T. Report", (in Genus of Mem! (See 445.10)) would be useful here in exam telling what type of problem param is — "what kind of" PST, PST: is. Then somehow, we have a funct that looks at the pair of categories, & finds ξ way to assign μ, α to them. This assignment could be done by purely statistical studies of past category pairs. 430.15

19 20 → Another approach to NMTM is to make a TSG for NMTM problems & try to find a TSG of MTM problems that would be a good fit to it.

22 BZZ is an imp. kind of NMTM. Could we get TM to discover BZZ & analyze/improve it? The Genus of "Context" of 402.12-22 two ways to Genus BZZ so that it is perhaps as general as it could be — But I really need to do a lot of investigation before I can do this well. I have listed 4 common kinds of context there. There is a BIG JUMP to General forms of S & E — first has to be filled. (Note 425.10 General predictor for context)

28: 427.10 SN Evolve: John Nash's work on "local equilibria" in Game Theory has been applied to Org evolution — in which para determines the equilibrium pt. This

30 analysis Must be relevant to my (understanding) analysis of SOMACS (426.08)

31 (N2509) GENERAL State of Project! If BZZ doesn't work well to implement phase 1, then, I will improve BZZ w/o use AZ or some other "compression" scheme. Similarly, T: redundancy reduction may be easier diff, but I can compute CJS, using any particular method for "redundancy reduction" (= "degeneracy explain" = degen. explain). Major problems (not necessarily in order of diffy or "unsolvedness"): Biggathra modified CJS.

- 35 1) Does BZZ implement degen & other AZ tricks adequately? If not, fix (Also see (22)) → 430.34
- 2) Good ways to deal w/ "degeneracy" → Also other Math problems of (08)
- 3) Good TSG for NMTM — Is long of "solve", "improve" etc. (interesting?) (Abstract)
- 4) Link (3) to NMTM problems / TSG's — in particular (10)

tijeras = scissors

00:427.19 : If BZZ did not do something like that, it would be a rather inefficient code.

01 say $B = b_1 b_2 b_3 b_4$; $C = c_1 c_2 c_3$. To code "B, C: sum of B is RPN, not Polish!")

$b_1 b_2 b_3 b_4 c_1 c_2 c_3$ & sum: we code $b_1 b_2 b_3 b_4$, then $c_1 c_2 c_3$ (end symbol) \rightarrow (Noty. 103)
then sum... we code $b_1 \dots b_4$ (indep string) to code c_1 , we can use context

(if any) $b_1 \dots b_4$ & c_2 has context $b_1 \dots b_4 c_1$ acc / c_3 / "A" symbol
has context $b_1 \dots c_3$, it ~~should be~~ "sum" has context $b_1 \dots b_4 c_1 \dots c_3$.

I think this implements OSL - e.g. if sum is 4, end of ...

sum had occurred only once before, it would however pc increased over c_3 never having
around before. \rightarrow T. pc of BAC should be same as pc of CAB - This would be
True for $CB=00$, but for $CB \leq 00$ it will often not be true.

How would it work if I coded it backward (Polish notation)?

? \rightarrow If I did BZZ on various Corp - would doing Reverse make any difference? \rightarrow
Try English try List Expressions, "C" genes, etc.

To code a seq. of characters, make a Basic table each character followed by a Carri. Res.

When it reverses order: check that Basic doesn't screw up; look at tax codes

On t. problem of 427.15-19; 428.00ff: I could just look at the function trees that I

want to code & devise a sequential method of coding them, that is somewhat like BZZ

[If the "improved BZZ" has been programmed for speed/low memory use; try to get copy; & use its tricks to speed up my own program]

The ideas of 00-010 on the sequential coding of trees, would seem to be fine here!

Since functions are represented by Trees, wouldn't a stack be a v.g. to store & represent, & code them? One normally uses a stack to do Tree Stack. Here we have to trace that represents a function; & also a tree that represents our search over all possl.

Function trees.

SN On symmetry of function "Lsch"; By Not recognizing equivalent redundant ("degenerate")

forms, we ~~lose~~ lose twice ① for redundancy r , we waste time by testing something r times.

② If the redundancy is soln. is r , we ↑ its CJS by a factor of r by not considering redundancy.

The 2 effects do not multiply exactly. I think the second effect is more important, since search time \propto CJS: ~~But first effect also wastes time~~ The first effect does waste

time, here. If r redundancy of a tree is r , & its row pc is p_0

t. take CC_0 for a row test, then we wasteful testing, we spend time $\frac{CC_0}{r}$ if on that case. If we recognize r redundancy, we can skip $r-1$ out - so we do save time.

The fact that its pc has increased w. redundancy recognition does not affect this

result - we do save that could otherwise waste search, here, because it gets higher pc.

\rightarrow Also, I think the normalization constant is modified when we reject a bunch of r

$$P_n(x) = \sum_{i=1}^n 2^{-i} \delta(x, x_i)$$

$$= \sum_{i=1}^n 2^{-i} |S_i(x)|$$

Spec
426.07
or
426.07

00 : Org. Evolution: Recently (on Internet ~~in~~ NZ0,03): Evolution of microbes in 3

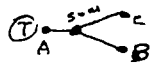
runs: Variations betw. 3 populations were "very small": Certain problems never solved in slightly different ways — I would expect more diversity ~~but~~ but betw. 3 results

Hvr, since ~~the~~ Environments did not change, the problems presented were minimal, (if at all). Some kinds of "pabs" that ought to be tried: Slow/rapid change of pH, temperature, nutrient composition, poisons,

T. reason for lack of diversity betw. 3 populations: that "problems" solved only used a few solns. By ~~then~~ trying a greater diversity of problems, more significant results

Mite be obtained. Hvr, ~~since~~ Env. was constant, we had "neutral (mutational) variations". 429.28

1) BZZ is Redundancy & adequacy for representing subsets, etc.
Say I'm looking for cases of "sum & mul, (.)". If "sum & mul" occurs.
(. is betw. mul) (i.e. really), I would miss it — So How do I do (w. this)?

2) General uncertainty of whether BZZ gets imp: redund obtained by "definitions".
e.g.  similar context of token "T": A sum B is a context, but say "c" was

by pc definition ... i.e. it occurred frequently in previous corpus.
BZZ would seem to not use this fact — which should \uparrow compression.

Well perhaps it does use this fact. B & C are coded as compactly as poss, then sum B, C is coded 428.00

20 : ~~420.40~~ $R(n) = 2 \times R(n-1) \left(\frac{1}{k} + \frac{1}{k^2} + \frac{2}{k^3} + \frac{5}{k^4} + \frac{14}{k^5} \dots \right)$
 $1 = 2 \times \frac{1}{k}$
 $1 = 2 \times \frac{1}{k} \left(\frac{1}{2} \right)$

It is sum \rightarrow slowly
 $\propto n$ for large k !

$1 = \frac{n}{k}$ so $k = \frac{n}{2}$ — for large n $\frac{n}{2}$ — seems r/t to for very large n ,
but for small n , k seems to be $\approx 3, (1, 1), \dots$ No! $k \neq \frac{n}{2}$ for large n ; ~~the~~ $k = \frac{n}{2}$ was obtained by

assuming k was constant for large n :

say $R(n)/R(n-1) \approx k$ for large n , $R(n) = 2 \times \frac{R(n)}{k} \left(1 + \frac{1}{k} + \frac{1}{k^2} + \frac{1}{k^3} + \dots \right)$

say $R(n) = k \cdot R(n-1)$ and k is function of n . $R(n) = k \cdot R(n-1) = R(n-1)$

$R(n) \approx \left(R\left(\frac{n}{2}\right) \right)^2 \cdot (n)$ $n = 2^k$ $R(2^{k+1}) = \left(R(2^k) \right)^2 \cdot (2^k)$

try $R(n) \approx R\left(\frac{n}{2}\right)^2 \cdot n$ $n = 2^k$ $R(2^{k+1}) = R(2^k)^2 \cdot 2^k = \left(R(2^k) \cdot 2^{\frac{k}{2}} \right)^2$
 $R(2^{k+2}) = \left(R(2^k)^2 \cdot 2^k \right)^2 \cdot 2^{k+1} = R(2^k)^4 \cdot 2^{3k+1}$

exponents $R(2^k) = 2, 4, 8, \dots$

" of 2 $k, 9k+1, 7k+3, 15k+7, 31k+15$
 $k, (2^2-1)k+2^2-1, (2^3-1)k+2^3-1, (2^4-1)k+2^4-1, \dots$

Take $R(0) = 1$

$r = R(0) : r, r^2, r^4, r^8, r^{16}$ $R(2^{k+1}) = R(2^k)^2 \cdot 2^k$ $R(1) = R(0)^2 \cdot 2^0 = \frac{r^2 \cdot 2^0}{2^0}$
 $\rightarrow \frac{r^4 \cdot 2^0 \cdot 2^1}{R(1)} \left| \begin{array}{l} r^8 \cdot 2^0 \cdot 2^2 = r^8 \cdot 2^2 \\ r^{16} \cdot 2^0 \cdot 2^3 = r^{16} \cdot 2^3 \end{array} \right. \left| \begin{array}{l} k=2 \\ k=3 \end{array} \right.$

TM3

Org Evoln (00)

Adequacy of BZZ .20
Function of Backtracking .20

Adequacy
Adequacy

3 or ev. Expts.

.00: (4.14.31) : I had this idea of Org. Evoln. w. a sudden gross change of envt., leading to Mass Extinctions" — corresponding to a SUMAC Given a diff. problem, so few of its old ~~old~~ codes survive.

.07 → In Org. evoln., the remaining life forms have not complexity in their evolutionary search, so it is essentially, "Universal" If envt. changes slowly, envt. it will regain its former diversity.

Org. Evoln
427.01

.08 In a good Sumac (≡ suitably designed Sumac), the same sort of thing should occur: There should be cut backtracking (or Equivalent) after a hard problem to restore "diversity" (≡ no. of somewhat different solns to present total corpus).

.09 [A Sumac can also be designed so it only has to solve all the past problems. — SOMETHOW or, use a ~~win~~ window rather than a Rect. window] → 429.28 → consumed org. evolutn.

.10 As a rule of SP: Global warming is a slow change of envt. — most ecological systems will adapt; Man will also adapt, but on his usual way of ways between coalitions. Due to the present in near future technology of war, this can be very disastrous — perhaps ^{Much} more than Global warming.

.20: .09: Essentially, the Function of Backtracking is any technique that ~~restores~~ restores diversity to an adequate level. (BZZ may be able to do this). Compare Backtracking Method & BZZ in effectiveness: Can we get cross fertilization from the 2 technique types?

.24: (425.40) Spec: The function that takes $(O_1, Q_1, \dots, O_n, Q_n, A_n)$ & proposes a pd on A_n is the most General/Context internal, & I have considered it — But it needs to be etcd. — i.e. its "too complex" — i.e. say we are looking for a function that maps $(O_1, Q_1, A_1, \dots, O_n, Q_n, A_n)$ to O_{n+1} — we put the argument in as one of the "initial arguments" for A_n (copy) & look for function that map it to A_{n+1} (probability). This function must work for $n=1$ to n . As stated, the CJS would be far too large for n of practical size — but it may work for $n=1$ or 2 or 3. Using Phase Functions as initial approxs, the keys to get function for $n=4$, then $n=5$, etc.

.24-.32 Any Elemental could furnish sub-functions that could make the job of more feasible. The Method I'm currently considering for doing phase, involving token freqs, data (params) & BZZ is expossi elem. of the problem & aff

TM3

TSQI ¹⁰⁰ **: V.g.!** ^{0.00} → But note 0.06

Genz, CONTEXT: RESOLN. ^{0.57}

↳ internal

Geno
suffin.

37.37
11
48.375 v. >
54.548

425

41 p 120
25.6 only
5.3 part
5. fiber
42.9

4926 : I that of TM loosely t. deduc of various expressions in Maple, Macsyma, Mathematica, ...

as a good basis of a TSQ. On a more practical level, doing those deduc, would

Train Me a lot in TSQ construction! — So: V.g. Idea! ^{Much} ^{More!} → (429.02) →

The Ques are MTM problems, They can be a basis for N MTM problems as well (see 420.03 - 19)

→ One Q is how to find N MTM problems that are reasonable continns of MTM problems — i.e. — ~~420.18-19!~~

I May Invent this: Find imp N MTM Problems, then look for MTM problems that use n cases.

Maybe a "R" functions of \mathbb{S} of "T-report" would be a good example: Tho I haven't yet defined

probabilistic R's, it would seem like a v.g. Idea to do so.

A Quick Attempt! A R funct is a p.d on all Q's, w. peak values on certain set of $Q \in \mathbb{S}$. — An immediate problem! Do we want a p.d (normalizable) or proby values that don't have to be normalizable? — Could non-normalizable proby of this kind be applied to situations \mathbb{S} like (uniformity $x \in \mathbb{R}$ $x \in \mathbb{R}$ $x \in \mathbb{R}$)?

Anyway: Also assoc w. each R funct, is a p.d on $P(A|Q)$ distributions. ~~XXXXXXXXXX~~

Again, I'm not sure I want (or need) this p.d to be normalizable.

T. way these R functs are used: A Q_{int} comes in: All R functs have wts. (probab) for that Q_{int} also have wts for all $P(A|Q_{int})$. From this info, we get a bunch of wtd. P_i functs \rightarrow we could get a proby def. over all A by using a suitable wtd member overall to $P_i(A|Q_i)$ — ~~or~~

From a practical point view, the computations of 16-19 may take too long so

We have to use Approxns.

Re: to. normal/non-normal of 11-15: T. pc's mita to normal vert. For R's, ^(I'm not sure this is a way.)

so $\sum R_i(Q) = 1$. for each possl Q. This makes R categories mut. exclusive \leftarrow (??)

Or there may be some other way which $R_i(Q)$ ~~can~~ can be normal.

I wonder about Fuzzy (Zadeh) sets. Are they "normalizable"? — My guess: If so, they are distant from proby in an essential way.

30 (424.37 spec) On the d-functns, R_i , as examples of contexts that are sets of Ngms, w. simple deduc.

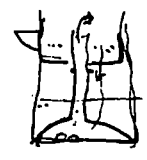
I was thinking of each R as being defined by a set of Q_i that it "accepted" — This is a peculiar kind of "context"! — But it is a legit "context" since it "conditions" a set of $P_i(A|Q)$ functs.

(N21.03) I may want to allow t. & d-R functs to overlap — so any "input" is a legit context ~~XXXXXXXXXX~~

35 (N22.03) On a General Q of How to implement various kinds of context implied by 402.19 (\exists t. full Univ. set of contexts), \leftarrow These are (so far) only "internal" contexts.

37 That 402.12-16 Don't nearly cover the set of cases implied by 402.19, means that Rest is a large Universe of contexts that I have yet to discover... so in teaching TM ^{435.19 more or less solve this problem.}

How to Inv. TSQ's, I will be able to employ new context types. ~~426.24~~ ~~426.24 spec~~ 426.24 spec



Spec 423.40: "In 2 kinds of PD": PDI ~~is the 2nd assigned to~~
 strings & sets of ~~codes~~ that describe. seems to all of the problems. In SUMAC, we try to find
 codes of hyperny that describe corpus. "2 kinds of PD" is one way to do this by using
 a gradually modified PD₂ (≠ PD₁, usually) to guide ~~the~~ modified "L search".
 The "close" to previous trials is obtained by BZZ ~~using~~ ^{operation} $O_1 \dots O_n$ corpus. This selection of trials,
 while still "Universal" is also heavily biased toward ϕ part. If ν , this bias is "right enough" so
 that (personally) specific "Backtracking" is not needed.

N2003 In SUMAC w. a finite word codes to formal, there is a sharp division (cutoff)
 between "stored" & "not stored" - which makes explicit backtracking necessary since the
 code needed for a new example (here) may not be in the stored set of codes.
 (Or, when, using BZZ, the set of codes stored is "soft", all possible codes
 are included to some extent. An rule ϕ is, "Is the Closeness Criterion used by
 BZZ using corpus $O_1 \dots O_n$ a good Closeness Criterion" -> to a near
 exclude the desirable codes for ϕ here, but we can give them very low pc's....
 Somehow, we want this latter tragedy to be very unlikely. - In general, the
 "Gardner's Rule" we can never be sure to avoid it!

In General, the problem of generating the best possible PD₂ (which is the
 "stored" set of codes) is a very important problem. BZZ w. corpus $O_1 \dots O_n$ is
 our guess.

SN Re: (19-20) one way to "adjust" the sharpness of PD₂'s "peaking"
 about the $O_1 \dots O_n$ corpus is to sharpen or flatten the individual Bernoulli Dist's
 (within BZZ/Ret substitute BZZ). (This method is not 100% perfect, but does move
~~the~~ the rate character ϕ - one easy way to do it is to \uparrow sizes of denominators
 of pc values. It amounts to stronger a priori equality of all tokens - a larger
 SSZ for the "equality" - Another possible way is to take ~~the~~ the roots of all finite pc
 values (i.e. renormalize, unless renormalization is automatic - which it often is) - \leftarrow the root may be a positive
 code for a step approach to accuracy is not needed).

Back to Weaknesses/Bottlenecks: 423.25, .27 seem most severe at present.
 This is about contexts that are OK of a set of Ngrams. 405.31, 397 I Rich 405.31-40 is about the only
 useful writing I've done on this: [A perhaps not very useful!]. I intended to idea that ordinary,
 linear numerical time series, linear (or h.h.) regression was example. - But I Ret this/ was more like
 a PST. The R's of S_{ϕ} of ϕ is an example. The R's relate a class of examples (S_{ϕ} 's)
 to a set of PD on ϕ should not map $\phi \rightarrow A$.

Maybe the R classes have sharp adjcs. ~~the~~ s function R's would be more general,
 probably more realistic - maybe easier to update (?). 425.30
 \rightarrow See 425.10 for a quick attempt at ~~the~~ prohibitive R functions.

TMB

00: (422.40) spec: So + Q is: Can we implement 422.38-40 by a suitable seq... or is it best to

"Tell" TM (ruffly) how to solve fi. problem by putting in O's in "R" form.

Perhaps think about .00-02 bit: There may be other w/o better way.

One way: Start out using T R system & get a bunch of O's.

Then just use BZZ w. currs of these O's to get PD2.

Similar technique to 420.02-07. A poss. trouble: If there are only O's in 1. currs that use R system, fi. best to use fi R system

Subsequently, will be very by (no never 100%): If we (no, we can (illegally) give O's currs w. R's loss w/ permit should have, to encourage "originality":

There may be a way to do this!

15: 422.11 Present bottlenecks: I'd like a clearer understanding of relationship of "pure context (BZZ) & "Definitions". Are they really almost equiv? Could it be that defs are actually more "correct" (better compression)?

Specify (often completely) fi. parsing which can be a big restriction. It looks like context includes definitions plus other codes in () So perhaps to max

(or only) advantage of defs, is lower C (which can ultimately translate to higher PC, by effectively ↑ in C B). Evaluating this trade off is how & when to do it, sounds diff (Maybe unsolvable by any exactness)

25: 422.20; 422.14: This is one aspect of Context that may not cover full contents of context of 402.12-22 402.19 in particular.

28: SN While BZZ may be "power" Equivalent to "defs" of simple & multi (25) Sorts of "deduction" may be diff. To what extent do these differing "pc's of definition" depend on initial choice of a pipe?

31: SN Re: 2 kinds of PD's v.s. "SUMAC" model of probabilistic knowledge Acquisition In Sumac, w. each new problem/Prob. a prob. is included to include w/c contributed by Prob. Problem Soln. If CB=∞, fi. system enters prob. best possl. predic. For finite CB, there are many possl. ways to proceed. One way is to start by retaining the best 100 codes for fi. currs. Systems could find in limited time. When a new problem is given, these 100 codes may be reduced to 10, 1 or none. In any case, we have to search for codes "close" to BZC ones that formerly worked, or if any known codes solve fi. problem, look for codes "close" to them, that also solve fi. problem. Keep "diversity level" at 100 (at least).

00: 421.40 : Introducing Definitions commits one to a certain parsing of the corpus.
 422 uses all poss. parsings (I guess) — so it is certainly better in that respect.
 On the other hand Definitions markedly reduce the no. of tokens in the corpus, so searching is much faster. To compromise: we don't use definitions, but only after they have a fair S&Z, & it is clear that they give very little info.
 [Even so, it may well be that we would want to retain alternative parsings as

Is this "summary detector, utilized" = some kind of "Heur P (418, 39)"? Is it kind of "Quick Abort"? (pre-concursion about @) Or, it is a way of modifying the PDZ int. sense of changing order of bits.

11 codes (or maybe alternatives O^i 's?)

10 **WEAKNESSES** At present it is not altogether clear to me what BZZ "covers" definitions completely (other than the CC aspect) → 423.15 → Misses 4.8.

12 Another (seems stronger) weakness of BZZ is it doesn't properly include contexts that are the "OR" of a set of NGRMs. An example of that kind of context is linear (is most non-linear) time series prediction (Also, "Mark function of S1 & Report" → 405.31-40 for context given. 405.33ff)

14 A third (More General) weakness of BZZ in this application is (421.19R) PDZ doesn't look at $Q_i A_j^{n+1}$ in deriving a PD for O_{n+1} . I think (non) linear regression does look at, but these techniques are more like PST's than PD's! Perhaps "weakness 3" is one part of the character of Phase 1 → 423.2

20 Disen of Weaknesses: (2) (12) can perhaps be fixed (if it is true & if it is serious). (3) Looks VERY SERIOUS. ^{Actually not} serious. How far one can go w/o. directly looking

22 at the problem is unclear. BUT we put in O^i , the res possibility of looking at the problem, is deciding what to do. — So essentially what we are doing is looking for a class of functions that can do this. Our inductive corpus is a set of O_j that have been successful in the past for smaller parts of the corpus. ^{in this kind of problem}

29 One move in direction of .22 is to "R" system of new § 1.2 of the part. ^{for QATM} Report.

30 The R's look at the problem & decide what class of P's to use. → This is a particular form for O^i 's. The "R" system is Genesis of it and of Much more of Phase 1

It is probably necessary to have a good understanding of .22-.29 when designing TSO's. → It is possible, that in view of .22-.29 (3) (15ff) is not really a legit weakness of BZZ !! I think I got involved w. this argument some time ago & finally realized what was going on! — It took quite a while to figure it out at that time.

32 N18-03 One (non-er) way to deal w. the problem of .22-.29 is to simply look for O^i 's that do what is needed, and hope the system can find suitable solns. — which will automatically include solns. for .22-.29 & etc.

00: 420.40! I want to go over the ideas of Context again: See how I'm using it as general way as poss. — see how Part B23 is doing it properly: (Not 420.07 rats on Context)

325.19-18 is 24 (approximately) much different ϵ cuz: of "Context"!

323.11-18 various some interesting ideas on centered to general Phase problem

Review this 322.31 ff. stuff carefully: It was interested in putting the O^2 cards in good order for (presumably) a Backtrack search. I need to re-familiarize myself w. these ideas. In view of my use of (u) B22, is an effective equivalent of "Backtrack" automatic & perhaps rather good?

Balance to 322.31 ff. against my more recent ideas on how Phase 1 should work.

A Quick Summary of ^{Not so Quick!} ~~Phase 1~~ T. most Recent ideas on How I want Phase 1 (or QATM) to work! QATM's Goal is to Maximize $\alpha = \sum_{j=1}^n O^j(A_j/Q_j)$: ... To find a set of O^j 's \rightarrow Maximize α subject to O^j constraints by AZ or by B22 or \approx B22. The constraints T. α is a function of $\{O^j\}$ set and constraints PD1 ... which we want to Maximize.

PD2 is the probab. of an O^j in view of the previous sets of O^j 's that have been "acceptable" .25

20: 420.10: This idea of using initial MTMTSQ as "springboard" for NMTM!

This may also work in a way for an RTM \rightarrow QATM or optimizing TM. (More "concrete" than RTM at advanced level). This presumes Part 4 or "Analogous" \leftarrow Too General? RTM has structures N to or elaborate to concs, contexts, functional pc's of

.13R (SN) that PD2 doesn't look like problem Part One has to solve, suggests Part 1 is not up. Method. to search for Opt!

25: .18 \rightarrow PD2 may also include any "external" or "internal" context into that the trainee wants to include. Concretely, the problems set a stage w. each of the part O^j 's is part of the context, but for Phase 1, we will not include it.

PD2 is used to search for O^j to maximize α of (.13R). B22 is used to obtain PD2 ~~and~~ (which is a D.f. on O_{n+1} , next problem)

T. Coups of \approx B22 is $\bigcup_{j=1}^n A Q_j$: \approx is a punctuation between O_j 's. Pre-coups of 3c is a linear string. The functions of \approx the O_j 's have been "desynchronized" by the use of 418.26 - 419.18 ... Note 419.37-38 on "how to do it". The O_j 's are bound also have their pc's \uparrow by 419.00-18.

\approx B22 doesn't use definitions as such. It uses pure ^{compression} ~~probability~~ (probably omits some ll coding that could be ... we also use symmetry (418.26-419.18) to give effects some effects of ll coding. \rightarrow 422.00

TM3

Get AL for AL Dish, plate
to low front tables
maybe 6-2

00:49:40: w. f. "Shortcut of 49.34 since probly will not know it & given TSO has exact size for d. needed contexts, to be "adequate TSO".

02 SN ~ 350.00 - 15, .30 off 2000s that a MTM TSO is not ordinarily a useful introduction to
03 \geq NMTM TSO - T. first uses 2 IU; T. second uses 3 IU. (Hvr, TT Value
of 6 MTM soln \geq TSO is that its PD2 gives us concs & contexts that r p's ch

06 Create ord directly for NMTM. any can we use concs, contexts designed for 2 IU & use them in 3 IU such.
07 \rightarrow Some imp't ideas on "Context" in PD2: $\geq 33.05, .09$ fact to 327.21: (322.31 - 328, 15)

09 .06 Details of .06 R: any Subtrees in functions of any args. can be used in
10 - functions of 3 args (subtrees are "contexts") - any part of the input function can
~~be used for 2 or 3 args~~
Woops! I don't know! Any subtree of 2 IU can be used as 3 IU
will subtrees of 3 IUs not be usable for 2 IU construction
Phases \rightarrow subtrees contain roots so r. 3rd input ("R"), R is fact. 3 not
relevant to \rightarrow (NMTM \rightarrow NMTM)
2 IU \rightarrow 3 IU

18 \rightarrow At present, I can't think of a ~~subtree~~ NMTM that would be a good construction
19 of a MTM. TSO. \rightarrow see 425.06-07 for an idea.

20 SN Q: What is the redundancy of a sound unusual N-variables?
and n variable symbols. The test symbols must be vars. In Polish, we need n-1 "sum symbols",
n-1 "sums" \geq n-2 vars. No. of first symbol must be "sum"
At no point can we have more W's than S's.
Another approach: A tree of n-1 s's. Each s can have coming out of
it either 1 or 2 s's.
A. Other approach: recursive soln.
Sum \rightarrow VV
n-2 S's; n-2 V's
Missing p's!
 \rightarrow 308, 304
Missing p's!
 \rightarrow 308, 304
Missing p's!
 \rightarrow 308, 304

30 R(k) = $\sum_{i=0}^k R(k-i) \cdot R(i)$?
k = n-1
For problem's use, we just compute R(k) & make a table for
TCU to use.
NB R(n) has to be mult by n! or (n!)! to include
N-23.03 If asymptotic ratio exists for n \rightarrow ∞ :
Seq: $1/k = 1/1 + 1/2 + 1/3 + 1/4 + 1/5 + 1/6 + 1/7 + 1/8 + 1/9 + 1/10 + \dots$
 $1/k + 1/(k+1) + 1/(k+2) + \dots = 2/k$
 $1/k + 1/(k+1) + 1/(k+2) + \dots = 2/k$
 $1/k + 1/(k+1) + 1/(k+2) + \dots = 2/k$
The seq. is defined by
 $R(0) = 1, R(n+1) = \sum_{i=0}^n R(n-i) R(i)$
If $R(0) = \alpha$, $R(1)$ is mult by $\alpha^2 = \alpha^2$
 $R(2)$ is mult by α^4 , $R(3)$ mult by α^8
 $R(n)$ is mult by α^{2^n}
I find that $R(n) < 1$
I find that $R(n) = 0$
 $\rightarrow 432.00$

TM3

00 : 418.40 : Re: Re repetition of symmetric functions identities of 418.26ff: They have 2 effects: (1) The careful repetition of trials (2) The pc's of these trials are not summed into PD₁. If a function has k equnt. codes, - we should test only one code & give that code a wt of $\frac{1}{k}$. So the equivalence data from techniques of 418.20-34 should be modified to mult. f.wts. of each Equivalence form by its Multiplicity ("redundance") of that form.

10 (2) Give much better pc's (\equiv values for PD₁). If k redundancy occurs in PD₂, it will multiply by k, the amount of $\binom{r}{k}$ we will spend generating & testing that cont. On 2 redundancy effects, in both PD₁ & PD₂, the most important of 6 effects may be a multiplying pc by its redundancy, k. It is notable that both effects of symmetry are imp. even if 8: Expressa. of interest occurs only one time in a cov. pos. If it occurs n times, its effect is k^n .

20 **SN** T SQ writing: A nice set of Concs to look, would be the meanings of various Expressn types in LISP, MAPLE, MACSYMA, MATHEMATICA, etc. eg.: Eval(x); Solve(x); Simplify(x); ? : differentiation, integration, etc. dot product, cart. product, Matrix x vector, Matrix x Matrix ...

26 Solve (n covs, n unk) \rightarrow linear, then non-linear. \rightarrow 425.00

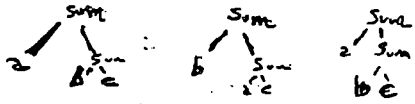
I will have to analyse these \uparrow concs to see what derive on adequate Conc net (a associated helping Contexts) - To an imp. extent. B22 does consider

30 Contexts - But I may need "external Contexts" or other kinds of Context that B22 is unable to furnish. Remember that "Generalized Context" is any aux info that causes the pc of a next token to be a "conditional" rather than unconditional probab.

34 405.31 is one kind of genus of "context": 402.12-22 gives various Context types, 401.15 a discn. of context. It is tempting to try a "reasonable T SQ", using a PD1 & PD2 both evaluated by B22, & hope that there is some "context" info to facilitate "learnability", but, it is more reasonable to try to get search contexts that I know of, that would be adequate for T SQ. - Not so easy to do! One of the imp. things in knowing whether a context is adequate, is knowing what its SSZ is.

(5000
411.26)

00: 417.40: On the commutativity problem; Consider to product $a b c$; $sum a sum b c$; $b sum a c$; $c sum a b$
we all legal — giving 3 way redundancy! So the Lexical order trick isn't solving that problem!



A different tack! We are trying to function in 2 pc. order.
 say $sum a sum b c$ occurs early in function, then $sum a sum b c$ is known likely to occur later.
 " $sum b sum a c$ " " " " " $sum b sum a c$ " " "
 — but $sum a sum b c$ will not occur (likely).
 To report write by first we see b (write P.D. quote all functions) but we repeat identical (functions many times). — (True, but not new)

So: if $sum a sum b c$ occurs, how to prevent $sum b sum a c$ & $sum c sum a b$ occurring?
 * $sum a sum b c$ perhaps convert immediately to $sum a b c$; sum is a symbolic over any finite w. occurs
 $sum a sum b c \rightarrow sum a (b) (c)$ if $sum a b c$ is right side w. sum ; we convert
 to $sum a (c) (b) (c)$ \uparrow say third variable sum w. sum ; we convert $(1, 2, 3)$ to $(2, 3)$
 sub go to previous rule.

More generally if any side of sum starts w. sum we do a conversion to sum ; in which τ falls how many ways start w. sum . Well! — perhaps 20 is way to deal w. some kinds of symmetry. As soon as certain seqs occur (int. generation of functions) T seq. is immediately rewritten.

In generating a function, each token goes on to stack, along w. certain "state" variable values. When we rewrite $2 \pm v$ fix, we pop. fr. stack a few times...

26 SN There are other equivalences that we may want to realize in our function generation:
 1. associativity of addition $sum a sum b c = sum sum a b c$. (This is replaced by $sum a b c$)
 distrib of mult. $mult a sum b c = sum mult a b c$

The trick of 20 never always rewrite a string into distributed form. This can lead to long expressions, but it does make equivalences clear.

There are many ~~other~~ symmetries — to associativity distributivity in Logical functions (like AND, OR, NOT, IF \Rightarrow (IF \Rightarrow \exists input (choice function). But may be able to be dealt w. by 20 off. Also Note TRIG Identities — 3 kinds of common identities

34 (N1503) Two imp. ideas! (1) T. idea that possibly equalize should vno. of knos. \exists
 In present case, this is done by writing all expressions in "standard forms".

37 \rightarrow This problem probably has been solved in "Symbolic Math. Technology" exp.
 38 \rightarrow Macysque, Maple, Mathematica.
 39 (2) Plus looks like a new type of flow:
 (1) big one,
 (2) quick - A bank,
 (3) changing P.D.

General Equality Problem
 Symbolic Math
 No w. type of heap.

SEM

.00:416.40! $C(z, T)$ is fast for small z , but drops rapidly for $z > eT$. (e is some constant)

.01 For study $C(z, T) = 0$ for $z > eT$, but more gradually C drops more slowly on T .

If C drops like $\frac{1}{z}$ for $z > eT$ then for most P_i , $\sum P_i C(z_i, T)$ will converge.

C includes a normal constant so $\frac{N(T)}{eT + z}$ $N(T)$ is a normal constant

$C(z, T) = \frac{N(T)}{eT + z}$ $N(T)$ is a normal constant
This $C(z, T)$ will work for most P_i 's. We may want $C(z, T)$ to \downarrow more on loss slowly w. T ... depending on the particular P_i dist.

If $P_i = \frac{1}{z}$, say $(z_i = (1, 2, \dots))$ then $C(z, T)$ can decrease much more slowly than $\frac{1}{z}$ for large z : $\frac{1}{(\ln(z))^2}$ would be ok. or slower things like $\frac{1}{(\ln(z)) \cdot (\ln(\ln(z)))^2}$ or whatever could be used.

In general, if we know the divergence rate of $\sum P_i$ we can design a $C(z, T)$ that causes $\sum P_i C(z, T)$ to converge w.r.t as slowly (or as rapidly) as we like.

The simplest $C(z, T)$ function is the step cutoff at $(.01 e)$

In Longshot^o in a lottery, new words are invented as $\frac{1}{N}$ - but

the situation w. corpus is different. We have a corpus of N words selected at "random" from a "very large" corpus. The no. of distinct words in the corpus will be a f function of N . For each value of N , we have a sample. By noting how the distribution varies w. N , perhaps we can guess at the dist. for $N = \infty$.

These may be used in those books I have on frags. of words in English text, to figure this out.
Type-Token Mathematics: Horden | Computational analysis of present day Amer. English
Lang. & Chances Chance

.27 Perhaps use BZZ w. GA Populations corpus, to create new trials.

The mechanics of using BZZ on an "unworded corpus", would have to be worked out.

One way: Each case (i.e. each case of success actual trial) starts w. symbol Δ .

.30 We look at corpus to get Δ d.f. of failure for that symbol - which is "say, Mark

Optimal Monte Carlo - using Δ as context, we obtain next symbol & act.
"Marko Carlo"

Formally, we ~~can~~ probly do BZZ or close to it: But from practical point

.33 view, it may not be so easy! One possibly easy way: Use as "sequencial corpus" all cases, bounded by Δ 's. \rightarrow we could use BZZ on Monte Carlo generate

Cands. We may be able to give wts on previous cases (depending on Pair Cases (I know how))

\rightarrow (.332) If sound similar (if not identical) foray plans to use BZZ in Phase 1. \rightarrow as 406.31 ff
(see: 406.31 ff ... jumpover & ~~data~~ analysis of "Lexical order" to X Allios ff. \rightarrow 438.10 for More on BZZ in GA

00 ⁵⁰⁰ 415.37 Re: proboset 415.37 : A possy: Look at present problem. from past problems select ^(neighbor) ~~(one)~~ (or more) closest to it. From population of solns. of past ~~problems~~ neighbors, move back to earlier populations; γ . But ~~not~~ "move back" is greater for neighbors that are further away.

05 Use same Mut, cross as last neighbor.

Cases where $\omega \rightarrow \infty$ will not work much. ~~T. case~~ we have a linear filter design problem. Say the only previous "cases" are filter designs but for much different Optza criteria — so we expect solns to be much different.

Another approach: Say we know as set of $\{Q_i, A_i\}_{i=1, \dots, n}$ pairs obtained by GA. Use some kind of rule induction to get a very broad A_{n+1} & if from Q_{n+1} : ~~broader~~ Q_{n+1} broad D.F. is the initial population for soln. Here Q_i decodes the ~~2nd~~ problem in terms of its "fitness func".

17 \rightarrow Another quite different approach to QA induction via GA!

We're looking for 1 (or more) $Q_i \rightarrow z_i = \prod_{j=1}^n Q_j(A_j | Q_j) = \max$ — An induction problem. We get solns for small n , via GA, then retain populations of "n" for soln of "n+1". This "soln" looks very much like a SUMAC soln. The Backtrack capabilities is obtained by retaining in the population, ~~several~~ ^{retaining} several \rightarrow approx solns for ~~retained~~ earlier problems. (How far back ~~problems~~ are depends on "Backtrack depth" we want to consider).

14-03 Discussion of GA: A prominent diffy is apparently extreme inefficiency.

First, we must assume all trials have bounded cc . We can get rid of by outlawing recursion or "loops" in perms. — This cuts down the "universality" a lot (!).

Using ~~from~~ ~~Monte Carlo~~ ~~made~~ ~~W. P.~~ P_i as prob of a trial is rather bad if P_i is prob of success of that trial, — which case is execution time = P_i of rate such.

On the other hand, P_i is guiding such that ~~need~~ not be P_i of ~~rate~~ as a cond, being a soln. It could be unimodal. — which is sort of optimum, but is convergent to converge P_i .

But if P_i is prob of soln. is no use of $\sqrt{P_i}$ for search, however we do P_i if

$\leq \sqrt{P_i}$ diverges? (I don't know how to do a Monte Carlo on P_i since probably useless $\leq P_i$ converges) One (not so good) way is to have P_i probly distrib. very w. each

trial: so, say $P_i = f(t, i)$ ($t=1, 2, \dots, \infty$). $\lim_{t \rightarrow \infty} f(t, i) \rightarrow k(i)$ ~~then~~ \tilde{P}_i ($\sum \tilde{P}_i = 1$)
 say $f(t, i) = \tilde{P}_i \cdot c(i, T)$ ($c(i, T)$ is a "convergent" function) so $\sum \tilde{P}_i c(i, T)$ converges
 "convergent"

05

10

17

20

30

Has
→ BICAKB ←
Foods of world ?
Deat (of).

.00: 41490 of having 'Several Children.

Each If one bids to solve a problem: One wins money & over solves it; One loses it & one doesn't. At "family time" one is given money: = Δ + one's chance to use for sold ground having children. (one's father can't say so).

Why have children & sell ground? (See literature on & life for reasons).

Anyway, there would be 2 tsq's (a) T. interaction, fight for life v.s. computers

(b) T. Governance problem-solving TSQ.

T. interaction of these 2 TSQ's would be of much interest.

T6 system is mind of Hillis's competitive GA's to solve a problem in computer design.

Some remarks on forgo: We could have parasites' feedback for problem solvers to get fed, then eat them!

How much interaction would there be between facilities in an organization for taking off predators/parasites & facilities for working problems?

We would want something corresponding to "publishing" that would encourage individuals to (share/distribute) knowledge.

More generally, we have various "Mature produ systems": How can we use one or more of them to construct a "Strong AI"?

Some Mature Systems:

1) GA's Sim, Annealing,

2) Ant-Like

3) "BZZ"

4) X (from 1 to 3) in a net (or LARGES Corp).

5) Hidden Markov Models?

6) ANN (or RNN?)

7) SVM for categorization (Vapnik)

8) Tree branch selection for categorization (ID3)

9) Linear/nonlinear Regression.

10) Theorem Provers?

11) Non-linear Optim. for Repts.

12) & such?

13) Univ. D.F.s

14) Clustering

Less Mature

1) TSQ ideas

2) MT (much xft)

3) HMM - because initial selection of states is not mechanized (?)

4) RNN (Crazy slow)

5) GA's - very slow

6) ANN

We start by asking how each & lacks full strong AI capability; then see if one or more of the technologies will fit in.

Well (1) GA: Given problem: If it's a NP problem, a good fitness func^{ness} has to be devised.

For optz problem (even inverts & d. continuity). After often need to modify fitness func^{ness}.

Selection of initial population → Diversify mut, & crossover rules → Spec 9/6/00

Org Evoln: 118

.00: 413.40 Not included in the "definition" routine: Now, it appears that classes of PC's of tokens are all varieties of context: that is, $B \geq Z$ should be able to deal w. P procedurally well — along w. definitions?

.10 **SN** I had Pizz idea of "SUMACS w. Backtrack". ~~Backtrack~~ with $CB = \infty$ we don't need backtrack. w. $CB < \infty$, we have a certain no of sols for each pt. in the TSO. ~~Answer~~ w.o. Backtrack, we start out w. K solns, & w. each new problem, f. no of solns \downarrow (i.e. $K \downarrow$) until it $\rightarrow \phi$... no soln.
 \circ However we deal w. this is that if at any pt. we have < 100 solns, we backtrack until f. no. of solns is ≥ 100 for 1000

.17
.18 In Organic Evolution we have a similar situation. We have a certain diversity population (an ecology in a forest). A "problem" is an external modification of the ecology. If modification is very large, we have "Mass Extinction" of species. If no other big changes occur, the no. of species will slowly \uparrow to perhaps its pre-catastrophic level.

Normally, the external ecology changes slowly enough that the population sub crossover can maintain diversity level.

If diversity level is low & a "Big External Mod of ecology occurs" — it could destroy all life — for perhaps a long time.

.30 — to compare how they deal w. new problems. A "Mass Extinction" corresponds to a "very diff problem" in SUMACS..... \rightarrow 426.00

There are several systems that can be usefully compared to the above 2.

Art life G.A. in various forms,

Could we modify A Life so it would be designed to solve problems of increasing difficulty: is it a TSO? One way might be to have a set of organisms. Each organism can recognize & solve one or more types of problems. It's a kind of "PANDAMONIUM". For solving a problem, an organism gets to participate

$60 \frac{k}{q} = 15k$

$52 \times 5 = 260 \text{ words/pr.}$

.00: 412.10 : Models to see what problems arise, & how they can be dealt w.
(Remember this is Phase 1, & Prediction doesn't have to be VERY good - only good
enough to get to Phase 2.)

.03 [SN] In Designing R recognizers; TM is first given indices that it tries to associate
w. a particular R's. Later to do "R categorization", we have this corpus of

cases in which we know Riter R's. ~~we can try standard categorization~~
OR SVM's (Support Vector Machines) or use BZZ to categorize
pgrams (like IDP) (IDP can make a word in this case) in a soft way & - like Li & Vitanyi did, by training on known corpus & adding

asking for additional coding cost of various poss. additions to that corpus.

.10 Therefore say each R_i has a corpus $[Q_i]_{n_i}$. We form one ~~corpus~~
each corpus $[Q_i]_{n_i}$ & see how much extra code length $(\rightarrow PC)$ is added when
we add this new ~~corpus~~ to each of $[Q_i]_{n_i}$ corpi. This gives us relative PC of belonging
to each of R_i 's

.15 **The Q is:** Are the regys recognized by BZZ good enough to give Good R's?
Also note in 2.1.2 off. report (on updating R's) Various Methods of updating

.20 are discussed & probably BZZ could be adapted to other models of Perm
- That section also has updating of p functions assoc. with R's - & BZZ
might be tried there as well.

.22: .02 T. idea of 412.36 - 413.02 is not bad. It may be that ~~some~~ several of
the unpleasant effects of 412.03 ~~are~~ really deliberate

much effect, & that the ones of much effect can be fixed - perhaps w. an
A.M. "soln". I could just examine some of the effects in details with
BZZ or PC's of deductions.

[N1203] BZZ does have most (if not all) regys of AZ, so if there are bugs in that applic.,
Pats would suggest bugs in my planned implementation w. AZ.

.30 I suspect that the "corpus" should only have changed O_i 's because the
situation where we need a new O_i is one of change. - Also,
the small changes of analog parents is auto made, we don't need
 $O_i \dots O_n$ corpus Guidance for it. - So don't put Part in $O_i \dots O_n$ corpus.

So superficially, it looks like a BZZ may be able to work Phase 1 (at least)
I suspect here that it is able to find a PC ($\in O_i$) for a single R (recog) func.
I previously had the idea that there were special (local) contexts, that were

.00 : 4.11.90 : **SN** A notable difference betw. PD_1 & PD_2 : PD_2 will usually have many O_i 's in its corpus: Usually they will differ only a little from one another.

PD_1 will have much smaller, the group, corpus. **NO!** PD_1 has n , Q_i, A_i pairs in its corpus. PD_2 has n Q_i, A_i, O_i triplets (sorted) in its corpus. O_i that occurs even when it does

.03 → **Q**: Should we consider in PD_2 's corpus every O_i that occurs? even when it does not change? ~~Usually it does not~~ We could have \geq kinds of (PD_1) corpus.
(See for algebra) : Many problems with short random noise as args or few probs w. long random noise args. In first case, many O_i 's, usually + same. In (b), every problem (usually) gives a new O_i . ~~Even when it does~~ In (a) we will usually have several in connect" O_i 's until it $\leq \epsilon$ gets large count, so that ϵ - "correct" O_i gets almost all wt.

Very small changes

.10 In the case of analog prediction, should we count small trivial changes in continuous params of C_i as a new addition to the corpus.

Should we just consider changed ~~new~~ O_i 's to be in PD_2 's corpus?

.15 **What about cases in which we have many sols for O_i , each w. its own wt.?**

km $\frac{5}{8}$ mi

.15 May cost needed like! If we consider all C_i 's that were found, each on its own wt ($\equiv PD_i$), we could consider all cases that arise (?) It contains lower wt. sols that could be used for Backtracking! Just how to use the info is unclear! Perhaps like GA Population!

10' 25 m
km
~.62 mi
.619

.20 Well, the wts of the various O_i are (almost) directly usable as sample wts. But Big problem is that each O_i gets lower & lower wt as i ↑!

$3\frac{1}{2}$ ft = meter
 $\frac{3.333 \text{ ft}}{5280} = 3$

Maybe normalize ~~each~~ the wts for each i ? But what wt. to associate w each O_i set? Keep the Eye on Goal: T. predn. of C_{n+1} at max wt.

.26 One way: Have exponentially ↓ wts as one goes into the past. Exponential "depression". Many diffs w. 26! Will repetitions of a problem class "increase" i , w. wt. into the past? (i=100)

Hotel
Aurora
Holiday Inn

.30 T. General Problem here is: Given O_i seq. (with or w.o. Q_i, A_i, i) To get good estimates (guesses, est. metrics) C_{n+1} . There are some that look at C_{n+1} & A_{n+1} and derive C_{n+1} using Probinfo.

Fill out paper.

Say we did know about sequences of identical O_i 's --- perhaps including info on ~~identical~~ Q_i, A_i 's for them.

11:50
78:50
Paris

.36 One way to deal w. this: Try diff't modes of dealing w. the problem e.g. Use only O_i 's that are not different from previous. ~~Use~~ C_{i-1} Use all O_i 's; use O_i 's in which the difference betw C_{i-1} & C_i is \geq certain amount. then try BZZ (or equnt) on each of these amount.

part

.00 : I do want to look at the "slow" statistical predictors for Compression, to see just how much better they are ... just what a difference regps they can detect.

NP A common commutative operation on many inputs, is addition & multiplication. We can just put the args. in Lex order $A \leq B \leq C \leq D \dots$

.05 **Q**: Does BZZ do anything ^(better than) beyond AZ? I was thinking of BZZ as doing + searching way for good coding. But it only works for sequential codes - what I've been able to do is usefully apply it to Polish formulas.

.10 How can it be applied to QATM (Phase 1)?

Perhaps it is usefully applicable to **PD₂**: which is (I think) sequential induction. The thing to be coded & retranslated is O_1, O_2, \dots, O_n . This is a sequence of Polish formulas. So viewed in this way, BZZ may be a "BIG DEAL" in getting TM off the ground!

In working w/ Polish formulas, it would be a good idea to use a "STACK" - Not only for evaln, but for a priv's & for constraints on symmetry.

The sequence O_1, O_2, \dots, O_n does not take to assoc.

problems into account. This is not a catastrophic criticism, but it would be nice if we could somehow include in to about the problems assoc. w. O_i .

.20 the sequence $\langle O_i, Q_i, A_i \rangle_{i=1-n}$ has 1. info, but would BZZ be able to use 1. info?

Would it be "good enough" for Phase 1, not to include QA info? \leftarrow I think so: there's a discussion on this Q.

The R functions $\approx \S 1.2$ of the IDSA report do implement a context of the problem itself (Q_i).

- Could BZZ be applied to recogn. of such contexts? My impression is that the R's are

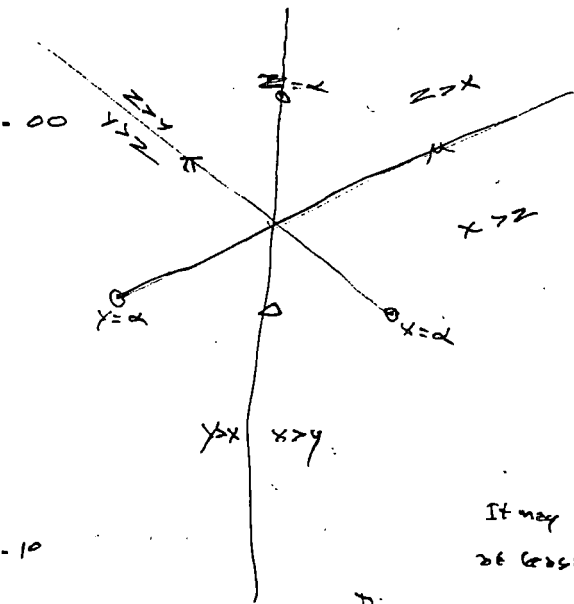
.26 a different kind of context than that which BZZ ~~works~~ works w. R. \rightarrow (18.00)

.30 **The recent** foray into symmetric funcs & the use of BZZ in Phase 1 of QATM

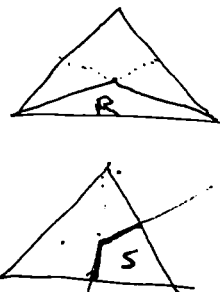
Q1 started around 405.3) in a discussion of whether definitions were okay or good or convenient because of speed or ... Could bro do very well w.o. them:

Q2 The kind of context at 405.31 is the kind used in BZZ. I was comparing that kind of context w. that of 405.34: A class of Token Contexts. The "R" funcs of Phase 1 QATM are one example of classes of context of Tokens. in this case: classes of Q's.
T. Q was: what was the efficacy of these 2 kinds of Contexts?

3 TM



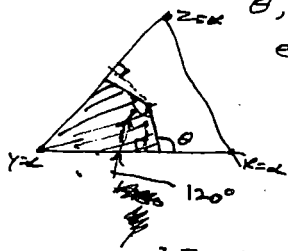
These 3 lines (actually planes) Δ , Π , Δ each divide space into 2 parts.



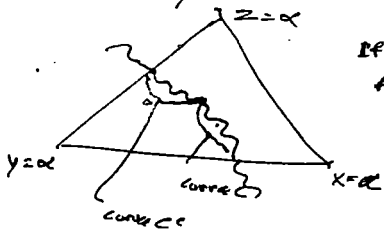
Region R is defined by $x > z$ $y > z$: in R, flipping x, y leaves it invariant.
 Region S is defined by $x > z$ $x > y$: in S, flipping z, y leaves it invariant.

It may be that if we have rotational symmetry, we'll have R or at least one flip symmetry exists (which we can choose).

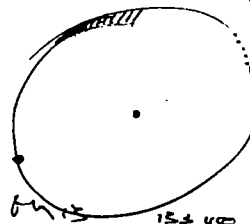
There are no π or 2π of rotationally symmetric shapes: e.g. for any value of θ , if rotation is invariant and will give rotational symmetry, but for each θ , certain flips will give invariance, others will not. In Δ figure (15C) some z, x flips leave it invariant, others don't.



If curves $\subset \Delta$ are invariant, but 120° displacement, we get rotational symmetry for Δ region Δ function R .



A circle is ^{perhaps} better space for ^{study} rotations, flips:




Well, I don't want to spend more time on this Δ Δ reference.

There are very probly mathematicians who know about symmetry groups (say n x n) that could answer all Q's directly!

So far, to get rot. symmetry we need for 3 axes, we could

$A \in B, A \in C$ ~~then~~ — this also gives ~~the~~ A, B flip & A, C flip symmetries, but not B, C flip symmetry.

Whether there ~~exist~~ functions w. rot symmetry and not even partial flip symmetries — is our question. $A \in B \in C$ gives like  only $\frac{1}{6}$ of 2π !

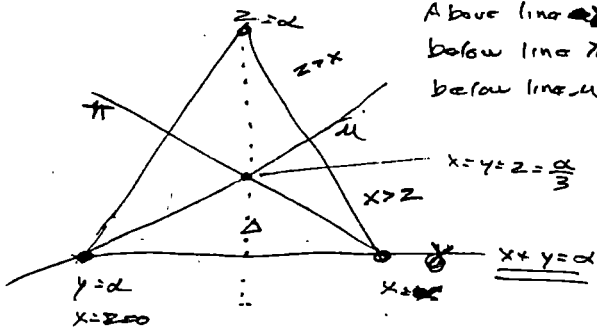
Anyway, for functions w. 2 axes (which is all we will have for Δ functions)

$A \in B$ is quite fine & will enable us to check all functions have a unique representation in Polish? So we can use Z_2 or whatever — or better "statistical" functions to detect regularities.

3 + 4

.00

A better way:



We want to region 'P' which is

Above line δ $x+y < \alpha$
 Below line π $x+y+z > \alpha$
 Below line μ $x+z > \alpha$
 $y+z > \alpha$
 $y+z > \alpha$

Since $\alpha \equiv x+y+z$

$x+y < x+y+z$ so $0 < z$
 $x+z > x+y+z$ $\left\{ \begin{array}{l} x+z > x+y+z \\ y > z \end{array} \right\} !?$

Equation of line μ : $x+y+z = \alpha$; $x=z$
 $y+z = \alpha$
 or $y+z = \alpha$

line π : $x+y+z = \alpha$

$y=z$
 $x+z = \alpha$
 or $x+z = \alpha$

$x+z > x+y+z$
 $z > y$

$z = \alpha$, $x=y=0$ is above π & μ
 $x+z = \alpha$ at this pt. so $x+z < \alpha$

line π is $y=z$ we want to side where $y = \alpha$; $x, z \geq 0$.

i.e. $y > z$

line μ $\rightarrow x > z$

line δ $z > 0$ (axis eqn. of 408.26L)

.10

.20

The condition $z > 0$ is probably not necessary (from graph of 408.26L)

$x > z, y > z$ look like $x > z, y > z$ locally, since we map Lex ordering into integer ordering.

This amounts to $B > A, C > A$ locally (probably $B > A, C > A$ is correct)

Since we want to region on or below

$B > A, C > A$ seem to give wrong results; i.e. 408.26

.30

The constraint $B > A$ was to sources \geq flip symmetry; i.e. $A < B = B < A$

Two flip symmetries gave a rotation, so that gave all 6 symmetries.

Maybe the $B > A$ & $C > A$ constraints don't combine that way!

Well, say $B > A$ & $C > A$. $A \rightarrow B \rightarrow C$ If we rotate, to! $B \rightarrow C \rightarrow A$, it violates condition

$C > A$ constraint.

If we flip $A > B$: $B < A < C$ violates constraint — so
 we don't have soln. to problem

On the other hand: ξ Geometric Soln. Seems to be correct!

So paradox!

By ~~the~~ 6 constraint $A \leq B \leq C$ (lexically) we make all permutations equal.
 However to ~~make~~ make $ABC \rightarrow BAC$ illegal; just do $A \leq B$ lexically.
 So we categorize single flip by lexical constraint.

To make all votes equal, but not any flips! How to do this?

What about ~~lex~~ constraints $A \leq B; A \leq C$. $\left\{ \begin{array}{l} \text{Note: 2 flips} \\ \text{over 2 votes} \end{array} \right.$

Actually there are more than 6 permutations! Add
 $AAB; ABB, ABA, AAA$; we assume A, B, C are distinct.

So total of 10 possys; w/ 3 inputs.

There are ~~more~~ 5 poss. symmetries: 3 flips; rotations; $\left\{ \begin{array}{l} 6 \text{ permutations} \rightarrow \text{equal} \\ \text{all } \rightarrow \text{equivalent} \end{array} \right.$

Lexical $A \leq B \leq C$ gives all 6 permutations equivalent.
 $A \leq B; B \leq C, A \leq C$ gives 3 flip equivalences.

How to get rot. symmetries?

$A \leq B \leq C$ so ~~we~~ $C \leq A \leq B$

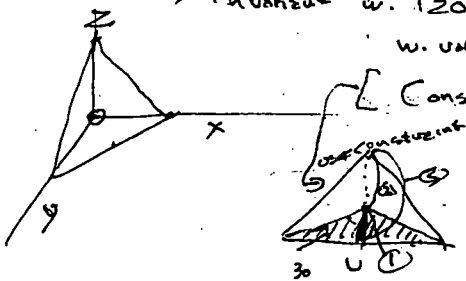
$A \leq B \leq C$
 $B \leq C \leq A$
 $C \leq A \leq B$

A possible way to deal w. this: ~~we want~~ to prevent rotary symmetries
 from occurring in our trial for A, B, C . Whenever trial occurs,
 we write out all permutations of that trial that we don't want to try (e.g. BAC).

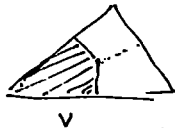
Could we do it with HASH coding?

Think of 3 axes in 3 space. Functions of x, y, z

Functions invariant w. 120° rotation about line $x=y=z$. We only have to "fill in"
 w. values, one third of space. How does characterise that space?

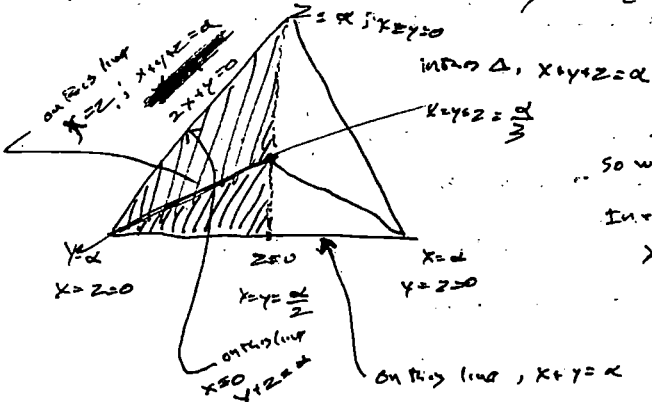


Consider a slice, $x+y+z = \alpha$.



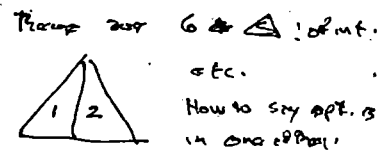
I think there are an ∞ of ways to draw the boundaries.

If A, B, C can be represented by integers, this analysis should work.



So want to prob. base $x+y+z = \alpha$
 $2x+y = \alpha$

In red Δ :
 $x < \frac{\alpha}{2}$



Group size 6 Δ of int. etc.
 How to say opt. is in one of them?

00: 406.40 Must be in Lexical Order.

Listing functions of Pairs sort in pc order: sounds diff (t, but perhaps isn't).

After "~~some~~" "sum ~~3~~" has been written, only legal symbols that can follow, are lexically ~~3~~ or later.

~~Can I use this trick for other symmetric functions?~~ Can I use this trick for other symmetric functions?

For order of args a, b, c being irrelevant, Pairs with order which can deal with it:

What about $f(a, b, c, d) = f(c, d, a, b)$? I.e. certain permutations etc.

07 08 09 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100

07 08 09 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100

What about general permutations? The idea is that if several permutations of the args are equal, then only one of those permutations is legal. (07-08 is an example - but can I generalize it to all possible permutations?).

Remember, All permutations form a group. Any subset of permutations form a subgroup.

13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 32 33 34 35 36 37 38 39 40 41 42 43 44 45 46 47 48 49 50 51 52 53 54 55 56 57 58 59 60 61 62 63 64 65 66 67 68 69 70 71 72 73 74 75 76 77 78 79 80 81 82 83 84 85 86 87 88 89 90 91 92 93 94 95 96 97 98 99 100

STUDY OF PERMUTATIONS: $ABC \rightarrow CAB$ are equal; then all powers of this permutation must be equal: i.e. BCA . To do/w. this $A \leq B \leq C$ lexically. I seem able to deal w. "special cases": But I'm not sure that the set of all perms is a group. Also e.g. in 13 & 14, I'm not sure that the set of all perms is a group. I've proposed ~~that~~ that a group rule is restrictive - so it may outlaw cases I want to have.

I.e. there are 3 equal permutations, but a total of 6 possible permutations:

20 ~~some~~ e.g. BAC, ACB, BCA are not in equivalence set of (ABC, BCA, CAB) rotations

A possible Approach. Show that all symmetric functions can be expressed as functions of symmetric 2 input functions.

In 20, there are 2 symmetry elements: each forms subgroup of 3 elements. 2 subgroups are indep. If we put these 2 symmetry elements in a group, they generate all 6 elements: No! if rotation element: applied 3 times, exchanges 3 permutations.

3 permutations. If we flip AB (in ABC) we get BAC ; if operator P applied again $\rightarrow ABC$; so all 3 reflection operators are self inverses.

So we have 4 subgroups: 1) the rotation group (ABC, BCA, CAB) 2) flip operators $BAC \leftrightarrow ACD$ 3) flip operators $ACB \leftrightarrow BAC$ 4) interchange or reflection subgroups.

So we have 4 subgroups: 1) the rotation group (ABC, BCA, CAB) 2) flip operators $BAC \leftrightarrow ACD$ 3) flip operators $ACB \leftrightarrow BAC$ 4) interchange or reflection subgroups.

2 flip operators $BAC \leftrightarrow ACD$. So if we apply P to $ABC \rightarrow BAC$ which is a flip operator. So any pair of flips generates the other set.

any flip is any set generators & a whole set.

Summing each flip forms 2 members & is a self inverse. 2 flip operators generate all 6 elements.

any pair of flips generates only 3 different values. Any combination of rotations generates only 3 different values.

ABC
BAC

00: 405.40 to Answer 405.30-40, I have to detail the way in which these contexts influence

pc's of Tokens (or Configs of Tokens (\equiv words))

04: 405. **SN** Also, I'm not so sure about 405.2-30 - (That it is poss. to have used contexts, w.o. using "definitions") - i.e. To use contexts for adequate predn, w.o. Ever using "Defns." - yet getting as good & usually better pc's. \rightarrow See (423.15 - 24): Aug. 5/84 analysis

09: 03 First I have to review use of defns in predn. My old idea of how to do defns was for at incorrect: i.e. I'm probly didn't analyze the cost of defining a token, properly! It has to do w. "punctuation cost". I was using a standard simple model for this cost, but I think it's more complex. About 4. simplest kind of punctuation is "comma" ... which tells where word boundaries are. There is the 2^{10^4} as a pc of integers, n so we could put in tokens both each word to both how long they are

Another way is to use symbols Δ . It's pc is $(\text{mean word length})^{-1}$. If Δ is pc for token by the factor $(1 - \frac{1}{n})$, n being mean word length, pc of Δ is $\frac{1}{n}$.

More generally, if we have a bunch of tokens w. pc's = $[pc_i]$, we can insert a comma symbol Δ to separate tokens. If we can assign a pc to Δ that maximizes the pc of the corpus. More generally given $[pc_i]$ set, how can we design a prefix code w. that mean pc of message? $bc \equiv -\log_2 pc_i$ so $pc_i \equiv 2^{-bc}$

07: 30 There's a better way to code a message, with code (or maybe not quite legal, using uniform words of any length L , so tokens has just $pc_i \cdot 2^L$ words assigned to it.

T. result is that punctuation seems to cost nothing!!

08: 30 **SN** Does this punctuation cost occur in error way in "HMC problem (Newby's Code)?

Well, the prob. discussion is for a known $[pc_i]$, for "like" steady state situation. T. Q of interest is not that, but cost of defining new symbols. Is that prob by (clearly not all) on symbols w. zero freq. & "redundant" relevant? Maybe our encoder may had:

Test when one can stop before of known symbols. If a new symbol had a freq. of ϵ , then its prob of not having occurred until now, is $\approx \frac{1}{2}$ if ϵ has a constant value. We can make rules for computing ϵ that maximize pc of corpus. (possibly param. arbitrary) of rules, can be selected to maximize pc of more distant past corpus.

31 **SN** B22 (arithmetic coding) - i.e. v.a. compression code for Text (strings) or they mention statistical methods that are somewhat better, but much slower

These are for Text strings. Could I arrange for corpus codes so that reps would be always reflected in linear string compression codes?

35: One thing I'd have to deal w. is the commutativity of certain functions: 29 405.33-405.11

A B^2 is B^2 are equiv. strings. : Perhaps for commutative functs, to args. are always written in lexical order! Using Polish (or Reverse Polish) into mathematical

As soon as we write a name of a commutative funct. we know that it's args.

OSL(21,22)

$\ln 1.5 = 1 + 0.5 - \frac{0.5^2}{2} = 1.375 + \frac{0.5^3}{3.8} - \frac{1}{2} \cdot 0.5 - \frac{1}{48} = \frac{1}{64}$

$\frac{2.4}{15} \cdot \frac{6.8}{30} = .227$

$\frac{1}{64} = .015625$

$\frac{1}{64} = .015625$

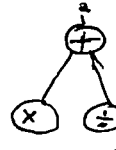
i.e. order is irrelevant

For non-commuting functions, each input is distinct

2.27
 $3 \overline{) 1.257}^{15}$



Say



has occurred, & we want to see if ~~that~~ Plus content of 2 has occurred next part.

$= 309$
 1.255
 1.2543
 1.2547416
 1.25495218

We look at all (+) nodes: we want (+) nodes w. input from (X) & (1/2) - order is irrelevant.

This may double work we heard in verifying "similarity", but this effect does not "mult by 2" every time it occurs in a tree. So it ~~is~~ seems to be a not bad effect.

There may be other diths in verifying dupl. of a certain tree in t. part.

Note! In ~~implementing~~ ~~context~~ in t. forgoing kind of way, it ~~seems~~ Seems to be unwise to make "Definitions": in fact it ~~is~~ has more freedom than a implementation involving "defs". T. advantage of Defs, is that it ~~is~~

$(1.25495218)^{1024}$
 406.35

~~Shortest~~ Shortest current code makes it faster easier to find new regys in it. The Dis advantage is that it commits one to a particular parsing of t. "known" code ... that will be "incorrect" in t. sense that there ~~is~~ exists parsing(s) that will eventually result in a shorter overall code.

Would t. forgoing implementn of Context facilitate OSL (isholmg)?

It would seem so!

A kind of compromise w. re Making Defs (Re. 12-20): Make definitions, but only after it has become clear that t. defs are very good: i.e. they ~~are~~ ↑ t. corpus p.c. much more than ~~any~~ any alternative set of Defs. This decreases t. probly that one will be able to get a better codes by using

a parsing of t. corpus incons w. those defs. → 406.04

The Context of 404.33 ff is one particular kind of context: A single context of Tokens ~~(adjacent)~~ adjacent to t. tokens to be predicted. A slitagen would allow t. confign to be not adjacent to t. predicted token

T. Genzn I'm thinking about is a class of configns of tokens that are defined in a "simple way", that condition. p.c. of a (usually adjacent) token. The previous kind of context could realize a class of contexts, but it would be much larger: ~~adjacent~~

WOOPS: ~~36~~ may be quite wrong! Say we have 10 contexts, each of same den length L, that predict l w. p.c. = p₀. Its 3 same as one ... ?? ... → 406.00

BZ2

EZTEN
try on Google.

Modified Lsqn.
 .00: 403.34 **SN on Lsqn** For random Lsqn: we do trial w. $PC = PC_2$ for duration Δ . For $CC = CC_2$ it takes $\frac{CC_2}{\Delta}$ trials. Each trial takes time Δ , so we need $\frac{1}{PC_2}$ trials to hit PC_2 once.
 So $\frac{CC_2}{\Delta PC_2}$ trials to hit it $\frac{CC_2}{\Delta}$ times. Each trial takes time Δ so total $< \frac{CC_2}{\Delta PC_2} \cdot \Delta = \frac{CC_2}{PC_2}$
 for soln. If we use $\tilde{PC}_2 (\equiv PD_2)$ to guide search to find solns.
 It will take $< \frac{\sum CC_i}{\sum PC_i}$ to find a particular soln. Then takes $\sum CC_i$ to generate all test, \tilde{PC}_2 is assigned \tilde{PC}_2 by following PD_2 .

.10: 403.40: Using PD_2 , the "soln" can be very by \tilde{PC}_2 : eg. $\tilde{PC}_2 = 1$ if we use promiscuous or Ad Hoc model of the corpus! (Promiscuous A.H. models are described in Sol/G4, for Inductive Grammars). Recently I showed (395.18-22) that GATM also has these 2 grammar types: sequential prod. also has Parse 2 types.
 However, Parse 2 "solns" while of very low CJS ($\frac{\sum CC_i}{1} = \sum CC_i$) usually give table when $PD_2 \equiv PD_1$, so the solution is normal $Lcost = \frac{CC_1}{PC_1}$ which may be fairly good PC, but $Lcost = CJS$ could be much too large. Often (usually?) in the CJS for PD_2 guidance $\frac{\sum CC_i}{\sum PC_i} = CJS$ will be acceptable, but PC_i will not be so good (never as good as that obtained by weight Lsqn).
 Perhaps Many or Most PD_2 's are alms w. $\tilde{PC}_2 = 1$. T. Compression Alms like LZ or BZ2 give $PC_i = 1$, so their PC_i 's are almost always better than those of the AH or promisc models.

.30 **SN** As corpus size \uparrow , the no of cones (i.e. the UNCONDITIONAL PC 's of the cones) must \uparrow .
 .31 w. \uparrow corpus size, we also have more contexts \uparrow in size for each context.
 Can affect .30 help ent w. validity of .30 to keep induction problem "feasible"?

.33 In counting how many times a Tree ~~tree~~ tree-corpus has occurred in the past, it's easy to take commutativity of functions (+, x, ...) into account. This seems to cause much of a search time. Can I derive a better representation, ... to help get rid of this litty?
 Consider: for multi-foreman: 2 inputs ~~but~~ but they are equal

TSQ writing: off: 419.20 ... TSQ

loop

20 : **SN** On writing TSQ's: I'd like to be able to write a TSQ, but I felt had adequate
 21 info, so that it should be irreducible w. acceptable CJS! Then mechanisms
 that are able to transfer info obtained by long to solve one problem, to be used to facilitate
 23 soln. of next & future problems. T. idea of mechanisms Genzel context (402.12-22)
 Should be helpful in devising these inductive devices

One way to write a TSQ: First make a seq. of problems that I feel ought to
 be adequate for long ^{to solve} some more diff. problem.

Next write down various ways in which t. (long or .0-.02) could take place:
 Then try to get TM able to use all of these mechanisms. In general, I should
 be able to get TM to learn any conc. that I "understand" by devising a
 suitable t.s.q. (That TSQ may have to have a very large CSZ, hvr)

~~Start w. ANL~~ Start w. ANL (to recursive eval). If I can't do it, Assume TM
 16 "understands" ANC ~~some~~ — Plus many I'm using RPM if it's taking to TM
 do define problems: so RPM is built into TM. — Its t. way we communicate w. TM.

20 : Perhaps best used Time in Costa Rica! Write very clear review of impt. ideas of last 100 pp.
 Write clear review of "T. 2 PD's" & "Context" first. This should be clear to me of
 { + 5 hrs or + 1 yr. }. Working backward to earlier ideas in these 100 pp. I. later ideas may
 significantly clarify/solve earlier problems.

22 **N.8.03** Review: P 403 is back about 100 pp. 6/8/04! understand what PD₂ is.

The "2 PD's" idea: This is (a b not counting) Relevant to IND (Induction Probs) only.

PD₁ is the 2-progs. This is a code & seq. w. a univ. lang. [AZ (≅ LISP), or FORTH
 ≅ OOPS, or whatever]. Because of t. way t. code is generated (see deriv of AZ

30 in IDSCA 2003 report: Appendix A for deriv of AZ), at any point in code sequence, we
 know it's PD₁ (≅ 2-prog). The goal is to find out or invent codes that
 desc. a particular corpus, but have as hy ≅ as possl.

To do this, one could do an L such w. PD₁ as t. "Guiding PD." This is "Normal L such"
 34 An alternative is "Manda" "PD₂ guided L such": (see 404.00-08 for Details): In any L such,
 if t. cond. → a soln, its CJS will be $\frac{CSZ}{PZ}$. $\tilde{C}C$ is cc needed to generate a test
 t. cond.: $\tilde{P}C$ is t. pc used in t. "Guiding PD" — so $\tilde{P}C$ may be PD₀ or PD₂ —
 depending on what was used for t. such.

From 2:16:30P to 10:30P 10:30P

00:40:40: Hw, since this is Only "Phase", External Context will be quite limited. The framework will specify it ~~more~~ narrowly at first: Later, many contexts ~~will~~ will be given to a problem: only a few (if any) will be relevant: TM has to determine which are relevant.

This narrowing down of relevant contexts is trivial for a new problem, since we have all the other contexts that have been already used, and TM can know which are relevant.

The problem w. many given contexts arises in "updating", re-determining which of all possible contexts are useful for helping specify the PST of a new problem:

Other (approaches) (problems) in selecting contexts occurs in the "5 update situations" in § 1.2 of the 2003 "Report" (EDSIA-16-03) (Here I'm regarding "R's" as "Contexts.")

So a list of Contexts: Types

- 1) Tokens, Reg. This is used PC (at first)
- 2) a) Data of the Mem. These are at first unconditional
- b) Other definitions ... (?)
- 3) The R type of context. § 1.2 of EDSIA-16-03
- 4) External context "This is a (Math problem / Ode/PDE...)"
- 5) T. (Genl. PD): T. procedure Tokens is a cond. PC; w. condition being all previous Data & Solns & Soln methods plus External Contexts.

6) ~~Any function like (5) but not universal, computable & not nearly any~~ Any function like (5) but ~~not universal, computable & not nearly any~~ ←

In 1, 2, 3, 4, 5, I want to see just how each case of type is implemented, discovered:

So I can parse them properly. (e.g. getting "R" from d to S contexts.)

N-703 In this study, Look at ~~the~~ $BZ3$ (or $ZB3$) ? Dit Simon BZ2?
 $BZ3$ But v.p. Compression PCM for
 strings, using "context". Site if I can ~~apply~~ apply any aspects of their techniques to more generalized "contexts": T. Contexts may use all "products" of the problem - i.e. Ray "nest". My hypothesis: Plates Janir
 even permissible contexts of larger, to compression Very sticky, slow, but worth loss, 2
 perhaps pointing to a Bug in their system. T. system itself ~~is~~ is pretty This seems unreasonable
~~improvable~~ improvable because it was oriented toward practical compression:
 So it leads to look for "Best" code, rather than use all possible codes.

Various kinds of Context can be used in both PD1 & PD2: But remember PD1 & PD2 have different Condi.

3TM

00: So o.k. Token freq. is a context.

What about "Definitions"? I guess they can be regarded as contexts, same as token freqs.
A slight difference: That token "contexts" are for tokens pc's or tokens only:

03 Definitions (e.g. "Nouns") involve pc's but extend further into the future.

So 400.39 - 401.03 says Context can include extra freqs & token definitions:

So: so far, Context includes all kinds of freqs ^{that} I'm including in PD

Modifns 888

10 400.39 ff is an impt New IDEA: I want to understand it better: — T. idea that

all freqs I've been considering in PD₂ are contexts. Actually, this isn't ~~the~~
much diff (if at all diff) from the use of the term "context" in text compression.

15 Superficially, it would seem that any PD would be expressible as context sensitive
p.e. — i.e. the pc of a token (or ngram) is expressed as a function of "what has gone before" —
(i.e. conditional PD, w.r.t. "condition" being the known past corpus.

So it seems clear that context can probably cover all modifns of the PD₂ —

20 Restricting ourselves to context discovery would be same as limiting ourselves to using
a PD₂ to guide Lsearch — which is only probably not bad but is not certainly not

the most general way to solve problems (or do induction) — e.g. Quick Abort & Long
during Lsearch are techniques not covered by Modifns of PD₂: See 389.18 (about F/N of 30/89)
for descr. of "(very) long Lsearch" & possible ways to include them in PD₂ This would
be a possibl. (possibly noisy) extension of Normal use of PD₂ for Lsearch guidance.

It is an augmentation of the simpler Phase 1 model. We have, perhaps, 2 alternatives of augmenting the simpler Phase 1 model or using a better/larger

30 TSC w. the simpler Phase 1 model. (i.e. no long Search Lsearch)

Might be a good idea to list a large no. of "Contexts": Try to group them,
to arrange the list so that it's clear that there are only a few (or one) way(s)
to generate contexts.

Well, the most general context is the most general PD₂, i.e. I know how to generate all
possibl. PD₂'s. The only hitch here is that some contexts are "external":

We will not think of them as possibl. arguments for the conditional aspect
of PD₂.

About "External" contexts: If External Context can be supplied in R/W,
then again this is a "LARGE space"! ☹

SEE 42519 for
Good examples of
Context: "Wtd
names"

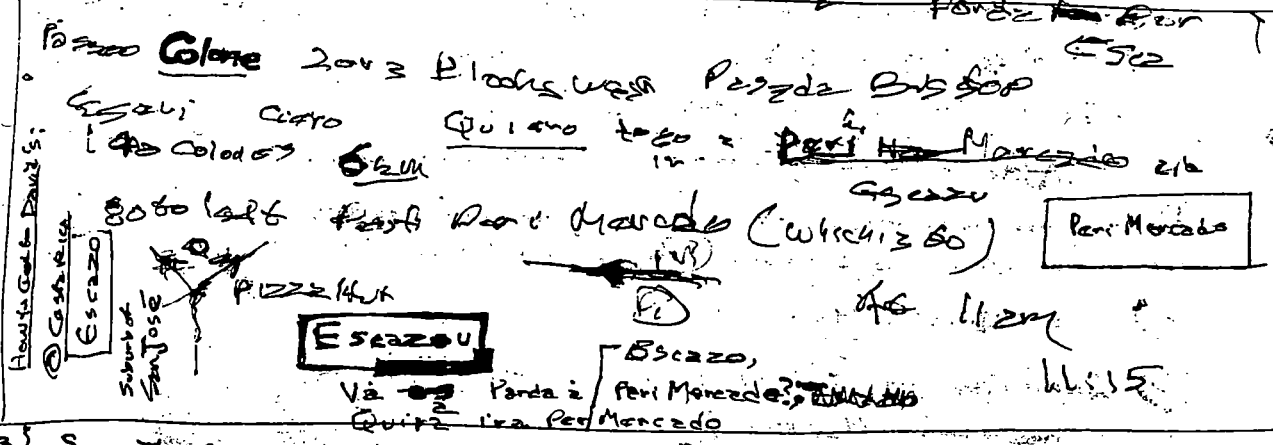
44

00: (399.08) ^{Spec} 399.08 is a good explanation of what PD₂ is about. It is not the most Genl. soln to a heuristic such problem, but it is meant to be a fairly good soln. (perhaps very et.) - Goodenough to Bootstrap to Phase 2. The idea of a PD₂ attempting to implement the ~~such~~ such problem of 399.08 (i.e. doing this w. a probability) is certainly not the best way to solve the problem: ~~Minimal differences!~~ (1) No long during such (2) No dump of "Experiments" (see 399.27-39).

To some extent 399.08 can learn better trials in a manner of the FN of 89!

[399.18-33] - by having the pd. being updated after each trial, since PD₂ should ~~can~~ depend on ~~some~~ info in the success & failures of any & all previous trials. The PD₂ seems to be getting very close to the methods of Phase 2. (I may, in fact, introduce Phase 2 in this way.)

Any Phase 1 tries to implement 399.08 in as simple a way as possible. Yet ~~good enough~~ good enough to get to Phase 2. (28)



28 (18) N6.03 So: The PD₂ idea is a PD that guides a modified L₁ such. It is not the most General Form of Heur. (won't be closer using PST's), but if we stick to that form we are restricting the form of our heur., but we (presumably) will not do badly. A good enuf (w. suitable Escazo) to get to Phase 2 (100-16 discuss possible arguments of the PD₂ approach to Phase 1. Depending on how good our PST's are & on how good PD₂ is, we may need to augment PD₂ in Phase ways to get to Phase 2.

37 So: Going back to 399.32 ft: Attempt to integrate "Content" = "Volendables" & "Escazo" tracks:

39 Token frequency is a context. For token α, (the no of times, α has occurred to corpus up to now) is a context, & we can compare with other token frequencies which would be contexts.

397.40: Well, Context can be defined in a very general way, involving Conditional PC's - which can in turn be defined by Codes.

A "definition" is a very general idea - defined by a code.

Frequency seems like a very cl. idea; but it seems to be very imp. in pc evaln.

So, again, 2 general types of PD: PD₁, = regys in ~~general~~ corpus:

PD₂ Regys ^{2 or 3} ~~is~~ successful & unsuccessful attempts to find codes for corpi. for Phase 1 → 400.00

So, one approach is to simply look at human methods of such for various problem types & see if I can find any good general characteristics.

T. lists of 398.00-.06 & 391.05-.40 2nd attempt in this direction.

Is full "Phase 2" a more general such method?

For some time I had a (Mistaken) idea that all hours could be expressed as modulus of PD₂ (The I didn't ~~understand~~ understand the PD₁, PD₂ dichotomy at that time). One main difficulty was that this didn't take "long during Lsrch" into account.

The Footnote of Sol 89 did take long during Lsrch into account, ~~but~~ - the ~~algorithm~~ such algm was quite close to a general PST, but I don't think I knew, at that time, just how to find & implement PST's of that kind.

Essentially, I had the PC of the codebook ~~to~~ (somehow optimized) function of all of TM's history (including "Process"). This is, indeed, close to 1. most general reasonable possibl., but I don't think I had figured

out any ways to ~~implement~~ find such "hours" (= PST's).

Adv. Int. sense of being able to implement "long during" such "T. FN of 89" was Very Good. 2? If does something a deficiency - but it would still work very well.

Q: Just how was the "FN of 89" to ~~be~~ implement "experiments" i.e. trials (or computations) made to gain info, not merely ~~to~~ to directly try to solve a problem? - "experiments" were used not be "codes". ~~Codebook~~ ~~to~~ solve a problem ~~was~~ attempts at solving of the problem ~~then~~ being worked on. (The I was certainly aware of the imp. of "experiments" for RTM.... Perhaps I had a way of ~~doing~~ the implementation. Now as a kind outgrowth of previous "experiments" of all kinds.

Anyway, the idea of PST ~~is~~ is very general & covers the "hours" of the "FN of 89" & doing experiments (as part of a PST).

00 397.40: Well, Context can be defined in a very general way, involving Conditional PC's - which can, in turn, be defined by Codes.

A definition is rather general idea - defined by a code.

Frequency seems like a very elo idea; but it seems to be very imp. in pc exam.

So, again, 2 general types of PD: PD₁: ways in ~~code~~ corpus:

PD₂ Repys ^{2nd ch} ~~success~~ successful (unsuccessful) attempts to find codes for Corp. for Phase 1 → 400.00

So, one approach is to simply look at human methods of such for various problem types & see if I can find any good general characteristics.

T. lists of 398.00-.06; 391.05-.40 2nd attempt in this direction.

Is full "Phase 2" the most general such method?

For some time I had (Mistaken) idea that all hours could be expressed as modulus of PD₂ (I didn't ~~understand~~ understand the PD₁ PD₂ dichotomy at that time). One uncertainty was that I didn't take "long during Lstch" into account.

Maybe not so wrong!
2000
436.11-.22

The Footnote of Sol 89 did take long during Lstch into account, ~~but~~ - the

~~algorithm~~ such algm was quite close to a general PST, but I don't think I knew, at that time, just how to find & implement PST's of that kind.

Essentially, I had the PC of it could be a ~~rather~~ (somehow optimized) function of all of TM's history (including "Yaces"). This is, indeed, close to the most general humble possibl., but I don't think I had figured

out any ways to ~~find~~ find such "hours" (≡ PST's).

Adv. Int. sense of being able to implement "long during such" T. FN of 89 was Very Good. 127 addresses something arbitrary - but it would still work very well.

Q: Just how was the "FN of 89" to ~~implement~~ implement "experiments" i.e. trials (or computations) made to gain info. not merely ~~to~~ to directly try to solve a problem? - The experiments ~~were~~ need not be "trials". ~~at that time~~ attempts at solving the problem ~~then~~ being worked on. (Tho I was certainly aware of the imp. of "experiments" for RTM..... Perhaps I had a way of ~~implementing~~ implementing

them as a kind outgrowth of previous experiences of all kinds.

Anyway, the idea of PST ~~is~~ is very general & covers the "hours" of the "FN of 89" is doing experiments (as part of a PST).

3PM



1 2 3 4 5 6 7
S S M To W

1905 = E. H. H. H.

10. 337.40 Analytical, it would be useful
In 396.07-11: The points: 1) Generalized "Context": The ^{general} definition of context
Context for tokens or for things larger than tokens.

- 2) Recogn. funcs.
- 3) Defns of pc's of Tokens.
- 4) r. SUMAC Concept.
- 5)

Since recogn. funcs are a form of context, (1;2) is (3) - and only 2 kinds of PD2. Sumac is related, but how to characterize t. relation? (32)

339104 SN Re: 390.35-391.04 (Context) In physics, Chem ("Natural Sciences") the no. of cones does not seem to grow much - The t. no. of cones is quite large; because used a very small no. of cones from which the others were derived, in a very h.c. manner. (So little of pc in t. m. N503) no. of cones, f. r. Re: Plus last: IN ANL (S200b), t. defns were quite sharp, yet we had a scaling "problem". So in t. "Natural Sciences" it is not at all clear that we don't have t. Usual "Scaling Effects"

SN On optzn problem v.s. Recult. lvy. As Best Model for Strong A.I.
T. oz problem seems simpler & if solved, could solve f. recult. lvy. problem.
If t. Recult. lvy. problem is solved (for a good trainer) t. oz problem could solve f. general optzn problem. - So both problems are solved f. other.
Q: One Q is: Is Oz really simpler (it seems to be) - T. division of f. system into sub-problems may have resulted in sub-probs that were not well defined (e.g. in T. Recognition function approach to QATM).
Another is T. general danger of RTM - uncertainty. The prob occurs in any h.c. level of problem solver as well, it seems to be more insidious in RTM's (i.e. the trainer will not be aware of what he himself really wants - T. goals given to RTM may have subgoals of unexpected side effects. - i.e. T. subgoals may be very bad for people in general.)

32 (08) So, I could do more on integrating ("R" Recognition) system & t. genl. context idea. This is closed up to Larry & Ken by 435.06-27 435.19 in Particular
Hvc., T. Defns. of pc's of Tokens seems to be a unrelated way to Get aprof (PD1) & Get parts of PD2 (parts) f. PD Rel Guides f. Struc.
Apparently, today is t. most General way to desc. regular exp. Context is a broad class of reg's (or parts of reg's): A status of tokens & pc's of tokens is a more imp class of reg's. But kitchen is as general as today.