

Definition of T.S. of Gops: 53.25

01. ((32.40 spec)): The t. T.S. is  $\Rightarrow$  There exists  $\exists$  seq. of Gops,  $G_{Op_i}$ , so that 1)  $G_{Op_i}$  can work from  $(P_0, A_0) \dots T_{Op_i}(P_i, A_i)$ .

(2)  $\forall \exists$ :  $G_{Op_{i+1}}$  can be obtained from  $G_{Op_i}$  by using addition of "information"  $P_{Op_{i+1}}$ . Also the cc (computing cost: usually in dollars) needed for  $G_{Op_{i+1}}$  to map  $X_{Op_i}$  into  $A_{Op_{i+1}}$  is  $CC_{Op_{i+1}}$ .

02. 3)  $\frac{CC_i}{PC_i} < A$ ; where  $A$  is a reasonable cc achievable with an available computer system. Here, more exactly,

From Usually,  $G_{Op_i} = G_{Op_{i+1}}$ : i.e. The same op is used to solve a sequence of problems.

Money owner willing to spend on soln.

04) say  $N$  is the average no. of probs solved by a gop  $G_{Op_i}$  before it fails & new  $G_{Op}$  must be found. Then I think we want ~~the total cost of solving N problems~~

What's the answer (so on. average using our soln. algos it

takes  $\sim CC = \frac{A}{N}$  to solve each problem  $\frac{A}{N}$  (if searching for solns takes more time than solving probs w.t. correct algos)

5) It may well be that there are other seqs. of Gops that can solve t. T.S. in t. manner of 1) ... 4), but there must be at least one seq. of

Gops that satisfy conditions 1), 2), 3), 4). (Also cc is vital).

T. range of  $G_{Op_i}$  is usually  $\ll PC_i$  of  $\frac{A}{N}$

6) Suppose that  $[G_{Op_i}]$  is a seq. of Gops that can solve t. T.S. in t. manner of 1) ... 4). Let  $[N_j]$  be a seq. of integers that tells how many probs are solved by each new gop before it fails. (so  $N = N_j$ )

We want  $\forall$  each  $N_j$  to be large enuf, so that when t. corresponding

Gop has solved  $N_j$  problems, there are ~~more~~ no more other Gops

$\frac{CC_i}{PC_i} < NA$  that have solved  $> N_j$  problems, ~~and hence~~

34 ④ There are few other Gops like this that can solve  $\geq N_j$  problems.

This is important if we are to retain the Gops that ~~have~~ have more soln. algos.

largest  $N_j$  value.

(perhaps easy to compute)

6) ② Is one essential feature of t. Alg.

6) ⑤ is a soft constraint. It gives a quantitative measure of soln. tot. except that it's violated

5.16.82 TS

Grant Operator

The Problem: To find a Gap w. max. pc., that can work all of t. problems. T.S. 135

.01 : An Algo. to work t. T.S. of 53.25 ( $\equiv 134.01 - .40$ ) . (i.e. to find this Gap).

.02 We start out with a Gop. ~~that can work all of t. problems~~

~~then~~ we try various modifications of it ~~that can work all of t. problems~~. These modifications have of various prob. of devn. — say  $p_{kA}$  for t.  $k^{\text{th}}$  poss. modifn.

.03 A set of / Gops;  $Gop_{0,k}$  is produced, is these are tried on  $\mathbb{Q}_0$ .

.04 We allow a cc of up to  $\mathbb{A} \cdot p_{kA}$  for the creation of  $Gop_{0,k}$ .

.05 We allow a cc of up to  $\mathbb{A} \cdot p_{kA}$  for the creation of  $Gop_{0,k}$ .

.06 and its soln. of  $\mathbb{Q}_0 \rightarrow A_0$ . If it takes more cc,

~~then Gop trial is a fail~~ or if it gets t. wrong  $A_0$ , then that

Gop is a failure, so go on to test t. next Gop.

Since  $\leq p_{kA} < 1$ ; t. total ~~cc~~ spent

on all these trials is  $\leq p_{kA} \cdot \mathbb{A} < \mathbb{A}$ .

We then take all t. Gops that solved  $\mathbb{Q}_0$  is we see how many

Subsequent problems each can solve — allowing  $\leq p_{kA} \cdot \mathbb{A}$  for each

~~problem.~~  $p_{kA}$  is certainly  $\geq$  cc needed to solve t. ~~first~~ problem.

~~Subsequent:~~ Any good Reasons for allowing  $\leq p_{kA}?$

.23 Say  $N_j$  is t. most problems any one of t. Gops can solve

before it fails. We then retain all Gops in memory that can

.25 Solve  $N_j$  probs. — say there are  $R$  such Gops.

(Solving for all Gops)

Say  $Q_s$  is t. first  $\mathbb{Q}$  for which all of these Gops has failed.

For each of these  $R$  Gops, we try modulus like .02 ff. ....

.24 (essentially, loop to .02 ~~but~~ but before that,  $Q_s \leftarrow Q_0$ ,  $A_s \leftarrow A_0$ )

{Total cc needed is  $\leq RA$

• If we run up against t.  $\mathbb{A} \cdot A$  but we can't solve even w. new Gop trials, then  $\mathbb{A} \leftarrow 2\mathbb{A}$  and loop to .09

Diffs of t. .01 ff algor.

If we assume that t. T.S. is of type 53.25 ( $134.01 - .40$ )

Then there are 2 major diffs:

.32 1) The number of  $\mathbb{A}$  Gops retained in memory tends to expand

exponentially w. t. no. of new ~~new~~ needed. ( $107.38, \frac{121.25}{121.30} \approx 2$  different (bad) effects).

.33 2) When a bunch of Gops are found ( $s$  in .23-.25) that can all

work t. next  $N_j$  probs, we must also test these Gops on all past

problems. The cc involved is  $\propto C^2$  (where  $C$  is t. no. of concepts

$\propto$  ( $\equiv$  new Gops) needed for t. T.S.) ( $107.07, 106.01 - .34$ ). Anything much more rapid than cc being linear in  $C$  is probably unacceptable.

Another ~~method~~ kind of activity that I haven't written much about recently ~~is~~ Part fits well into t. task net formalism is "Planning". It's a <sup>(sob)</sup> special way of breaking a problem into sub-goals. It ~~transforms~~ a problem into a task net.

On t. parti concept. of "Planning" I don't have ~~many~~ good ideas.

Newell may have something to say; Miller Galanter & Pribram? <sup>(bld MGP)</sup> (pp 177-194 maybe) relevant

A ~~rather general~~ method of making plans is by **Analogy**. Problem A has been solved. Problem X hasn't yet been solved. Prob A is analogous to X, in t. sense that there is a ~~transformation~~ function that transforms X to A. ~~Problem~~ There are elements in t. soln. of Problem A that have corresp. elements in t. domain of X. Those elements in X's space are used to ~~tentatively~~ sub-goals, to see if solving them would, indeed, solve X. If they would, then this set of goals in X's space would constitute a "**PLAN**".

T. larg. descn. of t. use of Analogy to devise Plans is very general.

— by itself, it is too general if it needs t. guidance obtained by experience to tell which kinds of analogy are good in this regard, & which ones can be used to plan what kinds of problems.

→ I suspect that most plans can be put in t. form of cribbing. <sup>10-20</sup>

I'm not sure, here!

Newell had t. idea of a "Plan space" as being a space in which t. problems to be solved were somehow simplified (say by leaving out certain ~~parts~~ parts of their descrs) — Then by solving t. ~~problem~~ (thus simplified) problem in t. plan space, one could guide t. corresp. path of work in t. true problem space. This view is a subclass of t. & more general Analogy method of <sup>10-20</sup>.

→ More generally a "plan" is any constraint on how one attempted to work on problem. ∵ it is part of t. descn. of t. method of working t. prob. However, to be a "plan" this descn. "morsel" must be shared by other problems. In this sense, this "morsel" would be "Named" if it were warranted — e.g. if it were used by enough probs. in t. past. As such, it would be assigned a PC & wt. for that PC. (so we'd know how to modify its PC. & when in t. future, we get more data on its % of success).

A "PLAN" is anything 2 diffnt problems solns (usually of different problems) have in common. A plan is somewhat reasonable if ~~is~~ sharable it  
shortens to ~~is~~ total dcrns of t. solns. of r. 2 problems.

"Sharing" the plan, gives it a pc. — but ordinarily, we also have some (conditional) prob. of t. effectiveness of a plan wrt. a particular problem. So, we "look at" t. problem using certain obs.. From the results of these obs, we obtain a (pc of that plan wrt that problem) ... a conditional probability of a name (of a plan).

The name of t. plan is part of t. dcrn. of t. soln. attempt. After choosing a plan, we fill in various other parts of t. dcrn. of t.  
soln. attempt ... as ~~demanded~~<sup>specified</sup> by the plan. Each of these parts of t. dcrn. has ~~is~~ its pc., & so at each point in our carrying out t. plan, we have to total pc to date ... as well as t. total cc to date. When, of course  $\frac{cc}{ps} > T_0$ , we drop that particular trial soln.

What all of t. forgg. looks like: say  $a, b, c, d$  are/elements in t. forgg. bag...  
 $a$  is simple enough, so that t. machine can code it w. direct L-srch.  
Plan to code  $b$ , we want to code for  $b$ , given a "good" code for  $a$  (not necessarily t. shortest).  
Then to code  $c$  we use " $a$  " "  $c$  " " good" codes for  $a \& b$ .

25 → A Plan is just another regularity type that has occurred in t. past. It is a particular peculiar regularity type! it involves, usually, t. structure of t. solns. to problems. As such, Plans are heuristics just like any others. They can be derived by TM<sub>2</sub> just like any other heuristic.

28 The nice thing about this: that it removes t. "PLAN" idea, from a allo device, to a well-integrated part of t. TM system.

This means that we can genz. plans, have t. same Genc for planning for other hours,

29 Perhaps Newishibas! That hours are perhaps only raggs, in which one has considered cc — i.e. for each body of data is  $T_0$  level, there is a set of raggs. We can regard a hour as t. cc of a ragg that would normally be modifiable because of excessively by cc. — This is perhaps closer to t. older version of what a hour is.

30 If 29 is correct, this does give us a more integrated view of TM's behavior.

⇒ At present, My impression of t. reason for t. TM<sub>1</sub>, TM<sub>2</sub> dichotomy:

8.11.81 ~~TS~~: Index-filts: List of Impt. ideas referred to in text.

140

1) Backtracking 94.01 — discn.

refd!: (27.23, 27); 130.35

2) Subconscious Mind; — discns! 81.32-34; 89.22 ~ 40; 91.01-22  
(on Neg. Results for Animals, Humans) memdiscn.

referred to: 140.20

3) Updating & Concept Learning — 1980: 80TS 36.014) Heuristics: What are they? 80TS {77.25 — .39} : 81TS 86.10 - 19  
53.01

- 5) PLANS: a) What are they  $\Rightarrow$  80TS 76.01 - 77.39 All 80TS  
**[** b) What are some imp. plans: ①  $\frac{65.07}{68.10 \text{ ff}}$ ? Is t. prob. in & set I can solve? If not,  
 can I limit it into a prob. I can solve. ② GPS: ~~80TS 73.01~~ 73.01  
 ③ Shytle's Saint & Multiple. ~ to GPS but maybe better. (SAINT: 80TS 74)  
 ④ Subdivision of corpus = into "parts" at same level  $\Rightarrow$  Hierarchically!

130.25 — 144.90

(5) ANALOGY.

(6) PEM &amp; CPM's as Plans: 148.37 - 38; 150.14;

6) ~~Maxmathed~~: 149.16R; also 122.017) Probabilistic Algorithms: Probabilistic Proofs of Theorems in Geometry 76.01  
 Laplace's hypoth in physics: Moddy'd form: 100.01.8) A kind of CB for ~~unproven~~ UMC's that may make Lsm easier! 76.08

Mainly TS file Mainly 81TS  
2.20.82 DM 5cc (8.11.81) ~ P110 for z-buffer index - like file  
(TS) Some Notes topics of interest in TS 9.3.81 to 8.30.81: 13.01.80 164.40 looks like  
Back tracking : 94.01, 96.38; 127.23, 92.30 { 164 was last page  
Nug. Recurrent (why bad) : 112.  
Subconscious Mind 91

Details of  $\Theta$  operator is D(O) 106.09

Gahl. Admin of TH work 113. is on what kind of paper to write.

An early mention of TS. Diffy that I worked on mainly ff: 116.25: (~ 160ff0)  
(Not yet solved). This was Pic loop v.s. B12 form of "Eval" operator  
by 1.9.83 I had solved it. Got 1 soln., then 2nd soln. was  
replaced by Pic "pluperfect TS." concept of 82TS 166.32 ~ 174.40

Mark McPhail 122; 149.16 R

(Calc. of rel. p.c. of 12.10., 23; 126.30; 127.17 (Loop) Step rule (No loop ( $\exists B^{12}$ ))

method of eval. of alg. expressions: 118.17 - ff: Note for Pic more efficient accurate work on this &

B: My version

•  $H(x/y) \neq H(x/y)$  my version! Variation of Chaitin: 131.01 ff

• Operator induction or extrapolating in F.O device 123.10.

Dividing up t. corpus for Length 139.01, 202.04 (~ 162.04)

• On "Plans": (Analogy pp 76, 77) This is in  $\leq$  TS! (See discussions on Plans)  
pw (Plan weighting) prob: 150.83, 158.00

Art on Evoln. of Intelligence: SA. 1965 pp 92-101

Associativity in RPN: 42.00 RPN intent 40.27 - 43.21

Assigning hype's to new dofs: 48.28 (also refs to 80TS).

Substitution: In learning of: ~ 51.01  
60.24 Bibliography

Problem #1 How to parse an expression so it's in a certain  
form: 65.27

Optimum order of probs for TS: 66.27

T. "Generalization" heuristic. 67.01

GPS Heuristics for undsol prob. Heurs: 73.01

UNARY FUNCTS, HYPC OBS; Boolean, Numerical 93.01  
constants: A List: 78.01  
out-Applied Numbers 122.25

Operator Induction in TS

117.19

• 81TS 159.30:

A kind of  
summary of  
• B12 v.s.  
Loop soln.  
problem.  
Also note →  
81TS 208.00-209.40

Reviews: 60.01

• 60.17: work on T.S. up to that pt. Bibliography  
various concepts worked on.

71.01ff An outline of an Alg. TS.

74.01: "On browsing thru 80TS: some implications  
refg.

• 79.01 List of probs (in TS) not solved or worked  
on: (few hrs, hr!)

99.01 Review of t. state of t. prob. at that pt.  
99.01 A list of different types of problems. (all equivalent cont'd to  
a certain type.)

~140 (8.11.81): A list of mft. dofs in TS file (Mainly 81TS)

110.01 A doc. of t. TS being worked on (in English)  
111.01

5.I.82 TS

0.4.: So +. Many diff'gs of f. methods of  $\alpha$  "w. doubling" (of 104.33) ( $\alpha$  also  $\gamma$  in 107.17)

107

① 105.10 : Doubling / May not just double cc of solving t. T.S.

② 115.18 : T. cc of solving t. T.S. is  $\propto$  +. Cost of f. most diff' concept

$\frac{121.01}{166.23} \leftarrow$  in particular it has  $\frac{167.02}{}$  This seems wasteful. One would like A to vary over t. corpus. (125.25 may be helpful!)

0.7  $\rightarrow$  ③ 106.4-34 The  $c^2$  effect: If we have to check out all trial GOP's  
most serious. But t. entire list of t. T.S., +. cc for each will be Solving t. prob's  $\propto$  (w.o. new  
+. c.c. of search) will be  $\propto c^2$ , so it will be, for large enough c (c = no. of concepts in t. T.S.)  
T. solutions may be related to that of ②: i.e. "localization" of prob. areas

1.7 ④ 105.02 If we add to refinement of  $\alpha$  to  $\alpha$ , so that f. prob's off. T.S.

need not be in exactly t. right order (see 98.03-20), then we must characterize t. params of this improvement before we can decide that a particular value of A has failed & we must  $\alpha < 2A$

5.2.82 2:18A

1.22 ⑤  $\rightarrow$  104.26 T. General Q of simply retaining only t. GOPs that solve f. most problems.  
This is an imp't. Q. that I may have some solns, but I'm not sure as to f. reasons!

$\rightarrow$  continues to 103.40

2.4 ⑥ 102.03 : 101.37-102.03 This is t. Q of just how TM can solve successive approx. prob's or Leach prob's. First take  $\alpha$  squintly  
different cc for each problem. 102.04 to n. 29 is a pretty big.  $\rightarrow$  120.31  
but I don't yet understand this or some imp't. assoc. Q's.  
— i.e. what various pc's are f. probabilities of.

1.8  $\rightarrow$  120.31  
seems to be reasonably clear

SN ⑦ (64.35ff) is an Elzn. of t. GOP concept. It has assoc. w. its some (apparently)  
modelling of t.  $\alpha$  & t. search techniques. I'll have to look into this after getting  
 $\alpha, \gamma$  straightened out better. Also see if C can lead to more prob's  
of elzn.

⑧ 97.35 : Re  $\gamma$  : can A.H. (spurious) solns. screw up t. use of  
 $\gamma$  to deal w. slightly m'ordant T.S.'s.  $\rightarrow$  see 113.38

3.5 ⑨ 110.10 T. needed concept may not always be the one that solves t. most problems  
immediately. See discussion following. Very imp't.  $\rightarrow$  see 114.23 for unif. w. ⑩

3.7 ⑩ 115.35 A certain concept needed is not directly included in t. T.S.

3.8 ⑪ 121.25 : Many ~~replicates~~  $\rightarrow$  11 codes for t. same GOP. Gives spuriously low cc thresholds. —  $\rightarrow$  also see 141.09  
Wastes enormous amt. of cc. T. no. of redundant codes is exponential in t. no. of concepts in t. corpus!

5.20.80:

Impf. Problems, projects:

- 1) Work on TM < perhaps organize past writings better.  
Make list of headings: Use files under a) headings  
b) Time of creation.
- 2) Work on good Computer System }  
}  $\beta$ ugs  
Hardware .....

- 3) Get ~~scribble~~ I.B. in good shape! Maintain it:

Get logging boys prob. straightened out.

- 4) Vacation travel.

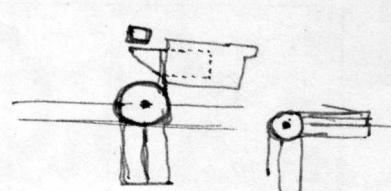
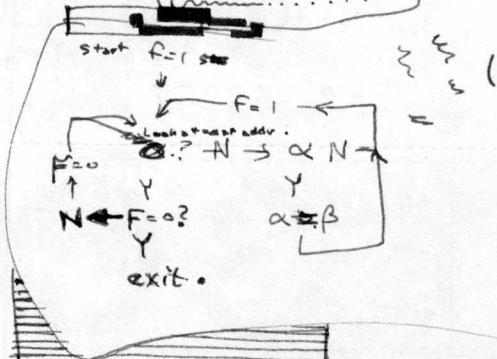
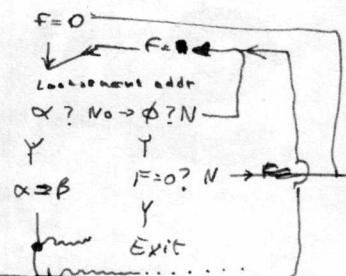
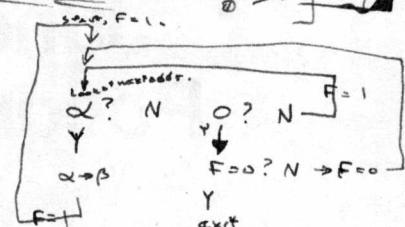
- 5) ~~scribble~~ Brain Drugs, Hyperbaric  $\text{O}_2$ , etc. [get good tests for changes.]

- 6) Clean up, straighten' out Comb. Apt.

- 7) Get new place to live.

- 8) Get General System to decide on what to work on next.

- 9) (part of <sup>1,2,6,8</sup> ~~scribble~~) Devise filing system for objects, documents.



6.5.80

Genl. Admin., Work Records, etc.

At all times, have in ~~the~~ rapid access storage. See current most impt. unsolved probs. I keep these in mind at all times.

This makes it poss'l. for me to work on them all t. time, & spend lots of pure "mental doodling" time on them. Also, it insures my spending more time on the truly <sup>imp</sup> problems.

Periodically (say every week), I can write a proj. report on t. ~~5~~ or ~~10~~ top problems, detailing poss'l. sub.probs., etc.

■ "Paging problems": (1) What in list of ~~the~~ most impt. projects, How to decide ~~on~~ on ordering (2) for "most impt problems" What to decide to put in Rapid access storage.

(3) In paper (or other files): I will have <sup>a pile of</sup> 100 or 200 pp. of readily accessible ~~the~~ pages. Which set pages shall I put there?

~~Recency of last use~~ Recency of last use is a factor. Also great importance.

N2482

Can Admin. (Plan).

I really need some means to get more work done on TM.

The main ~~sources~~ Energy Sinks for me!

- 1) Load
- 2) ~~me~~ 4 computers.

Time spent on me. could be useful toward TM, but so far, it hasn't been.

Bottlenecks in ~~me~~: ① reliability ② Disc. ③ cpu.

② Getting good, easily used, operating system: say FORTRAN

Also fast address to user hardware system

When I get O.S. working O.K., I will work on RTM in Accade

frame form. Discuss w. Tomas Witten whom is more advanced form.

Or, work on HR or SM (say, option ~~1~~, or any very new kind of)  
(but offer)

$$T = \tau_{AS}$$

$$E = 1 \text{ kV}$$

$$E_0 = \frac{1}{2} ET^2$$

$$= \frac{1}{2} \times 5 \times E(t) dt$$

$$= \frac{1}{2} \times \frac{1}{2} E^2$$

$$T = \frac{1}{2} \times K \cdot 10^{-12}$$

$$\alpha = \frac{20}{K \cdot 10^{12}} \\ = 2 \times 10^{10}$$

$$\frac{2 \times 10^{10}}{R \cdot 3 \cdot 4} \times K \cdot 10^{-24}$$

$$\frac{12}{10^{13} \times 10^{-24}} \\ = 10^{-12} \text{ volt!}$$

$$2V, .352$$

$$\sim 2.5 \text{ m.}$$

$$\frac{2 \times 10^{10}}{10^9} E t^4$$

$$t = 10^{-3}, E = 1000 \\ \rightarrow 1 \text{ volt out!}$$

Problems: Go to higher level goal.

Interim Soln:  $\boxed{4 \text{ hrs/d. on TM}}$  (again!) : Plan at

least 4 hrs./day.

→ Try to devise more comfortable work place & 26B.  
Get more comfortable chair.

Perhaps tie ~~foam~~ foam cushion to cherry wood chair..

3/19/83 TM:

.01 A general off.  $\pi$ -order alphabet for CBI — so that  $\textcircled{2}$  T.V.H.  
method ~~implies~~ gives ~~a~~ more like a continuous distribution of  
probs.

Use a large alphabet w. r symbols. Assoc w. t-<sup>2</sup> symbol's F. prob p<sub>i</sub>;  $\sum p_i = 1$ .

One apparently good way to do this: use  $p_i = \epsilon$  for  $i = k$ ,  
w. small  $\epsilon$ . Small  $p_n = 1/(n-1)\epsilon$ ;  $\log_2 p_n$  is very  
close to  $\phi - i \approx -(n-1)\epsilon$ ;  $\log_2 (n-1)\epsilon \approx$   $\log_2 n$  due  
quantization error of  $\pi$ -alphabet.

I guess  $r=2$  is workable!

Now: How to map the original corpus into the  
new alphabet? or into strings in t. new alphabet

(SN)

42588 Haven't been able to figure out meaning of  $\phi$  !

on 817

Maybe codification  
of TH  $\phi$ 's  
not reliable!